



iSeries

WebSphere® Development Studio C/C++ Language Reference

SC09-4815-00







iSeries

WebSphere® Development Studio C/C++ Language Reference

SC09-4815-00

Note!

Before using this information and the product it supports, be sure to read the general information under "Notices" on page 339.

First Edition (May 2001)

This edition applies to Version 5, Release 1, Modification 0, of IBM WebSphere Development Studio for iSeries (program 5722-WDS), ILE C/C++ compilers, and to all subsequent releases and modifications until otherwise indicated in new editions. This edition replaces ILE C for AS/400 Language Reference (SC09–2711–01) and VisualAge C++ for OS/400 C++ Language Reference (SC09–2121–00).

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About This Reference

The C language documented in this reference is consistent with the one described in *Programming languages* – *C* (ISO/IEC 9899:1990).

The C++ language documented in this reference is consistent with the one described in *Programming languages* – C++ (ISO/IEC 14882:1998).

Highlighting Conventions

Bold Identifies commands, keywords, files, directories, and other items

whose names are predefined by the system.

Italics Identify parameters whose actual names or values are to be

supplied by the programmer. *Italics* are also used for the first mention of new terms that are defined in the glossary.

Example Identifies examples of specific data values, examples of text similar

to what you might see displayed, examples of portions of program code, messages from the system, or information that you should

actually type.

Examples are intended to be instructional and do not attempt to minimize run time, conserve storage, or check for errors. The examples do not demonstrate all of the possible uses of C and C++ language constructs. Some examples are only code fragments and will not compile without additional code.

How to Read the Syntax Diagrams

• Read the syntax diagrams from left to right, from top to bottom, following the path of the line.

The >>— symbol indicates the beginning of a command, directive, or statement.

The → symbol indicates that the command, directive, or statement syntax is continued on the next line.

The —— symbol indicates that a command, directive, or statement is continued from the previous line.

The → symbol indicates the end of a command, directive, or statement.

Diagrams of syntactical units other than complete commands, directives, or statements start with the ►── symbol and end with the ──► symbol.

Note: In the following diagrams, statement represents a C or C++ command, directive, or statement.

Required items appear on the horizontal line (the main path).

▶>—statement—required item—

Optional items appear below the main path.

Reading the Syntax Diagrams



• If you can choose from two or more items, they appear vertically, in a stack. If you must choose one of the items, one item of the stack appears on the main path.



If choosing one of the items is optional, the entire stack appears below the main path.



The item that is the default appears above the main path.

$$\rightarrow$$
—statement— $alternate_item$ — \rightarrow

An arrow returning to the left above the main line indicates an item that can be repeated.

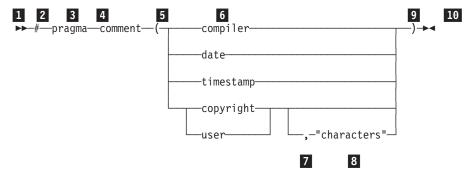


A repeat arrow above a stack indicates that you can make more than one choice from the stacked items, or repeat a single choice.

- Keywords appear in nonitalic letters and should be entered exactly as shown (for example, extern).
 - Variables appear in italicized lowercase letters (for example, identifier). They represent user-supplied names or values.
- If punctuation marks, parentheses, arithmetic operators, or other such symbols are shown, you must enter them as part of the syntax.

The following syntax diagram example shows the syntax for the #pragma comment directive. See "Pragma Directives (#pragma)" on page 176 for information on the **#pragma** directive.

Reading the Syntax Diagrams



- 1 This is the start of the syntax diagram.
- 2 The symbol # must appear first.
- 3 The keyword pragma must appear following the # symbol.
- 4 The name of the pragma comment must appear following the keyword pragma.
- 5 An opening parenthesis must be present.
- **6** The comment type must be entered only as one of the types indicated: compiler, date, timestamp, copyright, or user.
- **7** A comma must appear between the comment type copyright or user, and an optional character string.
- **8** A character string must follow the comma. The character string must be enclosed in double quotation marks.
- **9** A closing parenthesis is required.
- 10 This is the end of the syntax diagram.

The following examples of the **#pragma comment** directive are syntactically correct according to the diagram shown above:

```
#pragma comment(date)
#pragma comment(user)
#pragma comment(copyright, "This text will appear in the module")
```

Reading the Syntax Diagrams

Chapter 1. Scope and Linkage

Scope

The area of the code where an identifier is visible is referred to as the *scope* of the identifier. The following are the kinds of scopes:

- Local
- Function
- Function prototype
- Global or Global namespace
- Namespace
- C++ Class

The scope of a name is determined by the location of the name's declaration.

In all declarations the identifier is in scope before the initializer. The following example demonstrates this:

```
int x;
void f() {
  int x = x;
}
```

The x declared in function f() has local scope, not global namespace scope.

A function name that is first declared as a friend of a class is in the innermost nonclass scope that encloses the class. If the friend function is a member of another class, it has the scope of that class. The scope of a class name first declared as a friend of a class is the first nonclass enclosing scope.

The implicit declaration of the class is not visible until another declaration of that same class is seen.

RELATED REFERENCES

- "Local Scope"
- "Function Scope" on page 2
- "Function Prototype Scope" on page 2
- "Global Scope" on page 2
- "Chapter 10. Namespaces" on page 177
- "Class Scope" on page 3

Local Scope

A name has *local scope* or *block scope* if it is declared in a block. A name with local scope can be used in that block and in blocks enclosed within that block, but the name must be declared before it is used. When the block is exited, the names declared in the block are no longer available.

Parameter names for a function have the scope of the outermost block of that function. Also if the function is declared and not defined, these parameter names have function prototype scope.

SEC++ If a local variable is a class object with a destructor, the destructor is called when control passes out of the block in which the class object was constructed.

When one block is nested inside another, the variables from the outer block are usually visible in the nested block. However, if the declaration of a variable in a nested block has the same name as a variable that is declared in an enclosing block, the declaration in the nested block hides the variable that was declared in the enclosing block. The original declaration is restored when program control returns to the outer block. This is called block visibility.

RELATED REFERENCES

- "Block Statement" on page 145
- "Function Prototype Scope"
- "Destructors" on page 266

Function Scope

The only type of identifier with function scope is a label name. A label is implicitly declared by its appearance in the program text and is visible throughout the function that declares it.

A label can be used in a goto statement before the actual label is seen.

RELATED REFERENCES

• "Labels" on page 143

Function Prototype Scope

In a function declaration (also called a *function prototype*) or in any function declarator — except the declarator of a function definition — parameter names have function prototype scope. Function prototype scope terminates at the end of the nearest enclosing function declarator.

RELATED REFERENCES

"Function Declarations" on page 124

Global Scope

A name has *global scope* if the identifier's declaration appears outside of any block. A name with global scope and internal linkage is visible from the point where it is declared to the end of the translation unit. (A translation unit is a source code file after preprocessing with include files.)

■C++■ A name has *global namespace scope* if the identifier's declaration appears outside of all blocks and classes. A name with global namespace scope and internal linkage is visible from the point where it is declared to the end of the translation unit.

A name with global (namespace) scope is also accessible for the initialization of global variables. If that name is declared extern, it is also visible at link time in all object files being linked.

RELATED REFERENCES

• "Chapter 10. Namespaces" on page 177

- "Internal Linkage" on page 4
- "extern Storage Class Specifier" on page 25

Class Scope

C++ A name declared within a member function hides a declaration of the same name whose scope extends to or past the end of the member functionÆs class.

The scope of a declaration that extends to or past the end of a class definition also extends to the regions defined by its member definitions and any portion of the declarator part of such definitions which follows the identifier.

The name of a class member has *class scope* and can only be used in the following

- In a member function of that class
- In a member function of a class derived from that class
- After the . (dot) operator applied to an instance of that class
- After the . (dot) operator applied to an instance of a class derived from that class, as long as the derived class does not hide the name
- After the -> (arrow) operator applied to a pointer to an instance of that class
- After the -> (arrow) operator applied to a pointer to an instance of a class derived from that class, as long as the derived class does not hide the name
- After the :: (scope resolution) operator applied to the name of a class
- After the :: (scope resolution) operator applied to a class derived from that class.

RELATED REFERENCES

- "Chapter 12. Classes" on page 199
- "Member Functions" on page 211
- "Derivation" on page 233
- "Dot Operator ." on page 81
- "Arrow Operator ->" on page 82
- "C++ Scope Resolution Operator ::" on page 77
- "Scope of Class Names" on page 203

Name Hiding

If a class name or enumeration name is in scope and not hidden it is visible. A class name or enumeration name can be hidden by an explicit declaration of that same name — as an object, function, or enumerator — in a nested declarative region or derived class. The class name or enumeration name is hidden wherever the object, function, or enumerator name is visible. This process is referred to as name hiding.

In a member function definition, the declaration of a local name hides the declaration of a member of the class with the same name. The declaration of a member in a derived class hides the declaration of a member of a base class of the same name.

Suppose a name x is a member of namespace A, and suppose that the members of namespace A are visible in a namespace B because of a using declaration. A declaration of an object named x in namespace B will hide A::x. The following example demonstrates this:

#include <iostream> #include <typeinfo> using namespace std;

```
namespace A {
  char x;
};

namespace B {
  using namespace A;
  int x;
};

int main() {
  cout << typeid(B::x).name() << endl;
}</pre>
```

The following is the output of the above example:

int

The declaration of the integer x in namespace B hides the character x introduced by the using declaration.

RELATED REFERENCES

- "Chapter 12. Classes" on page 199
- "Member Functions" on page 211
- "Member Scope" on page 213
- "Chapter 10. Namespaces" on page 177
- "Using Directive" on page 182

Program Linkage

Linkage determines whether identifiers that have identical names refer to the same object, function, or other entitiy, even if those identifiers appear in different translation units. (A *translation unit* is a source code file after preprocessing with include files.) The linkage of an identifier depends on how it was declared. There are three kinds of linkage:

- If an identifier x has *internal linkage*, every appearance of x within one translation unit refers to the same entity.
- If an identifier x has *external linkage*, every appearance of x across any translation unit (of the same program) refers to the same entity.
- If an identifier x has *internal linkage*, every appearance of x refers to a unique entity.

RELATED REFERENCES

- "Internal Linkage"
- "External Linkage" on page 5
- "No Linkage" on page 6

Internal Linkage

The following kinds of identifiers have internal linkage:

- Objects, references, functions or function templates explicitly declared static.
- Objects or references declared in namespace scope (or global scope in C) with the specifier **const** and neither explicitly declared **extern**, nor previously declared to have external linkage.
- Data members of a anonymous union.
- Identifiers declared in the unnamed namespace.

A function declared inside a block will usually have external linkage. An object declared inside a block will usually have external linkage if it is specified extern. If a variable that has static storage is defined outside a function, the variable has internal linkage and is available from the point where it is defined to the end of the current translation unit.

A class that has no static members or noninline member functions, and that has not been used in the declaration of an object or function or class is local to its compilation unit.

If the declaration of an identifier has the keyword **extern** and if a previous declaration of the identifier is visible at namespace or global scope, the identifier has the same linkage as the first declaration.

RELATED REFERENCES

- "static Storage Class Specifier" on page 28
- "volatile and const Qualifiers" on page 50
- "extern Storage Class Specifier" on page 25
- "Global Scope" on page 2
- "Anonymous Unions" on page 45
- "Unnamed Namespaces" on page 180

External Linkage

The following kinds of identifiers with namespace scope (or global scope in C) have external linkage:

- An object, reference, or function unless it has internal linkage.
- C++ A named class or enumeration.
- An enumerator of an enumeration that has external linkage.
- · A template, unless it is a function tempalte with internal linkage
- A namespace, unless it is decalred in an unnamed namespace

The following also have external linkage:

- Member functions, static data members, classes, or enumerations if the class that they belong to has external linkage.
- Identifiers with namespace scope or local scope that have the keyword extern in their declarations.
- Static class members and noninline member functions

If a previous declaration of an object or function is visible in an enclosing scope, the identifier has the same linkage as the first declaration. For example, a variable or function that is first declared with the keyword static and later declared with the keyword extern has internal linkage. However, a variable or function that has no linkage and later declared with a linkage specifier will have the linkage you have specified.

RELATED REFERENCES

- "Scope" on page 1
- "static Storage Class Specifier" on page 28
- "extern Storage Class Specifier" on page 25
- "Chapter 12. Classes" on page 199
- "Enumerations" on page 46
- "typedef" on page 29
- "Chapter 10. Namespaces" on page 177

Program Linkage

- "Static Members" on page 218
- "Inline Functions" on page 141

No Linkage

The following kinds of identifiers have no linkage:

- Names that have neither external or internal linkage
- Names declared in local scopes (with exceptions like certain entities declared with the extern keyword)
- Identifiers that do not represent an object or a function, including labels, enumerators, **typedef** names that refer to entities with no linkage, type names, function parameters, and template names

You cannout use a name with no linkage to declare an entity with linkage. For example, you cannot use the name of a class or enumeration or a typedef name referring to an entity with no linkage to declare an entity with linkage. The following example demonstrates this:

```
int main() {
  struct A { };
// extern A a1;
  typedef A myA;
// extern myA a2;
}
```

The compiler will not allow the declaration of a1 with external linkage. Class A has no linkage. The compiler will not allow the declaration of a2 with external linkage. The typedef name a2 has no linkage because A has no linkage.

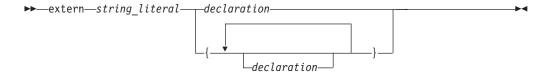
RELATED REFERENCES

- "Program Linkage" on page 4
- · "Scope" on page 1
- "typedef" on page 29
- "static Storage Class Specifier" on page 28
- "extern Storage Class Specifier" on page 25

Linkage Specifications — Linking to Non-C++ Programs

You can link C++ object modules to object modules produced using other source languages such as C by using a *linkage specification*.

The syntax is:



The *string_literal* is used to specify the linkage associated with a particular function. For example:

```
/**
    ** This example illustrates linkage specifications
    **/
extern "C" int printf(const char*,...);
```

Linkage Specifications

```
void main(void)
      printf("hello\n");
```

Here the *string_literal*, "C", tells the compiler that the routine printf(const char*,...) is a C library function.

Note: This example is not guaranteed to work on all platforms. The only safe way to declare a function from the C library in a C++ program is to include the appropriate header. In this example you would substitute the line of code with **extern** with the following line:

```
#include <stdio.h>
```

String literals used in linkage specifications should be considered as case-sensitive.

All platforms support the following values for string_literal

"C++" Unless otherwise specified, objects and functions

have this default linkage specification.

"C" Indicates linkage to a C procedure

Linkage Specifications

Chapter 2. Lexical Elements

This document contains discussions of the basic lexical elements and conventions of the C and C++ programming languages: tokens, character sets, comments, identifiers, and literals.

Tokens

Source code is treated during preprocessing and compilation as a sequence of *tokens*. A token is the smallest independent unit of meaning of a program as defined by the compiler. There are five different types of tokens:

- Identifiers
- Keywords
- Literals
- Operators
- Punctuators

Adjacent identifiers, keywords and literals must be separated with white space. Other tokens should be separated by white space to make the source code more readable. White space includes blanks, horizontal and vertical tabs, new lines, form feeds and comments.

RELATED REFERENCES

- "Identifiers" on page 14
- "Keywords" on page 15
- "Literals" on page 16
- "Chapter 5. Expressions and Operators" on page 71
- "Source Program Character Set"

Source Program Character Set

The following lists the basic *source character set* that must be available at both compile and run time:

• The uppercase and lowercase letters of the English alphabet

```
a b c d e f g h i j k l m n o p q r s t u v w x y z A B C D E F G H I J K L M N O P Q R S T U V W X Y Z
```

• The decimal digits 0 through 9

```
0 1 2 3 4 5 6 7 8 9
```

• The following punctuators (A *punctuator* is a character that has syntactic and semantic meaning, but does not specify an operation that produces a value. Depending on the context, a punctuator can also be an operator.):

```
! " # % & ' ( ) * + , - . / :
; < = > ? [ \ ] _ { } ~
```

- The caret () character in ASCII (bitwise exclusive OR symbol) or the equivalent not (¬) character in EBCDIC
- The split vertical bar (|) character in ASCII, which may be represented by the vertical bar (|) character on EBCDIC systems
- The space character
- The control characters representing new-line, horizontal tab, vertical tab, and form feed, and end of string (NULL character)

RELATED REFERENCES

"Chapter 5. Expressions and Operators" on page 71

Escape Sequences

You can represent any member of the execution character set by an escape sequence. They are primarily used to put nonprintable characters in character and string literals. For example, you can use escape sequences to put such characters as tab, carriage return, and backspace into an output stream.



An escape sequence contains a backslash (\) symbol followed by one of the escape sequence characters or an octal or hexadecimal number. A hexadecimal escape sequence contains an x followed by one or more hexadecimal digits (0-9, A-F, a-f). An octal escape sequence uses up to three octal digits (0-7). The value of the hexadecimal or octal number specifies the value of the desired character or wide character.

Note: The line continuation sequence (\ followed by a new-line character) is not an escape sequence. It is used in character strings to indicate that the current line continues on the next line.

The escape sequences and the characters they represent are:

Escape Sequence	Character Represented
\a	Alert (bell, alarm)
\ b	Backspace
\f	Form feed (new page)
\n	Newline
\r	Carriage return
\t	Horizontal tab
\v	Vertical tab
\ 1	Single quotation mark
\ "	Double quotation mark
\?	Question mark
\\	Backslash

The value of an escape sequence represents the member of the character set used at run time. Escape sequences are translated during preprocessing. For example, on a system using the ASCII character codes, the value of the escape sequence \x56 is the letter V. On a system using EBCDIC character codes, the value of the escape sequence \xE5 is the letter V.

Use escape sequences only in character constants or in string literals. An error message is issued if an escape sequence is not recognized.

In string and character sequences, when you want the backslash to represent itself (rather than the beginning of an escape sequence), you must use a \\ backslash escape sequence. For example:

```
cout << "The escape sequence \\n." << endl;</pre>
```

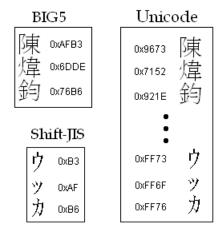
This statement results in the following output: The escape sequence \n .

The Unicode Standard

The Unicode Standard is a standardized character code designed to encode international texts for display and storage. It uses a unique 16-bit value to represent each individual character. The Unicode standard includes the following:

- Alphabets used in Europe, Africa, and Asia
- Standard characters from China, Japan, Korea, and Taiwan
- Mathematical operators
- Technical symbols

The following diagram illustrates how Unicode assigns a unique 16-bit value to each character:



0xAFB3 represents the 原 character in BIG5 and the characters ウッ in Shift-JIS.

Unicode assigns each character with a unique code point. In this case 原 is assigned with 0x9673 and ウツ with 0xFF73 and 0xFF6F.

Although the 16-bit architecture of Unicode can handle more than 65,000 different characters, the Unicode Standard can extend to handle an additional one million characters by the surrogate extension mechanism. This mechanism uses two 16-bit values to represent one character. The Unicode Standard has not used any of these surrogates. (The current standard contains 38,885 characters.)

The Unicode Standard lets you dynamically compose accented characters. In the Unicode Standard, a character and an accent are separate characters. In other character encodings such as ASCII, you select from a set of accented characters.

The Unicode Standard supports bidirectional ordering of languages. Bidirectional language ordering occurs when a script uses two or more languages with different dominant directions. For example, a script would have bidirectional language ordering if it mixes Arabic (which reads from right-to-left) with Greek (which reads from left-to-right). The Unicode Standard includes characters that specify a change of direction.

You can represent Unicode characters in your program by using one of the two following forms, where each *x* is a hexadecimal digit:

\uxxxx \Uxxxxxxxx

The first form, \uxxxx, represents a Unicode character that uses one 16-bit value. The second form, \Uxxxxxxxx, represents a character that uses two 16-bit values.

Trigraph Sequences

Some characters from the C and C++ character set are not available in all environments. You can enter these characters into a C source program using a sequence of three characters called a *trigraph*. The trigraph sequences are:

??=	#	pound sign
??([left bracket
??)]	right bracket
??<	{	left brace
??>	}	right brace
??/	\	backslash
??'	^	caret
??!		vertical bar
??-	~	tilde

The preprocessor replaces trigraph sequences with the corresponding single-character representation.

Comments

A comment is text replaced during preprocessing by a single space character; the compiler therefore ignores all comments.

There are two kinds of comments:

- The /* (slash, asterisk) characters, followed by any sequence of characters (including new lines), followed by the */ characters. This kind of comment is commonly called a C-style comment.
- The // (two slashes) characters followed by any sequence of characters. A new line not immediately preceded by a backslash terminates this form of comment. This kind of comment is commonly called a single-line comment or a C++ comment. A C++ comment can span more than one physical source line if it is joined into one logical source line with line-continuation (\) characters. The backslash character can also be represented by a trigraph.

You can put comments anywhere the language allows white space. You cannot nest C-style comments inside other C-style comments.

Multibyte characters can also be included within a comment.

Note: The /* or */ characters found in a character constant or string literal do not start or end comments.

In the following program, the second printf() is a comment: #include <stdio.h>

```
int main(void)
```

```
printf("This program has a comment.\n");
  /* printf("This is a comment line and will not print.\n"); */
return 0;
}
```

Because the second printf() is equivalent to a space, the output of this program is:

This program has a comment.

Because the comment delimiters are inside a string literal, printf() in the following program is not a comment.

```
#include <stdio.h>
int main(void)
{
    printf("This program does not have \
/* NOT A COMMENT */ a comment.\n");
return 0;
}
```

The output of the program is:

This program does not have /* NOT A COMMENT */ a comment.

You cannot nest C-style comments. Each comment ends at the first occurrence of */.

In the following example, the comments are highlighted:

```
/* A program with nested comments. */
```

```
#include <stdio.h>
int main(void)
{
    test_function();
    return 0;
}
int test_function(void)
{
    int number;
    char letter;
/*
number = 55;
letter = 'A';
/* number = 44; */
*/
return 999;
}
```

In test_function, the compiler reads the first /* through to the first */. The second */ causes an error. To avoid commenting over comments already in the source code, you should use conditional compilation preprocessor directives to cause the compiler to bypass sections of a program. For example, instead of commenting out the above statements, change the source code in the following way:

```
/* A program with conditional compilation to avoid nested comments. */
#define TEST_FUNCTION 0
#include <stdio.h>
int main(void)
{
   test_function();
```

Comments

```
return 0;
int test_function(void)
   int number;
   char letter;
 #if TEST FUNCTION
  number = 55;
  letter = 'A';
  /*number = 44;*/
 #endif /*TEST_FUNCTION */
```

You can nest single line comments within C-style comments. For example, the following program will not output anything:

```
#include <stdio.h>
int main(void)
  printf("This line will not print.\n");
  // This is a single line comment
  // This is another single line comment
  printf("This line will also not print.\n");
  return 0;
}
```

RELATED REFERENCES

"Trigraph Sequences" on page 12

Identifiers

Identifiers consist of an arbitrary number of letters or digits. They provide names for the following language elements:

- Functions
- Objects
- Labels
- Function parameters
- Macros and macro parameters
- Typedefs
- Enumerated types and emumerators
- Classes and class members
- Templates
- Template parameters
- Namespaces
- Struct and union names

An identifier has the form:



Case Sensitivity and Special Characters in Identifiers

The compiler distinguishes between uppercase and lowercase letters in identifiers. For example, PROFIT and profit represent different identifiers.

Avoid creating identifiers that begin with an underscore () for function names and variable names.

The first character in an identifier must be a letter. The _ (underscore) character is considered a letter; however, identifiers beginning with an underscore are reserved by the compiler for identifiers at global namespace scope.

Identifiers that contain two consecutive underscores or begin with an underscore followed by a capital letter are reserved in all contexts.

You should always include the appropriate headers when using standard library functions.

Keywords

Keywords are identifiers reserved by the language for special use. Although you can use them for preprocessor macro names, it is poor programming style. Only the exact spelling of keywords is reserved. For example, auto is reserved but AUTO is not. The following lists the keywords common to both the C and C++ languages:

double	long	typedef
else	register	union
enum	return	unsigned
extern	short	void
float	signed	volatile
for	sizeof	while
goto	static	
if	struct	
int	switch	
	else enum extern float for goto if	else register enum return extern short float signed for sizeof goto static if struct

The C++ language also reserves the following keywords:

asm	export	private	true
bool	false	protected	try
catch	friend	public	typeid
class	inline	reinterpret_cast	typename
const_cast	mutable	static_cast	using
delete	namespace	template	virtual
dynamic_cast	new	this	wchar_t
explicit	operator	throw	

In addition, the following alternative representations of operators and punctuators are also reserved in C and C++:

and	bitor	not_eq	xor	
and_eq	compl	or	xor_eq	
bitand	not	or_eq		

Identifiers

Alternative Tokens

C and C++ provide alternative representations for some operators and puctuators. The following table lists the operators and punctuators and their alternative representation:

Operator or Punctuator	Alternative Representation
{	<%
}	%>
[<:
]	:>
#	%:
##	%:%:
&&	and
	bitor
	or
•	xor
	comp1
&	bitand
&=	and_eq
=	or_eq
_=	xor_eq
!	not
!=	not_eq
	-

Literals

A literal does not change its value while the program is running. The value of any literal must be in the range of representable values for its type. The following are the available types of literals:

- Integer
- Character
- Floating-point
- String
- C++ Boolean

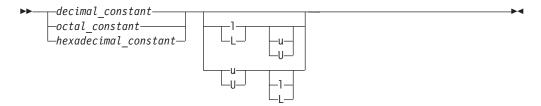
The C language uses the term *constant* in place of the term *literals*.

RELATED REFERENCES

- "Integer Literals" on page 17
- "Floating-Point Literals" on page 18
- "Character Literals" on page 20
- "String Literals" on page 20
- "Boolean Literals" on page 22

Integer Literals

Integer literals can represent decimal, octal, or hexadecimal values. They are numbers that do not have a decimal point or an exponential part. However, an integer literal may have a prefix that specifies its base, or a suffix that specifies its type.



The data type of an integer literal is determined by its form, value, and suffix. The following table lists the integer literals and shows the possible data types. The smallest data type that can represent the constant value is used to store the constant.

Possible Data Types
int, long int, unsigned long int
int, unsigned int, long int, unsigned long int
int, unsigned int, long int, unsigned long int
unsigned int, unsigned long int
long int, unsigned long int
unsigned long int
long long int, unsigned long long int
unsigned long long int

A plus (+) or minus (-) symbol can precede an integer literal. The operator is treated as a unary operator rather than as part of the literal.

RELATED REFERENCES

- "Decimal Integer Literals"
- "Hexadecimal Decimal Literals" on page 18
- "Octal Decimal Literals" on page 18
- "Integer Variables" on page 34

Decimal Integer Literals

A decimal integer literal contains any of the digits 0 through 9. The first digit cannot be 0.



Integer literals beginning with the digit 0 are interpreted as an octal integer literal rather than as a decimal integer literal.

Literals

The following are examples of decimal literals:

```
485976
-433132211
+20
5
```

A plus (+) or minus (-) symbol can precede the decimal integer literal. The operator is treated as a unary operator rather than as part of the literal.

Hexadecimal Decimal Literals

A hexadecimal decimal literal begins with the 0 digit followed by either an x or X, followed by any combination of the digits 0 through 9 and the letters a through f or A through F. The letters A (or a) through F (or f) represent the values 10 through 15, respectively.



The following are examples of hexadecimal integer literals:

0x3b24 0XF96 0x21 0x3AA 0X29b 0X4bD

Octal Decimal Literals

An octal decimal literal begins with the digit 0 and contains any of the digits 0 through 7.



The following are examples of octal decimal literals:

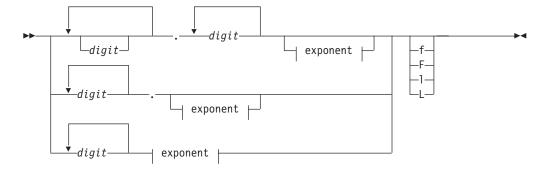
0125 034673 03245

Floating-Point Literals

A *floating-point literal* consists of the following:

- an integral part
- · a decimal point
- · a fractional part
- · an exponent part
- · an optional suffix

Both the integral and fractional parts are made up of decimal digits. You can omit either the integral part or the fractional part, but not both. You can omit either the decimal point or the exponent part, but not both.



Exponent:



The magnitude range of **float** is approximately 1.2e-38 to 3.4e38. The magnitude range of **double** or **long double** is approximately 2.2e-308 to 1.8e308. If a floating-point constant is too large or too small, the result is undefined by the language.

The suffix f or F indicates a type of float, and the suffix l or L indicates a type of long double. If a suffix is not specified, the floating-point constant has a type double.

A plus (+) or minus (-) symbol can precede a floating-point literal. However, it is not part of the literal; it is interpreted as a unary operator.

The following are examples of floating-point literals:

Floating-Point Constant	Value
5.3876e4	53,876
4e-11	0.0000000004
1e+5	100000
7.321E-3	0.007321
3.2E+4	32000
0.5e-6	0.0000005
0.45	0.45
6.e10	6000000000

When you use the **printf** function to display a floating-point constant value, make certain that the **printf** conversion code modifiers that you specify are large enough for the floating-point constant value.

RELATED REFERENCES

- "Floating-Point Variables" on page 33
- "Unary Expressions" on page 86

Character Literals

A character literal contains a sequence of characters or escape sequences enclosed in single quotation mark symbols, for example 'c'. A character literal may be prefixed with the letter L, for example L'c'. A character literal without the L prefix is an ordinary character literal or a narrow character literal. A character literal with the L prefix is a wide character literal. An ordinary character literal that contains more than one character or escape sequence (excluding single quotes ('), backslashes (\) or new-line characters) is a multicharacter literal.

Character literals have the following form:



At least one character or escape sequence must appear in the character literal. The characters can be from the source program character set, excluding the single quotation mark, backslash and new-line symbols. A character literal must appear on a single logical source line.

A character literal that contains only one character has type **char**, which is an integral type. A multicharacter literal has type int.

A character literal has type int.

A wide character literal has type wchar_t. A multicharacter literal has type int.

The value of a narrow or wide character literal containing a single character is the numeric representation of the character in the character set used at run time. The value of a narrow or wide character literal containing more than one character or escape sequence is implementation-defined.

You can represent the double quotation mark symbol by itself, but you must use the backslash symbol followed by a single quotation mark symbol (\' escape sequence) to represent the single quotation mark symbol.

You can represent the new-line character by the \n new-line escape sequence.

You can represent the backslash character by the \\ backslash escape sequence.

The following are examples of character literals:

```
'a'
L'0'
```

RELATED REFERENCES

• "char and wchar_t Type Specifiers" on page 32

String Literals

A string literal contains a sequence of characters or escape sequences enclosed in double quotation mark symbols.



A string literal with the prefix L is a wide string literal. A string literal without the prefix L is an ordinary or narrow string literal.

The following are examples of string literals:

```
char titles[] = "Handel's \"Water Music\"";
char *mail_addr = "Last Name First Name MI Street Address \
   City Province Postal code ";
char *temp_string = "abc" "def" "ghi"; /* *temp_string = "abcdefghi\0" */
wchar t *wide string = L"longstring";
```

A null ('\0') character is appended to each string. For a wide string literal, the value '\0' of type **wchar_t** is appended. By convention, programs recognize the end of a string by finding the null character.

Multiple spaces contained within a string literal are retained.

To continue a string on the next line, use the line continuation sequence (\ symbol immediately followed by a new-line character). A carriage return must immediately follow the backslash. In the following example, the string literal second causes a compile-time error.

Another way to continue a string is to have two or more consecutive strings. Adjacent string literals will be concatenated to produce a single string. If a wide string literal and a narrow string literal are adjacent to each other, the resulting behavior is undefined. The following example demonstrates this:

```
"hello " "there" /* is equivalent to "hello there" */
"hello " L"there" /* the behavior is undefined */
"hello" "there" /* is equivalent to "hellothere" */
```

Characters in concatenated strings remain distinct. For example, the strings "\xab" and "3" are concatenated to form "\xab3". However, the characters \xab and 3 remain distinct and are not merged to form the hexadecimal character \xab3.

Following any concatenation, '\0' of type **char** is appended at the end of each string. C++ programs find the end of a string by scanning for this value. For a wide string literal, '\0' of type **wchar_t** is appended. For example:

The type of a narrow string literal is array of **const char** and the type of a wide string literal is array of **const wchar_t**. Both types have **static** storage duration.

The type of a narrow string literal is array of **char** and the type of a wide string literal is array of **wchar_t**.

Literals

Use the escape sequence \n to represent a new-line character as part of the string. Use the escape sequence \\ to represent a backslash character as part of the string. You can represent a single quotation mark symbol either by itself or with the escape sequence \'. You must use the escape sequence \" to represent a double quotation mark.

For example:

```
** This example illustrates escape sequences in string literals
#include <iostream>
using namespace std;
int main ()
      char *s = "Hi there! \n";
      cout << s;
      char *p = "The backslash character \\.";
      cout << p << endl;</pre>
      char *q = "The double quotation mark \".\n";
      cout << q;
}
```

This program produces the following output:

```
Hi there!
The backslash character \.
The double quotation mark ".
```

RELATED REFERENCES

- "char and wchar_t Type Specifiers" on page 32
- "volatile and const Qualifiers" on page 50
- "static Storage Class Specifier" on page 28

Boolean Literals

There are only two boolean literals: true and false. These literals have type bool and are not lvalues.

RELATED REFERENCES

- "Boolean Variables" on page 33
- "Lvalues and Rvalues" on page 75

Chapter 3. Declarations

A *declaration* establishes the names and characteristics of data objects and functions used in a program. A *definition* allocates storage for data objects or specifies the body for a function. When you define a type, no storage is allocated.

Declarations Overview

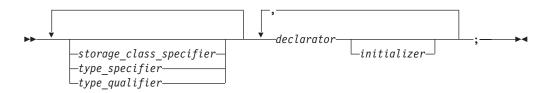
Declarations determine the following properties of data objects and their identifiers:

- Scope, which describes the visibility of an identifier in a block or source file.
- Linkage, which describes the association between two identical identifiers.
- Type, which describes the kind of data the object is to represent.

The lexical order of elements of a declaration for a data object is as follows:

- Storage duration and linkage specification.
- Type specification.
- Declarators, which introduce identifiers and make use of type qualifiers and storage qualifiers.
- Initializers, which initialize storage with initial values.

All data declarations have the form:



The following table shows examples of declarations and definitions. The identifiers declared in the first column do not allocate storage; they refer to a corresponding definition. In the case of a function, the corresponding definition is the code or body of the function. The identifiers declared in the second column allocate storage; they are both declarations and definitions.

Declarations	Declarations and Definitions
extern double pi;	double pi = 3.14159265;
<pre>float square(float x);</pre>	<pre>float square(float x) { return x*x; }</pre>
struct payroll;	<pre>struct payroll {</pre>

RELATED REFERENCES

- "Scope" on page 1
- "Program Linkage" on page 4
- "Storage Class Specifiers" on page 24
- "Type Specifiers" on page 30

Declarations

- "Chapter 4. Declarators" on page 55
- "Initializers" on page 56
- "Chapter 7. Functions" on page 123

Objects

An object is a region of storage that contains a value or group of values. Each value can be accessed using its identifier or a more complex expression that refers to the object. In addition, each object has a unique data type. Both the identifier and data type of an object are established in the object declaration.

The data type of an object determines the initial storage allocation for that object and the interpretation of the values during subsequent access. It is also used in any type-checking operations.

Both C and C++ have built-in, or fundamental data types and user-defined data types. Standard data types include signed and unsigned integers, floating-point numbers, and characters. User-defined types include enumerations, structures, unions, and classes.

An instance of a class type is commonly called a *class object*. The individual class members are also called objects. The set of all member objects comprises a class object.

RELATED REFERENCES

"Chapter 12. Classes" on page 199

Storage Class Specifiers

The storage class specifier used within the declaration determines whether:

- The object has internal, external, or no linkage
- The object is to be stored in memory or in a register, if available
- The object receives the default initial value 0 or an indeterminate default initial value
- The object can be referenced throughout a program or only within the function, block, or source file where the variable is defined
- The storage duration for the object is *static* (storage is maintained throughout program run time) or automatic (storage is maintained only during the execution of the block where the object is defined)

For a function, the storage class specifier determines the linkage of the function.

Declarations with the auto or register storage-class specifier result in automatic storage. Those with the **static** storage-class specifier result in static storage.

Most local declarations that do not include the extern storage-class specifier allocate storage; however, function declarations and type declarations do not allocate storage.

The only storage-class specifiers allowed in a namespace or global scope declaration are static and extern.

This section describes the following storage class specifiers:

- auto
- extern

- C++ mutable
- register
- static
- typedef

RELATED REFERENCES

- "auto Storage Class Specifier"
- "extern Storage Class Specifier"
- "mutable Storage Class Specifier" on page 27
- "register Storage Class Specifier" on page 27
- "static Storage Class Specifier" on page 28
- "typedef" on page 29
- "Program Linkage" on page 4

auto Storage Class Specifier

The auto storage class specifier lets you declare a variable with automatic storage. A variable x that has automatic storage is deleted when the block in which x was declared exits.

You can only apply the auto storage class specifier to names of variables declared in a block or to names of function parameters. However, these names by default have automatic storage. Therefore the storage class specifier auto is usually redundant in a data declaration.

Initialization

You can initialize any auto variable except parameters. If you do not initialize an automatic object, its value is indeterminate. If you provide an initial value, the expression representing the initial value can be any valid C or C++ expression. For structure and union members, the initial value must be a valid constant expression if an initializer list is used. The object is then set to that initial value each time the program block that contains the object's definition is entered.

Note that if you use the goto statement to jump into the middle of a block, automatic variables within that block are not initialized.

Storage

Objects with the auto storage class specifier have automatic storage duration. Each time a block is entered, storage for auto objects defined in that block is made available. When the block is exited, the objects are no longer available for use.

If an auto object is defined within a function that is recursively invoked, memory is allocated for the object at each invocation of the block.

RELATED REFERENCES

- "Block Statement" on page 145
- "goto Statement" on page 159
- "Function Declarations" on page 124

extern Storage Class Specifier

The extern storage class specifier lets you declare objects and functions that several source files can use. All object declarations that occur outside a function and that

Storage Class Specifiers

do not contain a storage class specifier declare identifiers with external linkage. All function definitions that do not specify a storage class define functions with external linkage.

Solution In C++, the **extern** storage class specifier can only be applied to names of objects or functions. Using the extern specifier with type declarations is illegal.

An extern variable, function definition, or declaration also makes the described variable or function usable by the succeeding part of the current source file. This declaration does not replace the definition. The declaration is used to describe the variable that is externally defined.

If a declaration for an identifier already exists at file scope, any extern declaration of the same identifier found within a block refers to that same object. If no other declaration for the identifier exists at file scope, the identifier has external linkage.

An **extern** declaration can appear outside a function or at the beginning of a block. If the declaration describes a function or appears outside a function and describes an object with external linkage, the keyword **extern** is optional.

If you do not specify a storage class specifier, the function has external linkage. It is an error to include a declaration for the same function with the storage class specifier static before the declaration with no storage class specifier because of the incompatible declarations. Including the extern storage class specifier on the original declaration is valid and the function has internal linkage.

In C++, an **extern** declaration cannot appear in class scope.

Initialization

You can initialize any object with the extern storage class specifier at namespace (or global scope in C). You can initialize an extern object with an initializer that must either:

- Appear as part of the definition and the initial value must be described by a constant expression. OR
- Reduce to the address of a previously declared object with static storage duration. You may modify this object with pointer arithmetic. (In other words, you may modify the object by adding or subtracting an integral constant expression.)

Note:

If you do not explicitly initialize an extern variable, its initial value is zero of the appropriate type. Initialization of an extern object is completed by the time the program starts running.

Storage

All extern objects have static storage duration. Memory is allocated for extern objects before the main function begins running. When the program finishes running, the storage is freed.

RELATED REFERENCES

- "External Linkage" on page 5
- "Internal Linkage" on page 4

- "static Storage Class Specifier" on page 28
- "Class Scope" on page 3
- "Chapter 10. Namespaces" on page 177

mutable Storage Class Specifier

The *mutable* storage class specifier is used only on a class data member to make it modifiable even though the member is part of an object declared as const. You cannot use the mutable specifier with names declared as **static** or **const**, or reference members.

```
class A
{
  public:
    A() : x(4), y(5) { };
    mutable int x;
    int y;
};
int main()
{
  const A var2;
  var2.x = 345;
  // var2.y = 2345;
}
```

In this example, the compiler would not allow the assignment var2.y = 2345 because var2 has been declared as const. The compiler will allow the assignment var2.x = 345 because A::x has been declared as **mutable**.

RELATED REFERENCES

- "static Storage Class Specifier" on page 28
- "volatile and const Qualifiers" on page 50
- "References" on page 68

register Storage Class Specifier

The **register** storage class specifier indicates to the compiler that a heavily used variable (such as a loop control variable) within a local scope data definition or a parameter declaration should be allocated a register to minimize access time.

It is equivalent to the **auto** storage class except that the compiler places the object, if possible, into a machine register for faster access. An object having the **register** storage class specifier must be defined within a block or declared as a parameter to a function.

Initialization

You can initialize any **register** object except parameters. If you do not initialize an automatic object, its value is indeterminate. If you provide an initial value, the expression representing the initial value can be any valid C or C++ expression. For structure and union members, the initial value must be a valid constant expression if an initializer list is used. The object is then set to that initial value each time the program block that contains the object's definition is entered.

Storage

Storage Class Specifiers

Objects with the **register** storage class specifier have automatic storage duration. Each time a block is entered, storage for register objects defined in that block are made available. When the block is exited, the objects are no longer available for

If a register object is defined within a function that is recursively invoked, the memory is allocated for the variable at each invocation of the block.

The register storage class specifier indicates that the object is heavily used and indicates to the compiler that the value of the object should reside in a machine register. Because of the limited size and number of registers available on most systems, few variables can actually be put in registers.

If the compiler does not allocate a machine register for a register object, the object is treated as having the storage class specifier auto.

Restrictions

You cannot use the register storage class specifier in namespace scope (or global scope in C) data declarations. In C programs, you cannot apply the address (&) operator to register variables. However, C++ lets you take the address of an object with the **register** storage class. For example:

```
register int i;
                 // valid in C++, but not in C
int* b = &i;
```

RELATED REFERENCES

- "Local Scope" on page 1
- "auto Storage Class Specifier" on page 25
- "References" on page 68

static Storage Class Specifier

The static storage class specifier lets you define objects with static storage duration and internal linkage, or to define functions with internal linkage.

The **static** storage class specifier can only be applied to the following names:

- Objects
- Functions
- Class members
- Anonymous unions

You cannot declare any of the following as **static**:

- Type declarations
- Function declarations within a block
- Function parameters

Objects with the static storage class specifier have static storage duration. The storage for a static variable is made available when the program begins running. When the program finishes running, the memory is freed.

For example, suppose a static variable x has been declared in function f(). When the program exits the scope of f(), x is not destroyed. The following example demonstrates this:

```
#include <stdio.h>
int f(void) {
 static int i = 0;
```

```
i++;
  return i;
}
int main(void) {
  int j;
  for (j = 0; j < 5; j++) {
    printf("Value of f(): %d\n", f());
  }
  return 0;
}</pre>
```

The following is the output of the above example:

```
Value of f(): 1
Value of f(): 2
Value of f(): 3
Value of f(): 4
Value of f(): 5
```

Because i is a static variable, it is not reinitialized to 0 on successive calls to f().

If you eplicitly declare an object, reference, function, or function template **static** in namespace or global scope, that entity will have internal linkage; you cannot use that entity in other translation units.

Initialization

- You initialize a **static** object with a constant expression, or an expression that reduces to the address of a previously declared **extern** or **static** object, possibly modified by a constant expression.
- You may initialize a **static** object with a non-constant expression.
- If you do not explicitly initialize a static (or external) variable, it will have a value of zero of the appropriate type.
- A static object of class type will use the default constructor if you do not initialize it.

Automatic and register variables that are not initialized will have undefined values.

RELATED REFERENCES

- "Internal Linkage" on page 4
- "extern Storage Class Specifier" on page 25
- "Objects" on page 24
- "Class Member Lists" on page 209
- "Anonymous Unions" on page 45

typedef

A **typedef** declaration lets you define your own identifiers that can be used in place of type specifiers such as **int**, **float**, and **double**. A **typedef** declaration does not reserve storage. The names you define using **typedef** are not new data types. They are synonyms for the data types or combinations of data types they represent.

Storage Class Specifiers

In C++, a **typedef** name must be different from any class type name declared within the same scope. If the **typedef** name is the same as a class type name, it can only be so if that typedef is a synonym of the class name. This condition is not the same as in C.

When an object is defined using a **typedef** identifier, the properties of the defined object are exactly the same as if the object were defined by explicitly listing the data type associated with the identifier.

Examples of typedef Declarations

The following statements declare LENGTH as a synonym for int and then use this typedef to declare length, width, and height as integer variables:

```
typedef int LENGTH;
LENGTH length, width, height;
```

The following declarations are equivalent to the above declaration:

```
int length, width, height;
```

Similarly, typedef can be used to define a class type (structure, union, or C++ class). For example:

```
typedef struct {
                int scruples;
                int drams;
                int grains;
               } WEIGHT;
```

The structure WEIGHT can then be used in the following declarations:

```
WEIGHT chicken, cow, horse, whale;
```

A C++ class defined in a **typedef** without being named is given a dummy name and the typedef name for linkage. Such a class cannot have constructors or destructors. For example:

```
typedef class {
               Trees();
              } Trees;
```

Here the function Trees() is an ordinary member function of a class whose type name is unspecified. In the above example, Trees is an alias for the unnamed class, not the class type name itself, so Trees() cannot be a constructor for that class.

RELATED REFERENCES

- · "Type Specifiers"
- "Arrays" on page 62
- "Structures" on page 36
- "Unions" on page 42
- "Chapter 12. Classes" on page 199
- "Constructors and Destructors Overview" on page 257

Type Specifiers

Solution In C++, types must be declared in declarations. They may not be declared in expressions.

Type specifiers indicate the type of the object or function being declared. The following are the available kinds of type specifiers:

- Simple type specifiers
- C++ Class specifiers
- Enumerated specifiers
- Elaborated type specifiers
- const and volatile qualifiers

RELATED REFERENCES

- "Simple Type Specifiers"
- "Declaring Class Types" on page 199
- "Enumerations" on page 46
- "Type Specifiers" on page 30
- "Type Specifiers" on page 30

Simple Type Specifiers

A simple type specifier either specifies a (previously declared) user-defined type or a fundamental type. A fundamental type is a type built-in to the language. The following describes how the fundamental types are categorized:

- Arithmetic types
 - Integral types
 - C++ bool
 - char
 - wchar_t
 - Signed integer types
 - signed char
 - short int
 - int
 - long int
 - Unsigned integer types
 - unsigned char
 - unsigned short int
 - · unsigned int
 - · unsigned long int
 - Floating-point types
 - float
 - double
 - long double
- void

RELATED REFERENCES

- "char and wchar_t Type Specifiers" on page 32
- "Boolean Variables" on page 33
- "Floating-Point Variables" on page 33
- "Integer Variables" on page 34
- "void Type" on page 35

char and wchar_t Type Specifiers

The **char** specifier has the following syntax:



The **char** specifier is an integral type.

A char has enough storage to represent a character from the basic character set. The amount of storage allocated for a **char** is implementation-dependent.

You initialize a variable of type char with a character literal (consisting of one character) or with an expression that evaluates to an integer.

Use **signed char** or **unsigned char** to declare numeric variables that occupy a single byte.

Solution For the purposes of distinguishing overloaded functions, a C++ char is a distinct type from signed char and unsigned char.

The wchar_t Type Specifier

The wchar_t type specifier has enough storage to represent a wide character literal. (A character literal that is prefixed with the letter L, for example L'x', is a wide character literal).

The wchar_t type specifier is an integral type.

Examples of the char Type Specifier

The following example defines the identifier end_of_string as a constant object of type **char** having the initial value \0 (the null character):

```
const char end of string = '\0';
```

The following example defines the unsigned char variable switches as having the initial value 3:

```
unsigned char switches = 3;
```

The following example defines string pointer as a pointer to a character: char *string_pointer;

The following example defines name as a pointer to a character. After initialization, name points to the first letter in the character string "Johnny":

```
char *name = "Johnny";
```

The following example defines a one-dimensional array of pointers to characters. The array has three elements. Initially they are a pointer to the string "Venus", a pointer to "Jupiter", and a pointer to "Saturn":

```
static char *planets[] = { "Venus", "Jupiter", "Saturn" };
```

RELATED REFERENCES

• "Character Literals" on page 20

• "Integer Variables" on page 34

Boolean Variables

Use the type specifier bool and the literals true and false to make boolean logic tests. A *boolean logic test* is used to express the results of a logical operation. For example:

```
bool f(int a, int b)
{
  return a==b;
}
```

If a and b have the same value, f() returns true. If not, f() returns false.

Variables of type bool can hold either one of two values: true or false. An rvalue of type bool can be promoted to an integral type. A bool rvalue of false is promoted to the value 0 and a bool rvalue of true is promoted to the value 1.

RELATED REFERENCES

- "Boolean Literals" on page 22
- "Integer Variables" on page 34

Floating-Point Variables

There are three types of floating-point variables:

- float
- double
- · long double

To declare a data object having a floating-point type, use the following **float** specifier:



The declarator for a simple floating-point declaration is an identifier. Initialize a simple floating-point variable with a float constant or with a variable or expression that evaluates to an integer or floating-point number. The storage class of a variable determines how you initialize the variable.

Examples of Floating-Point Data Types

The following example defines the identifier pi as an object of type **double**: double pi;

The following example defines the **float** variable real_number with the initial value 100.55:

```
static float real_number = 100.55f;
```

Note: If you do not add the **f** suffix to a floating-point literal, that number will be of type **double**. If you initalize an object of type **float** with an object of type **double**, the compiler will implictly convert the object of type **double** to an object of type **float**.

The following example defines the **float** variable float var with the initial value

```
float float_var = 1.43e-2f;
```

The following example declares the **long double** variable maximum: extern long double maximum;

The following example defines the array table with 20 elements of type **double**: double table[20];

RELATED REFERENCES

- "Floating-Point Literals" on page 18
- "Assignment Expressions" on page 107

Integer Variables

Integer variables fall into the following categories:

- integral types
 - ▶ C++ bool
 - char
 - wchar t
 - signed integer types
 - signed char
 - short int
 - int
 - long int
 - unsigned integer types
 - unsigned char
 - unsigned short int
 - unsigned int
 - unsigned long int

The default integer type for a bit field is **unsigned**.

The amount of storage allocated for integer data is implementation-dependent.

The **unsigned** prefix indicates that the object is a nonnegative integer. Each unsigned type provides the same size storage as its signed equivalent. For example, int reserves the same storage as unsigned int. Because a signed type reserves a sign bit, an unsigned type can hold a larger positive integer than the equivalent signed type.

The declarator for a simple integer definition or declaration is an identifier. You can initialize a simple integer definition with an integer constant or with an expression that evaluates to a value that can be assigned to an integer. The storage class of a variable determines how you can initialize the variable.

When the arguments in overloaded functions and overloaded operators are integer types, two integer types that both come from the same group are not treated as distinct types. For example, you cannot overload an int argument against a signed int argument.

Examples of Integer Data Types

The following example defines the **short int** variable flag:

```
short int flag;
```

The following example defines the **int** variable result: int result:

The following example defines the **unsigned long int** variable ss_number as having the initial value 438888834:

```
unsigned long ss_number = 438888834ul;
```

RELATED REFERENCES

- "Integer Literals" on page 17
- "Decimal Integer Literals" on page 17
- "Octal Decimal Literals" on page 18
- "Hexadecimal Decimal Literals" on page 18
- "Chapter 11. Overloading" on page 185

void Type

The **void** data type always represents an empty set of values. The only object that can be declared with the type specifier **void** is a pointer.

When a function does not return a value, you should use **void** as the type specifier in the function definition and declaration. An argument list for a function taking no arguments is **void**.

You cannot declare a variable of type **void**, but you can explicitly convert any expression to type **void**, but the resulting expression can only be used as one of the following:

- An expression statement
- The left operand of a comma expression
- The second or third operand in a conditional expression.

Example of void Type

In the following example, the function find max is declared as having type void.

Note: The use of the sizeof operator in the line find_max(numbers, (sizeof(numbers) / sizeof(numbers[0]))); is a standard method of determining the number of elements in an array.

```
/**
  ** Example of void type
  **/
#include <stdio.h>

/* declaration of function find_max */
extern void find_max(int x[], int j);

int main(void)
{
    static int numbers[] = { 99, 54, -102, 89 };
    find_max(numbers, (sizeof(numbers) / sizeof(numbers[0])));
    return(0);
}

void find_max(int x[], int j)
{ /* begin definition of function find_max */
    int i, temp = x[0];
```

```
for (i = 1; i < j; i++)
      if (x[i] > temp)
          temp = x[i];
   printf("max number = %d\n", temp);
} /* end definition of function find max */
```

RELATED REFERENCES

- "Pointers" on page 58
- "Comma Expression," on page 113
- "Conditional Expressions" on page 109
- "Function Declarations" on page 124

Structures

A structure contains an ordered group of data objects. Unlike the elements of an array, the data objects within a structure can have varied data types. Each data object in a structure is a member or field.

Use structures to group logically related objects. For example, to allocate storage for the components of one address, define the following variables:

```
int street no;
char *street name;
char *city;
char *prov;
char *postal code;
```

To allocate storage for more than one address, group the components of each address by defining a structure data type and as many variables as you need to have the structure data type.

In the following example, line int street_no; through to char *postal_code; declare the structure tag address:

```
struct address {
                 int street_no;
                 char *street name;
                 char *city;
                 char *prov;
                 char *postal code;
               };
struct address perm address;
struct address temp address;
struct address *p_perm_address = &perm_address;
```

The variables perm address and temp address are instances of the structure data type address. Both contain the members described in the declaration of address. The pointer p perm address points to a structure of address and is initialized to point to perm address.

Refer to a member of a structure by specifying the structure variable name with the dot operator (.) or a pointer with the arrow operator (->) and the member name. For example, both of the following:

```
perm address.prov = "Ontario";
p_perm_address -> prov = "Ontario";
```

assign a pointer to the string "Ontario" to the pointer prov that is in the structure perm address.

All references to structures must be fully qualified. In the example, you cannot reference the fourth field by prov alone. You must reference this field by perm_address.prov.

Structures with identical members but different names are not compatible and cannot be assigned to each other.

Structures are not intended to conserve storage. If you need direct control of byte mapping, use pointers.

You cannot declare a structure with members of incomplete types.

In C++ a structure is the same as a class except that its members and inheritance are public by default.

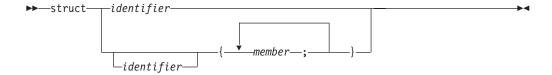
RELATED REFERENCES

- "Chapter 12. Classes" on page 199
- "Dot Operator ." on page 81
- "Arrow Operator ->" on page 82
- "Incomplete Types" on page 51

Declaring and Defining a Structure

A structure type definition describes the members that are part of the structure. It contains the struct keyword followed by an optional identifier (the structure tag) and a brace-enclosed list of members.

A structure definition has the form:



A structure declaration has the same form as a structure definition except the declaration does not have a brace-enclosed list of members.

The keyword struct followed by the identifier (tag) names the data type. If you do not provide a tag name to the data type, you must put all variable definitions that refer to it within the declaration of the data type.

The list of members provides the data type with a description of the values that can be stored in the structure.

A structure data member definition has the form of a variable declaration. However you may declare a bit-field as a member for a structure. A member that does not represent a bit field can be of any data type and can have the volatile or const qualifier.

Identifiers used as structure or member names can be redefined to represent different objects in the same scope without conflicting. You cannot use the name of a member more than once in a structure type, but you can use the same member name in another structure type that is defined within the same scope.

You cannot declare a structure type that contains itself as a member, but you can declare a structure type that contains a pointer to itself as a member.

RELATED REFERENCES

- "Declaring and Using Bit Fields in Structures" on page 39
- "volatile and const Qualifiers" on page 50

Defining a Structure Variable

A structure variable definition contains an optional storage class keyword, the struct keyword, a structure tag, a declarator, and an optional identifier. The structure tag indicates the data type of the structure variable.

The keyword **struct** is optional in C++.

You can declare structures having any storage class. Most compilers, however, treat structures declared with the register storage class specifier as automatic structures.

RELATED REFERENCES

- "auto Storage Class Specifier" on page 25
- "register Storage Class Specifier" on page 27

Initializing Structures

The initializer contains an = (equal sign) followed by a brace-enclosed comma-separated list of values. You do not have to initialize all members of a structure.

The following definition shows a completely initialized structure:

```
struct address {
                 int street no;
                 char *street_name;
                 char *city;
                 char *prov;
                 char *postal code;
static struct address perm_address =
               { 3, "Savona Dr.", "Dundas", "Ontario", "L4B 2A1"};
```

The values of perm address are:

```
Member
                             Value
perm address.street no
perm address.street name
                             address of string "Savona Dr."
perm address.city
                             address of string "Dundas"
perm address.prov
                             address of string "Ontario"
                             address of string "L4B 2A1"
perm address.postal code
```

The following definition shows a partially initialized structure:

```
struct address {
                 int street no;
                 char *street name;
                 char *city;
                 char *prov;
                 char *postal code;
struct address temp_address =
               { 44, "Knyvet Ave.", "Hamilton", "Ontario" };
```

The values of temp address are:

```
MemberValuetemp_address.street_no44temp_address.street_nameaddress of string "Knyvet Ave."temp_address.cityaddress of string "Hamilton"temp_address.provaddress of string "Ontario"temp address.postal codevalue depends on the storage class.
```

Note: The initial value of uninitialized structure members like temp_address.postal_code depends on the storage class associated with the member.

Declaring Structure Types and Variables

To define a structure type and a structure variable in one statement, put a declarator and an optional initializer after the type definition. To specify a storage class specifier for the variable, you must put the storage class specifier at the beginning of the statement.

Because this example does not name the structure data type, perm_address and temp_address are the only structure variables that will have this data type. Putting an identifier after **struct**, lets you make additional variable definitions of this data type later in the program.

The structure type (or tag) cannot have the **volatile** qualifier, but a member or a structure variable can be defined as having the **volatile** qualifier.

For example:

This example qualifies the structures file1 and file2, and the structure member subfile.code as **volatile**.

RELATED REFERENCES

- "Initializing Structures" on page 38
- "Storage Class Specifiers" on page 24
- "volatile and const Qualifiers" on page 50

Declaring and Using Bit Fields in Structures

A structure or a C++ class can contain *bit fields* that allow you to access individual bits. You can use bit fields for data that requires just a few bits of storage. A bit field declaration contains a type specifier followed by an optional declarator, a colon, a constant expression, and a semicolon. The constant expression specifies how many bits the field reserves.

Bit fields with a length of 0 must be unnamed. Unnamed bit fields cannot be referenced or initialized. A zero-width bit field causes the next field to be aligned on the next container boundary, where the container is the same size as the underlying type as the bit field. The padding to the next container boundary only takes place if the zero-width bit field has the same underlying type as the preceding bit-field member. If the types are different, the zero-width bit field has no effect.

The maximum bit-field length is implementation dependent.

For portability, do not use bit fields greater than 32 bits in size.

The following restrictions apply to bit fields. You cannot:

- Define an array of bit fields
- Take the address of a bit field
- Have a pointer to a bit field
- Have a reference to a bit field

In C, you can declare a bit field as type int, signed int, or unsigned int. Bit fields of the type int are equivalent to those of type unsigned int.

The default integer type for a bit field is **unsigned**.

A bit field cannot have the **const** or **volatile** qualifier.

The following structure has three bit-field members kingdom, phylum, and genus, occupying 12, 6, and 2 bits respectively:

```
struct taxonomy {
     int kingdom : 12;
     int phylum : 6;
     int genus : 2;
     };
```

Unlike ANSI/ISO C, C++ bit fields can be any integral type or enumeration type. When you assign a value that is out of range to a bit field, the low-order bit pattern is preserved and the appropriate bits are assigned.

If a series of bit fields does not add up to the size of an int, padding can take place. The amount of padding is determined by the alignment characteristics of the members of the structure.

The following example demonstrates padding. Suppose that an **int** occupies 4 bytes. The example declares the identifier kitchen to be of type struct on off:

```
struct on off {
                 unsigned light: 1;
                 unsigned toaster : 1;
                                       /* 4 bytes */
                 int count;
                 unsigned ac: 4;
                 unsigned: 4;
                 unsigned clock: 1;
                 unsigned: 0;
                 unsigned flag: 1;
                } kitchen;
```

The structure kitchen contains eight members totalling 16 bytes. The following table describes the storage that each member occupies:

Member Name	Storage Occupied		
light	1 bit		
toaster	1 bit		
(padding — 30 bits)	To the next int boundary		
count	The size of an int (4 bytes)		
ac	4 bits		
(unnamed field)	1 bit		
clock	1 bit		
(padding — 23 bits)	To the next int boundary (unnamed field)		
flag	1 bit		
(padding — 31 bits)	To the next int boundary		

All references to structure fields must be fully qualified. For instance, you cannot reference the second field by toaster. You must reference this field by kitchen.toaster.

The following expression sets the light field to 1:

```
kitchen.light = 1;
```

When you assign to a bit field a value that is out of its range, the bit pattern is preserved and the appropriate bits are assigned. The following expression sets the toaster field of the kitchen structure to θ because only the least significant bit is assigned to the toaster field:

```
kitchen.toaster = 2;
```

RELATED REFERENCES

- "Chapter 12. Classes" on page 199
- "Arrays" on page 62
- "Address &" on page 90
- "Pointers" on page 58
- "References" on page 68

Example Program Using Structures

The following program finds the sum of the integer numbers in a linked list:

```
name1.next num = &name2;
name2.next num = &name3;
name3.next_num = NULL;
while (recd pointer != NULL)
   sum += recd pointer->number;
   recd_pointer = recd_pointer->next_num;
printf("Sum = %d\n", sum);
return(0);
```

The structure type record contains two members: the integer number and next_num, which is a pointer to a structure variable of type record.

The record type variables name1, name2, and name3 are assigned the following values:

Member Name	Value
name1.number	144
name1.next_num	The address of name2
name2.number name2.next_num	203 The address of name3
name3.number name3.next num	488 NULL (Indicating the end of the linked list.)

The variable recd pointer is a pointer to a structure of type record. It is initialized to the address of name1 (the beginning of the linked list).

The while loop causes the linked list to be scanned until recd_pointer equals NULL. The statement:

```
recd pointer = recd pointer->next num;
```

advances the pointer to the next object in the list.

RELATED REFERENCES

- "Chapter 4. Declarators" on page 55
- "Initializers" on page 56
- "Incomplete Types" on page 51
- "Dot Operator ." on page 81
- "Arrow Operator ->" on page 82

Unions

A union is an object similar to a structure except that all of its members start at the same location in memory. A union can contain only one of its members at a time. The members of a union can be of any data type.

The storage allocated for a union is the storage required for the largest member of the union (plus any padding that is required so that the union will end at a natural boundary of its strictest member).

C++ In C++, a union is a limited form of the class type. It can contain access specifiers (public, protected, private), member data, and member functions, including constructors and destructors. It cannot contain virtual member functions or static data members. Default access of members in a union is public. A union cannot be used as a base class and cannot be derived from a base class.

C++ A C++ union member cannot be a class object that has a constructor, destructor, or overloaded copy assignment operator. Also, a union cannot have members of reference type. In C++, a member of a union cannot be declared with the keyword **static**.

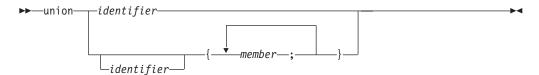
RELATED REFERENCES

- "Member Functions" on page 211
- "Constructors and Destructors Overview" on page 257
- "Virtual Functions" on page 249
- "Overloading Assignments" on page 190
- "static Storage Class Specifier" on page 28

Declaring a Union

A union type definition contains the union keyword followed by an identifier (optional) and a brace-enclosed list of members.

A union definition has the following form:



A union declaration has the same form as a union definition except that the declaration has no brace-enclosed list of members.

The *identifier* is a tag given to the union specified by the member list. If you specify a tag, any subsequent declaration of the union (in the same scope) can be made by declaring the tag and omitting the member list. If you do not specify a tag, you must put all variable definitions that refer to that union within the statement that defines the data type.

The list of members provides the data type with a description of the objects that can be stored in the union.

A union member definition has same form as a variable declaration.

You can reference one of the possible members of a union the same way as referencing a member of a structure.

```
For example:
union {
     char birthday[9];
     int age;
```

```
float weight;
      } people;
people.birthday[0] = '\n';
```

assigns '\n' to the first element in the character array birthday, a member of the union people.

A union can represent only one of its members at a time. In the example, the union people contains either age, birthday, or weight but never more than one of these. The printf statement in the following example does not give the correct result because people.age replaces the value assigned to people.birthday in the first line:

```
#include <stdio.h>
#include <string.h>
union {
 char birthday[9];
  int age;
 float weight;
} people;
int main(void) {
  strcpy(people.birthday, "03/06/56");
 printf("%s\n", people.birthday);
 people.age = 38;
 printf("%s\n", people.birthday);
```

The output of the above example will be similar to the following: 03/06/56

Defining a Union Variable

A union variable definition has the following form::

```
-union—union_data_type_name—identifier-
dashstorage class specifier
---initialiaztion value-
```

You must declare the union data type before you can define a union having that type.

You can define a union data type and a union of that type in the same statement by placing the variable declarator after the data type definition.

You can only initialize the first member of a union.

The following example shows how you would initialize the first union member birthday of the union variable people:

```
union {
      char birthday[9];
      int age;
      float weight;
      } people = {"23/07/57"};
```

To define union type and a union variable in one statement, put a declarator after the type definition. The storage class specifier for the variable must go at the beginning of the statement.

RELATED REFERENCES

- "Storage Class Specifiers" on page 24
- "Unions" on page 42
- "Type Specifiers" on page 30
- "volatile and const Qualifiers" on page 50

Anonymous Unions

An *anonymous union* is a union without a class name. It cannot be followed by a declarator. An anonymous union is not a type; it defines an unnamed object and it cannot have member functions.

The member names of an anonymous union must be distinct from other names within the scope in which the union is declared. You can use member names directly in the union scope without any additional member access syntax.

For example, in the following code fragment, you can access the data members i and cptr directly because they are in the scope containing the anonymous union. Because i and cptr are union members and have the same address, you should only use one of them at a time. The assignment to the member cptr will change the value of the member i.

An anonymous union cannot have protected or private members. A global or namespace anonymous union must be declared with the keyword **static**.

RELATED REFERENCES

- "static Storage Class Specifier" on page 28
- "Member Functions" on page 211

Examples of Unions

The following example defines a union data type (not named) and a union variable (named length). The member of length can be a **long int**, a **float**, or a **double**.

```
union {
     float meters;
     double centimeters;
     long inches;
} length;
```

The following example defines the union type data as containing one member. The member can be named charctr, whole, or real. The second statement defines two data type variables: input and output.

```
union data {
             char charctr;
             int whole;
             float real;
union data input, output;
```

The following statement assigns a character to input:

```
input.charctr = 'h';
```

The following statement assigns a floating-point number to member output: output.real = 9.2;

The following example defines an array of structures named records. Each element of records contains three members: the integer id_num, the integer type_of_input, and the union variable input. input has the union data type defined in the previous example.

```
struct {
         int id num;
         int type of input;
         union data input;
       } records[10];
```

The following statement assigns a character to the structure member input of the first element of records:

```
records[0].input.charctr = 'g';
```

RELATED REFERENCES

- "Initializers" on page 56
- "Structures" on page 36
- "Dot Operator ." on page 81
- "Arrow Operator ->" on page 82

Enumerations

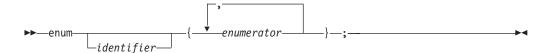
An enumeration data type represents a set of values that you declare. You can define an enumeration data type and all variables that have that enumeration type in one statement, or you can declare an enumeration type separately from the definition of variables of that type. The identifier associated with the data type (not an object) is called an *enumeration tag*.

RELATED REFERENCES

• "Type Specifiers" on page 30

Declaring an Enumeration Data Type

An enumeration type declaration contains the enum keyword followed by an optional identifier (the enumeration tag) and a brace-enclosed list of enumerators. Commas separate each enumerator in the enumerator list.



The keyword **enum**, followed by the identifier, names the data type (like the tag on a **struct** data type). The list of enumerators provides the data type with a set of values.

In C, each enumerator represents an integer value. In C++, each enumerator represents a value that can be converted to an integral value.

An enumerator has the form:



To conserve space, enumerations may be stored in spaces smaller than that of an **int**.

Enumeration Constants

When you define an enumeration data type, you specify a set of identifiers that the data type represents. Each identifier in this set is an *enumeration constant*.

The value of the constant is determined in the following way:

- 1. An equal sign (=) and a constant expression after the enumeration constant gives an explicit value to the constant. The identifier represents the value of the constant expression.
- 2. If no explicit value is assigned, the leftmost constant in the list receives the value zero (0).
- 3. Identifiers with no explicitly assigned values receive the integer value that is one greater than the value represented by the previous identifier.
- In C, enumeration constants have type int.
- In C++, each enumeration constant has a value that can be promoted to a signed or unsigned integer value and a distinct type that does not have to be integral. Use an enumeration constant anywhere an integer constant is allowed, or for C++, anywhere a value of the enumeration type is allowed.

Each enumeration constant must be unique within the scope in which the enumeration is defined. In the following example, second declarations of average and poor cause compiler errors:

```
func()
{
   enum score { poor, average, good };
   enum rating { below, average, above };
   int poor;
}
```

The following data type declarations list oats, wheat, barley, corn, and rice as enumeration constants. The number under each constant shows the integer value.

```
enum grain { oats, wheat=10, barley, corn=20, rice };
  /*
       0
                 10
                         11
```

It is possible to associate the same integer with two different enumeration constants. For example, the following definition is valid. The identifiers suspend and hold have the same integer value.

```
enum status { run, clear=5, suspend, resume, hold=6 };
```

RELATED REFERENCES

- "Integer Variables" on page 34
- "Integral and Floating-Point Promotions" on page 115

Defining Enumeration Variables

An enumeration variable definition has the following form:

```
-enum—enumeration data type name—identifier—
└storage class specifier
---enumeration constant-
```

You must declare the enumeration data type before you can define a variable having that type.

The initializer for an enumeration variable contains the = symbol followed by an expression enumeration_constant. In C++, the initializer must have the same type as the associated enumeration type.

The first line of the following example declares the enumeration grain. The second line defines the variable g food and gives g food the initial value of barley (2). enum grain { oats, wheat, barley, corn, rice };

The type specifier enum grain indicates that the value of g_food is a member of the enumerated data type grain.

C++ In C++, the **enum** keyword is optional when declaring a variable with enumeration type. However, it is required when declaring the enumeration itself. For example, both statements declare a variable of enumeration type:

```
enum grain g food = barley;
     grain cob_food = corn;
```

enum grain g_food = barley;

RELATED REFERENCES

• "Storage Class Specifiers" on page 24

Defining an Enumeration Type and Enumeration Objects

You can define a type and a variable in one statement by using a declarator and an optional initializer after the type definition. To specify a storage class specifier for the variable, you must put the storage class specifier at the beginning of the declaration. For example:

```
register enum score { poor=1, average, good } rating = good;
```

C++ also lets you put the storage class immediately before the declarator list. For example:

```
enum score { poor=1, average, good } register rating = good;
```

Either of these examples is equivalent to the following two declarations:

```
enum score { poor=1, average, good };
register enum score rating = good;
```

Both examples define the enumeration data type score and the variable rating. rating has the storage class specifier **register**, the data type enum score, and the initial value good.

Combining a data type definition with the definitions of all variables having that data type lets you leave the data type unnamed. For example:

defines the variable weekday, which can be assigned any of the specified enumeration constants.

Example Program Using Enumerations

The following program receives an integer as input. The output is a sentence that gives the French name for the weekday that is associated with the integer. If the integer is not associated with a weekday, the program prints "C'est le mauvais jour."

```
** Example program using enumerations
**/
#include <stdio.h>
enum days {
           Monday=1, Tuesday, Wednesday,
            Thursday, Friday, Saturday, Sunday
          } weekday;
void french(enum days);
int main(void)
  int num;
  printf("Enter an integer for the day of the week. "
          "Mon=1,...,Sun=7n");
   scanf("%d", &num);
  weekday=num;
   french (weekday);
   return(0);
void french(enum days weekday)
  switch (weekday)
     case Monday:
        printf("Le jour de la semaine est lundi.\n");
        break;
     case Tuesday:
        printf("Le jour de la semaine est mardi.\n");
        break;
     case Wednesday:
         printf("Le jour de la semaine est mercredi.\n");
```

```
break:
case Thursday:
   printf("Le jour de la semaine est jeudi.\n");
case Friday:
   printf("Le jour de la semaine est vendredi.\n");
   break;
case Saturday:
  printf("Le jour de la semaine est samedi.\n");
   break;
case Sunday:
   printf("Le jour de la semaine est dimanche.\n");
   break;
default:
   printf("C'est le mauvais jour.\n");
```

volatile and const Qualifiers

The volatile qualifier maintains consistency of memory access to data objects.

The volatile qualifier is useful for data objects having values that may be changed in ways unknown to your program (such as the system clock). Portions of an expression that reference **volatile** objects are not to be changed or removed.

The **const** qualifier explicitly declares a data object as a data item that cannot be changed. Its value is set at initialization. You cannot use const data objects in expressions requiring a modifiable lvalue. For example, a const data object cannot appear on the left-hand side of an assignment statement.

These type qualifiers are only meaningful in expressions that are lvalues.

For a volatile or const pointer, you must put the keyword between the * and the identifier. For example:

```
int * volatile x;
                        /* x is a volatile pointer to an int */
                      /* y is a const pointer to the int variable z */
int * const y = &z;
```

For a pointer to a volatile or const data object, the type specifier, qualifier, and storage class specifier can be in any order. For example:

```
volatile int *x;
                        /* x is a pointer to a volatile int */
int volatile *x;
                       /* x is a pointer to a volatile int */
const int *y;
                        /* y is a pointer to a const int */
int const *y;
                        /* y is a pointer to a const int */
```

In the following example, the pointer to y is a constant. You can change the value that y points to, but you cannot change the value of y:

```
int * const y
```

In the following example, the value that y points to is a constant integer and cannot be changed. However, you can change the value of y:

```
const int * y
```

For other types of volatile and const variables, the position of the keyword within the definition (or declaration) is less important. For example:

```
volatile struct omega {
                         int limit;
                         char code;
                      } group;
provides the same storage as:
struct omega {
                int limit;
                char code;
             } volatile group;
```

In both examples, only the structure variable group receives the volatile qualifier. Similarly, if you specified the const keyword instead of volatile, only the structure variable group receives the const qualifier. The const and volatile qualifiers when applied to a structure, union, or class also apply to the members of the structure, union, or class.

Although enumeration, class, structure, and union variables can receive the volatile or const qualifier, enumeration, class, structure, and union tags do not carry the volatile or const qualifier. For example, the blue structure does not carry the **volatile** qualifier:

```
volatile struct whale {
                           int weight;
                           char name[8];
                      } beluga;
struct whale blue;
```

The keywords volatile and const cannot separate the keywords enum, class, struct, and **union** from their tags.

You can declare or define a volatile or const function only if it is a C++ member function. You can define or declare any function to return a pointer to a volatile or **const** function.

You can put more than one qualifier on a declaration but you cannot specify the same qualifier more than once on a declaration.

An item can be both **const** and **volatile**. In this case the item cannot be legitimately modified by its own program but can be modified by some asynchronous process.

RELATED REFERENCES

- "Lvalues and Rvalues" on page 75
- "Pointers" on page 58
- "Chapter 12. Classes" on page 199
- "Structures" on page 36
- "Unions" on page 42

Incomplete Types

The following are incomplete types:

- Type void
- Array of unknown size
- Structure, union, or enumerations that have no member lists
- Pointers to class types that are declared but not defined

The following example is an incomplete type:

```
void *incomplete ptr;
struct dimension linear; /* no previous definition of dimension */
```

void is an incomplete type that cannot be completed. Incomplete structure or union and enumeration tags must be completed before being used to declare an object, although you can define a pointer to an incomplete structure or union.

An incomplete class declaration is a class declaration that does not define any class members. You cannot declare any objects of the class type or refer to the members of a class until the declaration is complete. However, an incomplete declaration allows you to make specific references to a class prior to its definition as long as the size of the class is not required.

For example, you can define a pointer to the structure first in the definition of the structure second. Structure first is declared in an incomplete class declaration prior to the definition of second, and the definition of oneptr in structure second does not require the size of first:

```
// incomplete declaration of struct first:
struct first;
// complete declaration of struct second:
struct second
 // pointer to struct first refers to
 // struct first prior to its complete
 // declaration:
 first* oneptr;
 // error, you cannot declare an object of
  // an incompletely declared class type:
 first one;
 int x, y;
};
// complete declaration of struct first:
struct first
  // define an object of class type second:
 second two;
 int z;
};
If you declare a class with an empty member list, it is a complete class declaration.
For example:
// incomplete class declaration:
class X;
// empty member list:
class Z {};
class Y
public:
  // error, cannot create an object of an
```

// incomplete class type:

X yobj;

```
// valid:
 Z zobj;
};
```

RELATED REFERENCES

- "void Type" on page 35
- "Arrays" on page 62
- "Structures" on page 36
- "Unions" on page 42
- "Chapter 12. Classes" on page 199
- "Incomplete Class Declarations" on page 204

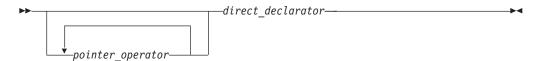
Chapter 4. Declarators

A *declarator* designates a data object or function. Declarators appear in most data definitions and declarations and in some type definitions.

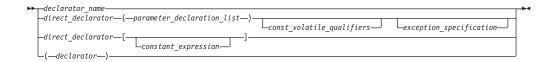
In a declarator, you can specify the type of an object to be an array, a pointer, or a reference. You can also perform initialization in a declarator.

A declarator has the form:

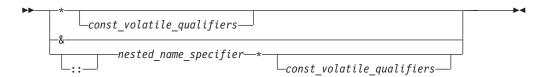
declarator



direct_declarator



pointer_operator



declarator_name



The following variables or delimiters are only available in C++:

- exception_specification
- nested_name_specifier
- :: (scope resolution operator)

The *const_volatile_qualifiers* variable represents one or a combination of **const** and **volatile**.

Declarators

In C, you cannot declare or define a volatile or const function. C++ class member functions can be qualified with **const** or **volatile**.

The following table illustrate some examples of declarators:

Example	Description
int owner	owner is an int data object.
int *node	node is a pointer to an int data object.
int names[126]	names is an array of 126 int elements.
<pre>int *action()</pre>	action is a function returning a pointer to an
	int.
volatile int min	min is an int that has the volatile qualifier.
int * volatile volume	volume is a volatile pointer to an int .
volatile int * next	next is a pointer to a volatile int.
<pre>volatile int * sequence[5]</pre>	sequence is an array of five pointers to
	volatile int objects.
extern const volatile int op_system_clock	op_system_clock is a constant and volatile integer with static storage duration and
	external linkage.

RELATED REFERENCES

- "volatile and const Qualifiers" on page 50
- "Enumerations" on page 46
- "Pointers" on page 58
- "Arrays" on page 62
- "Structures" on page 36
- "Unions" on page 42
- "Exception Specifications" on page 329
- "Scope of Class Names" on page 203

Initializers

An initializer is an optional part of a data declaration that specifies an initial value of a data object.

The initialization properties of each data type are described in the section for that data type.

The initializer consists of the = symbol followed by an initial expression or a braced list of initial expressions separated by commas. The number of initializers must not be more than the number of elements to be initialized. The initial expression evaluates to the first value of the data object.

To assign a value to an arithmetic or pointer type, use the simple initializer: = expression. For example, the following data definition uses the initializer = 3 to set the initial value of group to 3:

int group = 3;

For unions, structures, and aggregate classes (classes with no constructors, base classes, virtual functions, or private or protected members), the set of initial expressions must be enclosed in { } (braces) unless the initializer is a string literal. If the initializer of a character string is a string literal, the { } are optional. Individual expressions must be separated by commas, and groups of expressions can be enclosed in braces and separated by commas.

In an array, structure, or union initialized using a brace-enclosed initializer list, any members or subscripts that are not initialized are implicitly initialized to zero of the appropriate type.

An initializer of the form (expression) can be used to initialize fundamental types in C++. For example, the following two initializations are identical:

```
int group = 3;
int group(3);
```

You can also use the (expression) form to initialize C++ classes. However, the form with parentheses and the form with the assignment operator are not identical. The form with parenthesis calls the copy constructor of the class. The form with the assignment operator calls the copy assignment operator of the class.

You can initialize variables at namespace scope with nonconstant expressions. In C, you cannot do the same at global scope.

If your code jumps over declarations that contain initializations, the compiler generates an error. For example, the following code is not valid:

```
goto skiplabel; // error - jumped over declaration
int i = 3;
                 // and initialization of i
skiplabel: i = 4;
```

You can initialize classes in external, static, and automatic definitions. The initializer contains an = (equal sign) followed by a brace-enclosed, comma-separated list of values. You do not need to initialize all members of a class.

In the following example, only the first eight elements of the array grid are explicitly initialized. The remaining four elements that are not explicitly initialized are initialized as if they were explicitly initialized to zero.

```
static short grid[3][4] = \{0, 0, 0, 1, 0, 0, 1, 1\};
```

The initial values of grid are:

Element	Value	Element	Value	
grid[0] [0]	0	grid[1] [2]	1	
grid[0] [1]	0	grid[1] [3]	1	
grid[0] [2]	0	grid[2] [0]	0	
grid[0] [3]	1	grid[2] [1]	0	
grid[1] [0]	0	grid[2] [2]	0	
grid[1] [1]	0	grid[2] [3]	0	

RELATED REFERENCES

- "Copy Assignment Operators" on page 279
- "Assignment Expressions" on page 107
- "Arrays" on page 62
- "char and wchar_t Type Specifiers" on page 32
- "Simple Type Specifiers" on page 31

Initializers

- "Enumerations" on page 46
- "Floating-Point Variables" on page 33
- "Integer Variables" on page 34
- "Pointers"
- "Structures" on page 36
- "Unions" on page 42
- "Storage Class Specifiers" on page 24

Pointers

A pointer type variable holds the address of a data object or a function. A pointer can refer to an object of any one data type except to a bit field or a reference. Some common uses for pointers are:

- To access dynamic data structures such as linked lists, trees, and queues.
- To access elements of an array or members of a structure or C++ class.
- To access an array of characters as a string.
- To pass the address of a variable to a function. (In C++, you can also use a reference to do this.) By referencing a variable through its address, a function can change the contents of that variable.

RELATED REFERENCES

- "Calling Functions and Passing Arguments" on page 134
- "References" on page 68

Declaring Pointers

The following example declares proat as a pointer to an object having type long: extern long *pcoat;

If the keyword **volatile** appears before the *, the declarator describes a pointer to a volatile object. If the keyword volatile comes between the * and the identifier, the declarator describes a volatile pointer. The keyword const operates in the same manner as the **volatile** keyword described. In the following example, pvolt is a constant pointer to an object having type **short**:

```
short * const pvolt;
```

The following example declares pnut as a pointer to an **int** object having the **volatile** qualifier:

```
extern int volatile *pnut;
```

The following example defines psoup as a **volatile** pointer to an object having type float:

```
float * volatile psoup;
```

The following example defines pfowl as a pointer to an enumeration object of type bird:

```
enum bird *pfowl;
```

The next example declares pyish as a pointer to a function that takes no parameters and returns a char object:

```
char (*pvish)(void);
```

RELATED REFERENCES

• "volatile and const Qualifiers" on page 50

Assigning Pointers

When you use pointers in an assignment operation, you must ensure that the types of the pointers in the operation are compatible.

The following example shows compatible declarations for the assignment operation:

```
float subtotal;
float * sub_ptr;
// ...
sub_ptr = &subtotal;
printf("The subtotal is %f\n", *sub_ptr);
```

The next example shows incompatible declarations for the assignment operation:

```
double league;
int * minor;
// ...
minor = &league; /* error */
```

RELATED REFERENCES

- "Pointers" on page 58
- "Assignment Expressions" on page 107

Initializing Pointers

The initializer is an = (equal sign) followed by the expression that represents the address that the pointer is to contain. The following example defines the variables time and speed as having type **double** and amount as having type pointer to a **double**. The pointer amount is initialized to point to total:

```
double total, speed, *amount = &total;
```

The compiler converts an unsubscripted array name to a pointer to the first element in the array. You can assign the address of the first element of an array to a pointer by specifying the name of the array. The following two sets of definitions are equivalent. Both define the pointer student and initialize student to the address of the first element in section:

```
int section[80];
int *student = section;

is equivalent to:
int section[80];
int *student = &section[0];
```

You can assign the address of the first character in a string constant to a pointer by specifying the string constant in the initializer.

The following example defines the pointer variable string and the string constant "abcd". The pointer string is initialized to point to the character a in the string "abcd".

```
char *string = "abcd";
```

The following example defines weekdays as an array of pointers to string constants. Each element points to a different string. The pointer weekdays[2], for example, points to the string "Tuesday".

Initializers

```
static char *weekdays[] =
              "Sunday", "Monday", "Tuesday", "Wednesday",
              "Thursday", "Friday", "Saturday'
           };
```

A pointer can also be initialized to NULL using any integer constant expression that evaluates to 0, for example char * a=0;. Such a pointer is a NULL pointer. It does not point to any object.

RELATED REFERENCES

- "Initializers" on page 56
- "Arrays" on page 62

Restrictions on Pointers

You cannot use pointers to reference bit fields or objects having the register storage class specifier.

RELATED REFERENCES

"register Storage Class Specifier" on page 27

Using Pointers

Two operators are commonly used in working with pointers, the address (&) operator and the indirection (*) operator. You can use the & operator to refer to the address of an object. For example, the assignment in the following function assigns the address of x to the variable p to int. The variable p to int has been defined as a pointer:

```
void f(int x, int *p to int) {
 p_to_int = &x;
```

The * (indirection) operator lets you access the value of the object a pointer refers to. The assignment in the following example assigns to y the value of the object that p_to_float points to:

```
void g(float y, float *p to float) {
 y = *p_to_float;
```

The assignment in the following example assigns the value of z to the variable that *p to z references:

```
void h(char z, char *p_to_char) {
  *p_to_char = z;
```

RELATED REFERENCES

- "Address &" on page 90
- "Indirection *" on page 90

Pointer Arithmetic

You can perform a limited number of arithmetic operations on pointers. These operations are:

- · Increment and decrement
- · Addition and subtraction
- Comparison
- Assignment

The increment (++) operator increases the value of a pointer by the size of the data object the pointer refers to. For example, if the pointer refers to the second element in an array, the ++ makes the pointer refer to the third element in the array.

The decrement (--) operator decreases the value of a pointer by the size of the data object the pointer refers to. For example, if the pointer refers to the second element in an array, the -- makes the pointer refer to the first element in the array.

You can add an integer to a pointer but you cannot add a pointer to a pointer.

If the pointer p points to the first element in an array, the following expression causes the pointer to point to the third element in the same array:

```
p = p + 2;
```

If you have two pointers that point to the same array, you can subtract one pointer from the other. This operation yields the number of elements in the array that separate the two addresses that the pointers refer to.

You can compare two pointers with the following operators: ==, !=, <, >, <=, and >=.

Pointer comparisons are defined only when the pointers point to elements of the same array. Pointer comparisons using the == and != operators can be performed even when the pointers point to elements of different arrays.

You can assign to a pointer the address of a data object, the value of another compatible pointer or the NULL pointer.

RELATED REFERENCES

- "Pointers" on page 58
- "Increment ++" on page 87
- "Arrays" on page 62
- "Decrement --" on page 88
- "Chapter 5. Expressions and Operators" on page 71

Example Program Using Pointers

The following program contains pointer arrays:

```
** Program to search for the first occurrence of a specified **
  character string in an array of character strings.
*************************
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#define SIZE 20
int main(void)
  static char *names[] = { "Jim", "Amy", "Mark", "Sue", NULL };
  char * find name(char **, char *);
  char new name[SIZE], *name pointer;
  printf("Enter name to be searched.\n");
  scanf("%s", new name);
  name pointer = find name(names, new name);
  printf("name %s%sfound\n", new name,
```

```
} /* End of main */
      Function find name. This function searches an array of
      names to see if a given name already exists in the array.
      It returns a pointer to the name or NULL if the name is
**
** char **arry is a pointer to arrays of pointers (existing names) **
** char *strng is a pointer to character array entered (new name) **
char * find name(char **arry, char *strng)
  for (; *arry != NULL; arry++)
                                  /* for each name
                                                          */
     if (strcmp(*arry, strng) == 0) /* if strings match
       return(*arry);
                                   /* found it!
  return(*arry);
                                   /* return the pointer
} /* End of find name */
```

(name pointer == NULL) ? " not " : " ");

Interaction with this program could produce the following sessions:

Output Enter name to be searched.

Input Mark

Output name Mark found

or:

Enter name to be searched. Output

Input Deborah

name Deborah not found Output

RELATED REFERENCES

- "Chapter 4. Declarators" on page 55
- "volatile and const Qualifiers" on page 50
- "Initializers" on page 56
- "Address &" on page 90
- "Indirection *" on page 90

Arrays

An array is an ordered group of data objects. Each object is called an element. All elements within an array have the same data type.

Array elements cannot be of function data type or, in C++, of reference data type. You can, however, declare an array of pointers to functions.

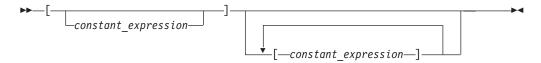
RELATED REFERENCES

- "Declaring Arrays" on page 63
- "Initializing Arrays" on page 64

Declaring Arrays

The array declarator contains an identifier followed by an optional *subscript declarator*. An identifier preceded by an * (asterisk) is an array of pointers.

A subscript declarator has the form:



The subscript declarator describes the number of dimensions in the array and the number of elements in each dimension. Each bracketed expression, or subscript, describes a different dimension and must be a constant expression.

The following example defines a one-dimensional array that contains four elements having type **char**:

```
char list[4];
```

The first subscript of each dimension is 0. The array list contains the elements:

- list[0]
- list[1]
- list[2]
- list[3]

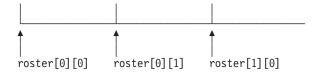
The following example defines a two-dimensional array that contains six elements of type **int**:

```
int roster[3][2];
```

Multidimensional arrays are stored in row-major order. When elements are referred to in order of increasing storage location, the last subscript varies the fastest. For example, the elements of array roster are stored in the order:

```
roster[0][0]
roster[0][1]
roster[1][0]
roster[1][1]
roster[2][0]
roster[2][1]
```

In storage, the elements of roster would be stored as:



You can leave the first (and only the first) set of subscript brackets empty in

- · Array definitions that contain initializations
- extern declarations
- · Parameter declarations

In array definitions that leave the first set of subscript brackets empty, the initializer determines the number of elements in the first dimension. In a one-dimensional array, the number of initialized elements becomes the total

Initializers

number of elements. In a multidimensional array, the initializer is compared to the subscript declarator to determine the number of elements in the first dimension.

An unsubscripted array name (for example, region instead of region[4]) represents a pointer whose value is the address of the first element of the array, provided the array has previously been declared.

Whenever an array is used in a context (such as a parameter) where it cannot be used as an array, the identifier is treated as a pointer. The two exceptions are when an array is used as an operand of the sizeof or the address (&) operator.

RELATED REFERENCES

- "extern Storage Class Specifier" on page 25
- "Function Declarations" on page 124
- "sizeof (Size of an Object)" on page 91
- "Address &" on page 90

Initializing Arrays

The initializer for an array contains the = symbol followed by a comma-separated list of constant expressions enclosed in braces ({ }). You do not need to initialize all elements in an array. Elements that are not initialized (in extern and static definitions only) receive the value 0 of the appropriate type.

Note: Array initializations can be either *fully braced* (with braces around each dimension) or unbraced (with only one set of braces enclosing the entire set of initializers).

The following definition shows a completely initialized one-dimensional array: static int number[3] = { 5, 7, 2 };

The array number contains the following values:

Element	Value
number[0]	5
number[1]	7
number[2]	2

The following definition shows a partially initialized one-dimensional array: static int number1[3] = { 5, 7 };

The values of number1 are:

Element	Value
number1[0]	5
number1[1]	7
number1[2]	0

Instead of an expression in the subscript declarator defining the number of elements, the following one-dimensional array definition defines one element for each initializer specified:

```
static int item[] = { 1, 2, 3, 4, 5 };
```

The compiler gives item the five initialized elements:

Element	Value
item[0]	1

item[1]	2
item[2]	3
item[3]	4
item[4]	5

You can initialize a one-dimensional character array by specifying:

- A brace-enclosed comma-separated list of constants, each of which can be contained in a character
- A string constant (Braces surrounding the constant are optional)

Initializing a string constant places the null character (\0) at the end of the string if there is room or if the array dimensions are not specified.

The following definitions show character array initializations:

```
static char name1[] = { 'J', 'a', 'n' };
static char name2[] = { "Jan" };
static char name3[4] = "Jan";
```

These definitions create the following elements:

Element	Value	Element	Value	Element	Value
name1[0]	J	name2[0]	J	name3[0]	J
name1[1]	a	name2[1]	a	name3[1]	a
name1[2]	n	name2[2]	n	name3[2]	n
		name2[3]	\0	name3[3]	\0

Note that the following definition would result in the null character being lost: static char name3[3]="Jan";

In C++, when initializing an array of characters with a string, the number of characters in the string — including the terminating '\0' — must not exceed the number of elements in the array.

You can initialize a multidimensional array by:

• Listing the values of all elements you want to initialize, in the order that the compiler assigns the values. The compiler assigns values by increasing the subscript of the last dimension fastest. This form of a multidimensional array initialization looks like a one-dimensional array initialization. The following definition completely initializes the array month_days:

```
static month_days[2][12] =
{
  31, 28, 31, 30, 31, 30, 31, 30, 31, 30, 31,
  31, 29, 31, 30, 31, 30, 31, 30, 31, 30, 31
};
```

Using braces to group the values of the elements you want initialized. You can
put braces around each element, or around any nesting level of elements. The
following definition contains two elements in the first dimension (You can
consider these elements as rows). The initialization contains braces around each
of these two elements:

```
static int month_days[2][12] =
{
    { 31, 28, 31, 30, 31, 30, 31, 30, 31, 30, 31 },
    { 31, 29, 31, 30, 31, 30, 31, 30, 31, 30, 31 }
};
```

Initializers

 Using nested braces to initialize dimensions and elements in a dimension selectively.

The following definition explicitly initializes six elements in a 12-element array: static int matrix[3][4] = {

```
{1, 2},
 \{3, 4\},\
 {5, 6}
};
```

The initial values of matrix are:

Element	Value	Element	Value	
matrix[0][0]	1	matrix[1][2]	0	
matrix[0][1]	2	matrix[1][3]	0	
matrix[0][2]	0	matrix[2][0]	5	
matrix[0][3]	0	matrix[2][1]	6	
matrix[1][0]	3	matrix[2][2]	0	
matrix[1][1]	4	matrix[2][3]	0	

You cannot have more initializers than the number of elements in the array.

RELATED REFERENCES

- "Initializers" on page 56
- "Arrays" on page 62
- "extern Storage Class Specifier" on page 25
- "static Storage Class Specifier" on page 28

Example Programs Using Arrays

The following program defines a floating-point array called prices.

The first for statement prints the values of the elements of prices. The second for statement adds five percent to the value of each element of prices, and assigns the result to total, and prints the value of total.

```
** Example of one-dimensional arrays
**/
#include <stdio.h>
#define ARR_SIZE 5
int main(void)
 static float const prices[ARR SIZE] = { 1.41, 1.50, 3.75, 5.00, .86 };
 auto float total;
 int i;
 for (i = 0; i < ARR SIZE; i++)
   printf("price = $%.2f\n", prices[i]);
 printf("\n");
 for (i = 0; i < ARR_SIZE; i++)
   total = prices[i] * 1.05;
```

```
printf("total = $%.2f\n", total);
  return(0);
This program produces the following output:
price = $1.41
price = $1.50
price = $3.75
price = $5.00
price = $0.86
total = $1.48
total = $1.57
total = $3.94
total = $5.25
total = $0.90
The following program defines the multidimensional array salary_tbl. A for loop
prints the values of salary_tbl.
 ** Example of a multidimensional array
 **/
#include <stdio.h>
#define ROW SIZE
#define COLUMN SIZE 5
int main(void)
  static int salary_tb1[ROW_SIZE][COLUMN_SIZE] =
    { 500, 550, 600, 650, 700
      600, 670, 740, 810, 880
    { 740, 840, 940, 1040, 1140
  int grade, step;
  for (grade = 0; grade < ROW SIZE; ++grade)</pre>
  for (step = 0; step < COLUMN SIZE; ++step)</pre>
     printf("salary_tbl[%d] [%d] = %d\n", grade, step,
             salary_tbl[grade] [step]);
   return(0);
This program produces the following output:
salary tb1[0] [0] = 500
salary_tbl[0][1] = 550
salary_tbl[0] [2] = 600
salary_tb1[0] [3] = 650
salary_tbl[0] [4] = 700
salary_tbl[1] [0] = 600
salary_tbl[1] [1] = 670
salary tbl[1][2] = 740
salary tbl[1][3] = 810
salary_tbl[1] [4] = 880
salary_tbl[2] [0] = 740
```

Initializers

```
salary_tbl[2] [1] = 840
salary_tbl[2] [2] = 940
salary_tbl[2] [3] = 1040
salary_tbl[2] [4] = 1140
```

RELATED REFERENCES

- "Arrays" on page 62
- "Pointers" on page 58
- "Array Subscript [] (Array Element Specification)" on page 80
- "String Literals" on page 20
- "Chapter 4. Declarators" on page 55
- "Initializers" on page 56
- "Chapter 6. Implicit Type Conversions" on page 115

Function Specifiers

The function specifiers **inline** and **virtual** are used only in C++ function declarations.

The function specifier inline is used to make a suggestion to the compiler to incorporate the code of a function into the code at the point of the call.

The function specifier virtual can only be used in nonstatic member function declarations.

RELATED REFERENCES

- "Function Declarations" on page 124
- "Inline Functions" on page 141
- "Virtual Functions" on page 249

References

A reference is an alias or an alternative name for an object. All operations applied to a reference act on the object the reference refers to. The address of a reference is the address of the aliased object.

A reference type is defined by placing the & after the type specifier. You must initialize all references except function parameters when they are defined.

Secause Because arguments of a function are passed by value, a function call does not modify the actual values of the arguments. If a function needs to modify the actual value of an argument, the argument must be passed by reference (as opposed to being passed by value). Passing arguments by reference can be done using either references or pointers. In C++, this is accomplished transparently.

Unlike C, C++ does not force you to use pointers if you want to pass arguments by reference. For example:

```
int f(int&);
void main()
      extern int i;
      f(i);
```

You cannot tell from the function call f(i) that the argument is being passed by reference.

References to NULL are not allowed.

RELATED REFERENCES

- "Objects" on page 24
- "Function Calls ()" on page 78
- "Pointers" on page 58

Initializing References

The object that you use to initialize a reference must be of the same type as the reference, or it must be of a type that is convertible to the reference type. If you initialize a reference to a constant using an object that requires conversion, a temporary object is created. In the following example, a temporary object of type **float** is created:

```
int i;
const float& f = i; // reference to a constant float
```

When you initialize a reference with an object, you *bind* that reference to that object.

Attempting to initialize a nonconstant reference with an object that requires a conversion is an error.

Once a reference has been initialized, it cannot be modified to refer to another object. For example:

Note that the initialization of a reference is not the same as an assignment to a reference. Initialization operates on the actual reference by initializing the reference with the object it is an alias for. Assignment operates through the reference on the object referred to.

A reference can be declared without an initializer:

- When it is used in an parameter declaration
- In the declaration of a return type for a function call
- · In the declaration of class member within its class declaration
- When the extern specifier is explicitly used

You cannot have references to any of the following:

- Other references
- Bit fields
- Arrays of references
- Pointers to references

Direct Binding

Suppose a reference r of type T is initialized by an expression e of type U.

References

The reference r is *bound directly* to e if the following statements are true:

- Expression e is an lvalue
- T is the same type as U, or T is a base class of U
- T has the same, or more, const or volatile qualifiers than U

The reference r is also bound directly to e if e can be implicitly converted to a type such that the previous list of statements is true.

RELATED REFERENCES

- "References" on page 68
- "Declaring Class Types" on page 199
- "Passing Arguments by Reference" on page 136
- "Pointers" on page 58
- "extern Storage Class Specifier" on page 25
- "volatile and const Qualifiers" on page 50
- "Chapter 4. Declarators" on page 55
- "Initializers" on page 56
- "Temporary Objects" on page 273

Chapter 5. Expressions and Operators

Expressions are sequences of operators, operands, and punctuators that specify a computation. The evaluation of expressions is based on the operators that the expressions contain and the context in which they are used.

An expression can result in an Ivalue, rvalue, or no value, and can produce side effects in each case.

C++ operators can be defined to behave differently when applied to operands of class type. This is called operator *overloading*. This chapter describes the behavior of operators that are not overloaded.

RELATED REFERENCES

- "Lvalues and Rvalues" on page 75
- "Overloading Operators" on page 187

Operator Precedence and Associativity

Two operator characteristics determine how operands group with operators: *precedence* and *associativity*. Precedence is the priority for grouping different types of operators with their operands. Associativity is the left-to-right or right-to-left order for grouping operands to operators that have the same precedence. An operator's precedence is meaningful only if other operators with higher or lower precedence are present. Expressions with higher-precedence operators are evaluated first.

For example, in the following statements, the value of 5 is assigned to both a and b because of the right-to-left associativity of the = operator. The value of c is assigned to b first, and then the value of b is assigned to a.

```
b = 9;
c = 5;
a = b = c;
```

Because the order of subexpression evaluation is not specified, you can explicitly force the grouping of operands with operators by using parentheses.

```
In the expression a + b * c / d
```

the * and / operations are performed before + because of precedence. b is multiplied by c before it is divided by d because of associativity.

The following table lists the C and C++ language operators in order of precedence and shows the direction of associativity for each operator.

The C++ scope resolution operator (::) has the highest precedence. The comma operator has the lowest precedence. Operators that have the same rank have the same precedence.

Operator Precedence and Associativity

Rank	Right Associative?	Operator Function	Usage
1	yes	global scope resolution	:: name_or_qualified name
1		c++ class or namespace scope resolition	class_or_namespace :: member
2		member selection	object . member
2		member selection	pointer -> member
2		subscipting	pointer [expr]
2		function call	expr (expr_list)
2		value constructon	type (expr_list)
2		postfix increment	lvalue
2		postfix decrement	lvalue ++
2	yes	□ C++ type identification	typeid (type)
2	yes	type identification at runtime	typeid (expr)
2	yes	compile time	static_cast < type > (expr)
2	yes	conversion checked at runtume	dynamic_cast < type > (expr)
2	yes	□ C++ unchecked conversion	reinterpret_cast < type > (expr)
2	yes	c++ const conversion	const_cast < type > (expr)
3	yes	size of object in bytes	sizeof (expr)
3	yes	size of type in bytes	sizeof type
3	yes	prefix increment	++ lvalue
3	yes	prefix decrement	lvalue
3	yes	complement	[~] expr
3	yes	not	! expr
3	yes	unary minus	- expr
3	yes	unary plus	+ expr
3	yes	address of	& lvalue
3	yes	indirection or dereference	* expr
3	yes	create (allocate memory)	new type
3	yes	create (allocate and initialize memoty)	new type (expr_list) type
3	yes	C++ create (placement)	new type (expr_list) type (expr_list)
3	yes	destroy (deallocate memory)	delete pointer
3	yes	C++ destroy array	delete [] pointer
3	yes	type conversion (cast)	(type) expr
4		member selection	object .* ptr_to_member

Operator Precedence and Associativity

Rank	Right Associative?	Operator Function	Usage
4		member selection	object ->* ptr_to_member
5		multiplication	expr * expr
5		division	expr I expr
5		modulo (remainer)	expr % expr
6		binary addition	expr + expr
6		binary subtraction	expr - expr
7		bitwise shift left	expr << expr
7		bitwise shift right	expr >> expr
8		less than	expr < expr
8		less than or equal to	expr <= expr
8		greater than	expr > expr
8		greater than or equal to	expr >= expr
9		equal	expr == expr
9		not equal	expr != expr
10		bitwise AND	expr & expr
11		bitwise exclusive OR	expr ^ expr
12		bitwise inclusive OR	expr expr
13		logical AND	expr && expr
14		logical inclusive OR	expr expr
16	yes	simplie assignment	lvalue = expr
16	yes	multiply and assign	lvalue *= expr
16	yes	divide and assign	lvalue I= expr
16	yes	modulo and assign	lvalue %= expr
16	yes	add and assign	lvalue += expr
16	yes	subtract and assign	lvalue -= expr
16	yes	shift left and assign	lvalue <<= expr
16	yes	shift right and assign	lvalue >>= expr
16	yes	bitwise AND and assign	lvalue &= expr
16	yes	bitwise exclusive OR and assign	lvalue ^= expr
16	yes	bitwise inclusive OR and assign	lvalue = expr
16		conditional expression	expr ? expr : expr
17	yes	C++ throw expression	throw expr
18		comma (sequencing)	expr, expr

The order of evaluation for function call arguments or for the operands of binary operators is not specified. Avoid writing such ambiguous expressions as:

```
z = (x * ++y) / func1(y);
func2(++i, x[i]);
```

In the example above, ++y and func1(y) might not be evaluated in the same order by all C language implementations. If y had the value of 1 before the first

Operator Precedence and Associativity

statement, it is not known whether or not the value of 1 or 2 is passed to func1(). In the second statement, if i had the value of 1, it is not known whether the first or second array element of x[] is passed as the second argument to func2().

The order of grouping operands with operators in an expression containing more than one instance of an operator with both associative and commutative properties is not specified. The operators that have the same associative and commutative properties are: *, *, *, and $\hat{}$ (or \neg). The grouping of operands can be forced by grouping the expression in parentheses.

Examples of Expressions and Precedence

The parentheses in the following expressions explicitly show how the compiler groups operands and operators.

```
total = (4 + (5 * 3));
total = (((8 * 5) / 10) / 3);
total = (10 + (5/3));
```

If parentheses did not appear in these expressions, the operands and operators would be grouped in the same manner as indicated by the parentheses. For example, the following expressions produce the same output.

```
total = (4+(5*3));
total = 4+5*3;
```

Because the order of grouping operands with operators that are both associative and commutative is not specified, the compiler can group the operands and operators in the expression:

```
total = price + prov tax + city tax;
```

in the following ways (as indicated by parentheses):

```
total = (price + (prov tax + city tax));
total = ((price + prov tax) + city tax);
total = ((price + city_tax) + prov_tax);
```

The grouping of operands and operators does not affect the result unless one ordering causes an overflow and another does not. For example, if price = 32767, prov tax = -42, and city tax = 32767, and all three of these variables have been declared as integers, the third satement total = ((price + city tax) + prov tax) will cause an integer overflow and the rest will not.

Because intermediate values are rounded, different groupings of floating-point operators may give different results.

In certain expressions, the grouping of operands and operators can affect the result. For example, in the following expression, each function call might be modifying the same global variables.

```
a = b() + c() + d();
```

This expression can give different results depending on the order in which the functions are called.

If the expression contains operators that are both associative and commutative and the order of grouping operands with operators can affect the result of the expression, separate the expression into several expressions. For example, the following expressions could replace the previous expression if the called functions do not produce any side effects that affect the variable a.

```
a = b();
a += c();
a += d();
```

Lvalues and Rvalues

Evey expression is either an lvalue or an rvalue.

An *lvalue* is an expression or function that represents an object that can be examined or changed. A *modifiable lvalue* is an expression representing an object that can be changed. It is typically the left operand in an assignment expression. For example, arrays and **const** objects are not modifiable lvalues, but **static int** objects are.

An *rvalue* is an epsression that cannot have a value assigned to it. A function that does not return a reference yields an rvalue. Rvalues always have complete types or the void type.

Certain operators require Ivalues for some of their operands. For example, all assignment operators evaluate their right operand and assign that value to their left operand. The left operand must evaluate to a reference to an object. The address operator (&) requires an Ivalue as an operand while the increment (++) and the decrement (--) operators require a modifiable Ivalue as an operand.

The following are examples of lvalues:

Expression	Lvalue
x = 42	x
*ptr = newvalue	*ptr
a++	a
int& f()	The function calll to f()

RELATED REFERENCES

- "Arrays" on page 62
- "volatile and const Qualifiers" on page 50
- "static Storage Class Specifier" on page 28

Integer Constant Expressions

An *integer compile-time constant* is a value that is determined during compilation and cannot be changed at runtime. An *integer compile-time constant expression* is an expression that is composed of constants and evaluated to a constant.

An integer constant expression is an expression that is composed of only the following:

- literals
- enumerators
- const variables
- static data members of integral or enumeration types
- casts to integral types
- · sizeof expressions

You must use an integer constant expression in the following situations:

• In the subscript declarator as the description of an array bound

Ivalue

- After the keyword case in a switch statement
- In an enumerator, as the numeric value of an enum constant
- · In a bit-field width specifier
- In the preprocessor #if statement (Enumeration constants, address constants, and **sizeof** cannot be specified in the preprocessor **#if** statement)

RELATED REFERENCES

- "Literals" on page 16
- "Enumerations" on page 46
- "volatile and const Qualifiers" on page 50
- "Static Members" on page 218
- "Cast Expressions" on page 108
- "sizeof (Size of an Object)" on page 91

Primary Expressions

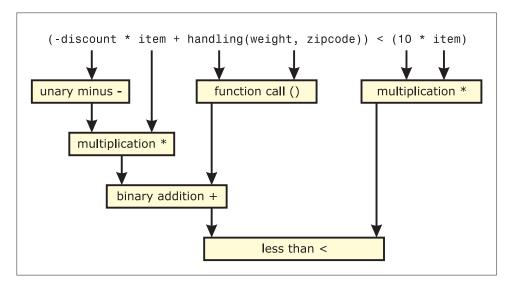
Primary expressions are literals, the C++ this pointer, parenthesized expressions, names, and names qualified by the scope resolution operator (::).

RELATED REFERENCES

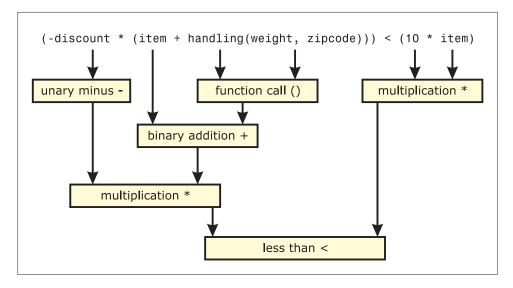
- "Literals" on page 16
- "The this Pointer" on page 216

Parenthesized Expressions ()

Use parentheses to explicitly force the order of expression evaluation. The following expression does not contain any parentheses used for grouping operands and operators. The parentheses surrounding weight, zipcode are used to form a function call. Note how the compiler groups the operands and operators in the expression according to the rules for operator precedence and associativity:



The following expression is similar to the previous expression, but it contains parentheses that change how the operands and operators are grouped:



In an expression that contains both associative and commutative operators, you can use parentheses to specify the grouping of operands with operators. The parentheses in the following expression guarantee the order of grouping operands with the operators:

```
x = f + (g + h);
```

C++ Scope Resolution Operator ::

The :: (scope resolution) operator is used to qualify hidden names so that you can still use them. You can use the unary scope operator if a namespace scope or global scope name is hidden by an explicit declaration of the same name in a block or class. For example:

```
int count = 0;
int main(void) {
  int count = 0;
  ::count = 1; // set global count to 1
  count = 2; // set local count to 2
  return 0;
}
```

The declaration of count declared in the main() function hides the integer named count declared in global namespace scope. The statemennt ::count = 1 accesses the variable named count declared in global namespace scope.

You can also use the class scope operator to qualify class names or class member names. If a class member name is hidden, you can use it by qualifying it with its class name and the class scope operator.

In the following example, the declaration of the variable X hides the class type X, but you can still use the static class member count by qualifying it with the class type X and the scope resolution operator.

```
#include <iostream>
using namespace std;
class X
{
```

```
public:
       static int count;
int X::count = 10;
                                        // define static data member
int main ()
       int X = 0;  // hides class type X
cout << X::count << endl;  // use static member of class X</pre>
```

RELATED REFERENCES

- "Scope of Class Names" on page 203
- "Chapter 10. Namespaces" on page 177

Postfix Expressions

Postfix operators are operators that appear after their operands. A postfix expression is a primary expression, or a primary expression that contains a postfix operator. The following summerizes the available postfix operators:

Rank	Right Associative?	Operator Function	Usage
2		member selection	object . member
2		member selection	pointer -> member
2		subscipting	pointer [expr]
2		function call	expr (expr_list)
2		value constructon	type (expr_list)
2		postfix increment	lvalue
2		postfix decrement	lvalue ++
2	yes	C++ type identification	typeid (type)
2	yes	type identification at runtime	typeid (expr)
2	yes	conversion checked at compile time	static_cast < type > (expr)
2	yes	conversion checked at runtume	dynamic_cast < type > (expr)
2	yes	□ C++ unchecked conversion	reinterpret_cast < type > (expr)
2	yes	C++ const conversion	<pre>const_cast < type > (expr)</pre>

Function Calls ()

A function call is an expression containing a simple type name and a parenthesized argument list. The argument list can contain any number of expressions separated by commas. It can also be empty.

```
For example:
stub()
overdue(account, date, amount)
notify(name, date + 5)
report(error, time, date, ++num)
```

There are two kinds of function calls: ordinary function calls and C++ member function calls. Any function may call itself except for the function main.

Type of a Function Call

The type of a function call expression is the return type of the function. This type can either be a complete type, a reference type, or the type **void**. A function call is an Ivalue if and only if the type of the function is a reference.

Arguments and Parameters

A *function argument* is an expression that you use within the parenthesis of a function call. A *function parameter* is an object or reference declared within the parenthesis of a function declaration or definition. When you call a function, the arguments are evaluated, and each parameter is initialized with the value of the corresponding argument. The semantics of argument passing are identical to those of assignment.

A function can change the values of its non-const parameters, but these changes have no effect on the argument unless the parameter is a reference type.

If you want a function to change the value of a variable, pass a pointer or a reference to the variable you want changed. When a pointer is passed as a parameter, the pointer is copied; the object pointed to is not copied.

Linkage and Function Calls

In C only, if a function definition has external linkage and a return type of **int**, calls to the function can be made before it is explicitly declared because an implicit declaration of extern int func(); is assumed. This is *not* true for C++.

Type Conversions of Arguments

Arguments that are arrays and functions are converted to pointers before being passed as function arguments.

Arguments passed to nonprototyped C functions undergo conversions: type **short** or **char** parameters are converted to **int**, and **float** parameters to **double**. Use a cast expression for other conversions.

The compiler compares the data types provided by the calling function with the data types that the called function expects and performs necessary type conversions. For example, when function funct is called, argument f is converted to a **double**, and argument c is converted to an **int**:

```
char * funct (double d, int i);
    /* ... */
int main(void)
{
    float f;
    char c;
       funct(f, c) /* f is a double, c is an int */
    return 0;
}
```

Evaluation Order of Arguments

The order in which arguments are evaluated is not specified. Avoid such calls as:

```
method(sample1, batch.process--, batch.process);
```

In this example, batch.process-- might be evaluated last, causing the last two arguments to be passed with the same value.

Example of Function Calls

In the following example, main passes func two values: 5 and 7. The function func receives copies of these values and accesses them by the identifiers: a and b. The function func changes the value of a. When control passes back to main, the actual values of x and y are not changed. The called function func only receives copies of the values of x and y, not the variables themselves.

```
** This example illustrates function calls
**/
#include <stdio.h>
void func (int a, int b)
  a += b;
  printf("In func, a = %d b = %d\n", a, b);
int main(void)
   int x = 5, y = 7;
  func(x, y);
printf("In main, x = %d y = %d\n", x, y);
  return 0:
```

This program produces the following output:

```
In func, a = 12 b = 7
In main, x = 5 y = 7
```

RELATED REFERENCES

- "Chapter 7. Functions" on page 123
- "Cast Expressions" on page 108
- "Pointers" on page 58
- "Program Linkage" on page 4

Array Subscript [] (Array Element Specification)

A postfix expression followed by an expression in [] (square brackets) specifies an element of an array. The expression within the square brackets is referred to as a subscript.

The expression a[b] is equivalent (by definition) to the expression *((a) + (b))(and because addition is associative, it is also equivalent to b[a]). Between expressions a and b, one must be a pointer to a type T, and the other must have integral or enumeration type. The result of an array subscript is an Ivalue. The following example demonstrates this:

```
#include <stdio.h>
int main(void) {
 int a[3] = \{ 10, 20, 30 \};
 printf("a[0] = %d\n", a[0]);
```

```
printf("a[1] = %d\n", 1[a]);
printf("a[2] = %d\n", *(2 + a));
return 0;
```

The following is the output of the above example:

```
a[1] = 20
a[2] = 30
```

The above restrictions on the types of expressions required by the subscript operator, as well as the relationship between the subscript operator and pointer arithmetic, do not apply if you overload operator[] of a class.

The first element of each array has the subscript 0. The expression contract[35] refers to the 36th element in the array contract.

In a multidimensional array, you can reference each element (in the order of increasing storage locations) by incrementing the right-most subscript most frequently.

For example, the following statement gives the value 100 to each element in the array code[4][3][6]:

```
for (first = 0; first < 4; ++first)
  for (second = 0; second < 3; ++second)
     for (third = 0; third < 6; ++third)
        code[first][second][third] = 100;
```

RELATED REFERENCES

- "Pointers" on page 58
- "Integer Variables" on page 34
- "Lvalues and Rvalues" on page 75
- "Arrays" on page 62
- "Overloading Subscripting" on page 193
- "Pointer Arithmetic" on page 60

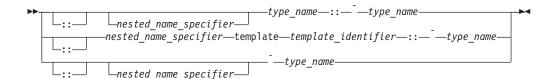
Dot Operator .

The . (dot) operator is used to access class, structure, or union members. The member is specified by a postfix expression, followed by a . (dot) operator, followed by a possibly qualified identifier or a pseudo-destructor name. The postfix expression must be an object of type class, struct or union. The name must be a member of that object.

The value of the expression is the value of the selected member. If the postfix expression and the name are lvalues, the expression value is also an lvalue.

Pseudo-destructors

A pseudo-destructor is a destructor of a non-class type named type_name in the following syntax diagram:



RELATED REFERENCES

- "Chapter 13. Class Members and Friends" on page 209
- "Unions" on page 42
- "Structures" on page 36
- "Scope of Class Names" on page 203

Arrow Operator ->

The -> (arrow) operator is used to access class, structure or union members using a pointer. A postfix expression, followed by an -> (arrow) operator, followed by a possibly qualified identifier or a pseudo-destructor name, designates a member of the object to which the pointer points. (A pseudo-destructor is a destructor of a non-class type.) The postfix expression must be a pointer to an object of type class, **struct** or **union**. The name must be a member of that object.

The value of the expression is the value of the selected member. If the name is an lvalue, the expression value is also an lvalue.

RELATED REFERENCES

- "Pointers" on page 58
- "Chapter 13. Class Members and Friends" on page 209
- "Unions" on page 42
- "Structures" on page 36
- "Dot Operator ." on page 81

static_cast Operator

The static_cast operator converts a given expression to a specified type.

```
Syntax — static_cast
```

```
▶ static cast <-- Type --> -- (--expression -- ) --
```

The following is an example of the **static_cast** operator.

```
#include <iostream>
using namespace std;
int main() {
 int j = 41;
 int v = 4;
 float m = j/v;
 float d = static cast<float>(j)/v;
 cout << "m = " << m << endl;
 cout << "d = " << d << endl;
```

The following is the output of the above example:

m = 10d = 10.25

In this example, m = j/v; produces an answer of type int because both j and v are integers. Conversely, d = static_cast<float>(j)/v; produces an answer of type float. The static_cast operator converts variable j to a type float. This allows the compiler to generate a division with an answer of type float. All static_cast operators resolve at compile time and do not remove any const or volatile modifiers.

Applying the **static_cast** operator to a null pointer will convert it to a null pointer value of the target type.

You can explicitly convert a pointer of a type A to a pointer of a type B if A is a base class of B. If A is not a base class of B, a compiler error will result.

You may cast an Ivalue of a type A to a type B& if the following are true:

- A is a base class of B
- You are able to convert a pointer of type A to a pointer of type B
- The type B has the same or greater **const** or **volatile** qualifiers than type A
- A is not a virtual base class of B

The result is an Ivalue of type B.

A pointer to member type can be explicitly converted into a different pointer to member type if both types are pointers to members of the same class. This form of explicit conversion may also take place if the pointer to member types are from separate classes, however one of the class types must be derived from the other.

reinterpret_cast Operator

A reinterpret_cast operator handles conversions between unrelated types.

Syntax — reinterpret_cast

▶─reinterpret_cast-<-*Type*->-(-expression-)------

The reinterpret_cast operator produces a value of a new type that has the same bit pattern as its argument. You cannot cast away a const or volatile qualification. You can explicitly perform the following conversions:

- · A pointer to any integral type large enough to hold it
- A value of integral or enumeration type to a pointer
- · A pointer to a function to a pointer to a function of a different type
- A pointer to an object to a pointer to an object of a different type
- A pointer to a member to a pointer to a member of a different class or type, if the types of the members are both function types or object types

A null pointer value is converted to the null pointer value of the destination type.

Given an Ivalue expression of type T and an object x, the following two conversions are synonymous:

- reinterpret_cast<T&>(x)
- *reinterpret cast<T*>(&x)

const_cast Operator

C++= A const_cast operator is used to add or remove a const or volatile modifier to or from a type.

Syntax — const_cast

```
▶►—const_cast—<—Type—>—(—expression—)-
```

Type and the type of expression may only differ with respect to their const and volatile qualifiers. Their cast is resolved at compile time. A single const_cast expression may add or remove any number of const or volatile modifiers.

The result of a **const_cast** expression is an rvalue unless *Type* is a reference type. In this case, the result is an Ivalue.

Types can not be defined within **const_cast**.

The following demonstrates the use of the **const_cast** operator:

```
#include <iostream>
using namespace std;
void f(int* p) {
 cout << *p << endl;
int main(void) {
 const int a = 10:
 const int* b = &a;
 // Function f() expects int*, not const int*
 // f(b);
 int* c = const_cast<int>(b);
 f(c);
 // Lvalue is const
 // *b = 20;
  // Undefined behavior
 // *c = 30;
 int a1 = 40;
 const int* b1 = &a1;
 int* c1 = const_cast<int>(b1);
 // Integer al, the object referred to by cl, has
 // not been declared const
 *c1 = 50;
 return 0;
```

The compiler will not allow the function call f(b). Function f() expects a pointer to an int, not a const int. The statement int* c = const_cast<int>(b) returns a pointer c that refers to a without the const qualification of a. This process of using const_cast to remove the const qualification of an object is called casting away constness. Consequently the compiler will allow the function call f(c).

The compiler would not allow the assignment *b = 20 because b points to an object of type **const int**. The compiler will allow the *c = 30, but the behavior of this statement is undefined. If you cast away the constness of an object that has been explicitly declared as **const**, and attempt to modify it, the results are undefined.

However, if you cast away the constess of an object that has not been explicity declared as **const**, you can modify it safely. In the above example, the object referred to by b1 has not been declared **const**, but you cannot modify this object through b1. You may cast away the constness of b1 and modify the value to which it referrs.

dynamic_cast Operator

The *dynamic_cast* operator performs type conversions at runtime.

The expression dynamic_cast<T>(v) converts the expression v to type T. Type T must be a pointer or reference to a complete class type or a pointer to void. If T is a pointer and the **dynamic_cast** operator fails, the operator returns a null pointer of type T. If T is a reference and the **dynamic_cast** operator fails, the operator throws the exception **std::bad_cast**. You can find this class in the standard library header **<typeinfo>**.

If T is a void pointer, then **dynamic_cast** will return the starting address of the object pointed to by v. The following example demonstrates this:

```
#include <iostream>
using namespace std;

struct A {
   virtual A() { };
};

struct B : A { };

int main() {
   B bobj;
   A* ap = &bobj;
   void * vp = dynamic_cast<void>(ap);
   cout << "Address of vp : " << vp << endl;
   cout << "Address of bobj: " << &bobj << endl;
}</pre>
```

The output of this example will be similar to the following. Both vp and &bobj will refer to the same address:

```
Address of vp : 12FF6C
Address of bobj: 12FF6C
```

The primary purpose for the **dynamic_cast** operator is to perform type-safe *downcasts*. A downcast is the conversion of a pointer or reference to a class A to pointer or reference to a class B, where class A is a base class of B. The problem with downcasts is that a pointer of type A* can and must point to any object of a class that has been derived from A. The **dynamic_cast** operator ensures that if you convert a pointer of class A to a pointer of a class B, the object that A points to belongs to class B or a class derived from B.

The following example demonstrates the use of the **dynamic_cast** operator:

```
#include <iostream>
using namespace std;
```

```
struct A {
 virtual void f() { cout << "Class A" << endl; }</pre>
struct B : A {
 virtual void f() { cout << "Class B" << endl; }</pre>
struct C : A {
 virtual void f() { cout << "Class C" << endl; }</pre>
void f(A* arg) {
 B* bp = dynamic_cast<b>(arg);
 C* cp = dynamic_cast<c>(arg);
 if (bp)
   bp->f();
 else if (cp)
    cp->f();
 else
    arg->f();
};
int main() {
 A aobj;
 C cobj;
 A* ap = \&cobj;
 A* ap2 = &aobj
 f(ap);
 f(ap2);
```

The following is the output of the above example:

Class C Class A

The function f() determines whether the pointer arg points to an object of type A, B, or C. The function does this by trying to convert arg to a pointer of type B, then to a pointer of type C, with the dynamic_cast operator. If the dynamic_cast operator succeeds, it returns a pointer that points to the object denoted by arg. If dynamic cast fails, it returns 0.

You may perform downcasts with the dynamic_cast operator only on polymorphic classes. In the above example, all the classes are polymorphic because class A has a virtual function. The dynamic_cast operator uses the runtime type information generated from polymorphic classes.

VAC++ You must indicate that you want the compiler to generate runtime type information with a compiler option.

Unary Expressions

A unary expression contains one operand and a unary operator. All unary operators have the same precedence and have right-to-left associativity.

As indicated in the following descriptions, the usual arithmetic conversions are performed on the operands of most unary expressions.

Unary Expressions

The following table summarizes the operators for unary expressions:

Rank	Right Associative?	Operator Function	Usage
3	yes	size of object in bytes	sizeof (expr)
3	yes	size of type in bytes	sizeof type
3	yes	prefix increment	++ lvalue
3	yes	prefix decrement	lvalue
3	yes	complement	[~] expr
3	yes	not	! expr
3	yes	unary minus	- expr
3	yes	unary plus	+ expr
3	yes	address of	& lvalue
3	yes	indirection or dereference	* expr
3	yes	create (allocate memory)	new type
3	yes	create (allocate and initialize memoty)	new type (expr_list) type
3	yes	create (placement)	new type (expr_list) type (expr_list)
3	yes	destroy (deallocate memory)	delete pointer
3	yes	C++ destroy array	delete [] pointer
3	yes	type conversion (cast)	(type) expr

In addition, postfix expressions are also unary expressions.

RELATED REFERENCES

- "Increment ++"
- "Decrement --" on page 88
- "Unary Plus +" on page 88
- "Unary Minus -" on page 89
- "Arithmetic Conversions" on page 121

Increment ++

The ++ (increment) operator adds 1 to the value of a scalar operand, or if the operand is a pointer, increments the operand by the size of the object to which it points. The operand receives the result of the increment operation. The operand must be a modifiable lvalue of arithmetic or pointer type.

You can put the ++ before or after the operand. If it appears before the operand, the operand is incremented. The incremented value is then used in the expression. If you put the ++ after the operand, the value of the operand is used in the expression before the operand is incremented. For example:

$$play = ++play1 + play2++;$$

is similar to the following expressions; play2 is altered before play:

Unary Expressions

```
int temp, temp1, temp2;
temp1 = play1;
temp2 = play2 + 1;
play1 = play1 + 1;
temp = temp1 + temp2;
play2 = temp2;
play = temp;
```

The result has the same type as the operand after integral promotion.

The usual arithmetic conversions on the operand are performed.

RELATED REFERENCES

- "Pointer Arithmetic" on page 60
- "Lvalues and Rvalues" on page 75
- "Arithmetic Conversions" on page 121

Decrement --

The -- (decrement) operator subtracts 1 from the value of a scalar operand, or if the operand is a pointer, decreases the operand by the size of the object to which it points. The operand receives the result of the decrement operation. The operand must be a modifiable lvalue.

You can put the -- before or after the operand. If it appears before the operand, the operand is decremented, and the decremented value is used in the expression. If the -- appears after the operand, the current value of the operand is used in the expression and the operand is decremented.

```
For example:
```

```
play = --play1 + play2--;
```

is similar to the following expressions; play2 is altered before play:

```
int temp, temp1, temp2;
temp1 = play1;
temp2 = play2 - 1;
play1 = play1 - 1;
temp = temp1 + temp2;
play2 = temp2;
play = temp;
```

The result has the same type as the operand after integral promotion, but is not an lvalue.

The usual arithmetic conversions are performed on the operand.

RELATED REFERENCES

- "Pointer Arithmetic" on page 60
- "Lvalues and Rvalues" on page 75
- "Arithmetic Conversions" on page 121

Unary Plus +

The + (unary plus) operator maintains the value of the operand. The operand can have any arithmetic type or pointer type. The result is not an Ivalue.

The result has the same type as the operand after integral promotion.

Note: Any plus sign in front of a constant is not part of the constant.

RELATED REFERENCES

"Lvalues and Rvalues" on page 75

Unary Minus -

The - (unary minus) operator negates the value of the operand. The operand can have any arithmetic type. The result is not an Ivalue.

For example, if quality has the value 100, -quality has the value -100.

The result has the same type as the operand after integral promotion.

Note: Any minus sign in front of a constant is not part of the constant.

RELATED REFERENCES

• "Lvalues and Rvalues" on page 75

Logical Negation!

The ! (logical negation) operator determines whether the operand evaluates to θ (false) or nonzero (true).

The expression yields the value 1 (true) if the operand evaluates to θ , and yields the value θ (false) if the operand evaluates to a nonzero value.

The expression yields the value true if the operatand evaluate to false (0), and yields the value false if the operand evaluates to true (nonzero). The operand is implicitly converted to **bool** and the type of the result is **bool**.

The following two expressions are equivalent:

```
!right;
right == 0;
```

RELATED REFERENCES

- "Lvalues and Rvalues" on page 75
- "Boolean Variables" on page 33

Bitwise Negation ~

The ~ (bitwise negation) operator yields the bitwise complement of the operand. In the binary representation of the result, every bit has the opposite value of the same bit in the binary representation of the operand. The operand must have an integral type. The result has the same type as the operand but is not an lvalue.

Suppose x represents the decimal value 5. The 16-bit binary representation of x is: 000000000000101

The expression ~x yields the following result (represented here as a 16-bit binary number):

1111111111111010

Note that the ~ character can be represented by the trigraph ??-.

The 16-bit binary representation of ~0 is:

Unary Expressions

111111111111111111

RELATED REFERENCES

- "Lvalues and Rvalues" on page 75
- "Trigraph Sequences" on page 12

Address &

The & (address) operator yields a pointer to its operand. The operand must be an lvalue, a function designator, or a qualified name. It cannot be a bit field, nor can it have the storage class register.

You may take the address of a register variable.

If the operand is an Ivalue or function, the resulting type is a pointer to the expression type. For example, if the expression has type int, the result is a pointer to an object having type int.

If the operand is a qualified name and the member is not static, the result is a pointer to a member of class and has the same type as the member. The result is not an lvalue.

If p_to_y is defined as a pointer to an **int** and y as an **int**, the following expression assigns the address of the variable y to the pointer p_to_y:

$$p_to_y = &y$$

You can use the & operator with overloaded functions only in an initialization or assignment where the left side uniquely determines which version of the overloaded function is used.

RELATED REFERENCES

- "Lvalues and Rvalues" on page 75
- "Chapter 7. Functions" on page 123
- "Pointers" on page 58
- "Overloading Functions" on page 185
- "register Storage Class Specifier" on page 27

Indirection *

The * (indirection) operator determines the value referred to by the pointer-type operand. The operand cannot be a pointer to an incomplete type. The operation yields an Ivalue or a function designator if the operand points to a function. Arrays and functions are converted to pointers.

The type of the operand determines the type of the result. For example, if the operand is a pointer to an **int**, the result has type **int**.

Do not apply the indirection operator to any pointer that contains an address that is not valid, such as NULL. The result is not defined.

If p to y is defined as a pointer to an int and y as an int, the expressions:

cause the variable y to receive the value 3.

RELATED REFERENCES

- "Arrays" on page 62
- "Chapter 7. Functions" on page 123
- "Pointers" on page 58

sizeof (Size of an Object)

The **sizeof** operator yields the size in *bytes* of the operand. Types cannot be defined in a **sizeof** expression. The **sizeof** operation cannot be performed on

- · A bit field
- · A function
- An undefined structure or class
- An incomplete type (such as **void**)

The operand can be the parenthesized name of a type or expression.

The compiler must be able to evaluate the size at compile time. The expression is not evaluated; there are no side effects. For example, the value of b is 5 from initialization to the end of program runtime:

```
#include <stdio.h>
int main(void){
  int b = 5;
  sizeof(b++);
  return 0;
}
```

The result is an integer constant.

The size of a **char** object is the size of a byte. For example, if a variable x has type **char**, the expression sizeof(x) always evaluates to 1.

The result of a size of operation has type size_t, which is an unsigned integral type defined in the <stddef.h> header.

The size of an object is determined on the basis of its definition. The **sizeof** operator does not perform any conversions. If the operand contains operators that perform conversions, the compiler does take these conversions into consideration. The following expression causes the usual arithmetic conversions to be performed. The result of the expression x + 1 has type **int** (if x has type char, **short**, or **int** or any enumeration type) and is equivalent to sizeof(int):

```
sizeof (x + 1);
```

Except in preprocessor directives, you can use a **sizeof** expression wherever an integral constant is required. One of the most common uses for the **sizeof** operator is to determine the size of objects that are referred to during storage allocation, input, and output functions.

Another use of **sizeof** is in porting code across platforms. You should use the **sizeof** operator to determine the size that a data type represents. For example: sizeof(int);

The result of a size of expression depends on the type it is applied to:

An array The result is the total number of bytes in the array. For example, in an array with 10 elements, the size is equal to 10 times the size of

Unary Expressions

a single element. The compiler does not convert the array to a

pointer before evaluating the expression.

A class The result is always nonzero, and is equal to the number of bytes

in an object of that class including any padding required for

placing class objects in an array.

A reference The result is the size of the referenced object.

RELATED REFERENCES

"Integer Constant Expressions" on page 75

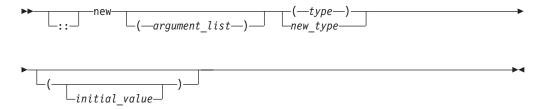
"Arrays" on page 62

• "Chapter 12. Classes" on page 199

"References" on page 68

C++ new Operator

Example 1 The **new** operator provides dynamic storage allocation. The syntax for an allocation expression containing the new operator is:



If you prefix **new** with the scope resolution operator (::), the **global operator new()** is used. If you specify an argument_list, the overloaded **new** operator that corresponds to that argument_list is used. The type is an existing built-in or user-defined type. A new_type is a type that has not already been defined and can include type specifiers and declarators.

An allocation expression containing the **new** operator is used to find storage in free store for the object being created. The new expression returns a pointer to the object created and can be used to initialize the object. If the object is an array, a pointer to the initial element is returned.

You can use **set_new_handler()** only to specify what **new** does when it fails.

You cannot use the **new** operator to allocate function types, **void**, or incomplete class types because these are not object types. However, you can allocate pointers to functions with the **new** operator. You cannot create a reference with the **new** operator.

When the object being created is an array, only the first dimension can be a general expression. All subsequent dimensions must be constant integral expressions. The first dimension can be a general expression even when an existing type is used. You can create an array with zero bounds with the **new** operator. For example: char * c = new char [0];

In this case, a pointer to a unique object is returned.

An object created with **operator new()** or **operator new[]()** exists until the **operator delete()** or **operator delete[]()** is called to deallocate the object's memory. A **delete** operator or a destructor will not be implicitly called for an object created with a **new** that has not been explicitly deallocated before the the end of the program.

If parentheses are used within a new type, parentheses should also surround the new type to prevent syntax errors.

In the following example, storage is allocated for an array of pointers to functions:

```
void f();
void g();

int main(void)
{
    void (**p)(), (**q)();
    // declare p and q as pointers to pointers to void functions
    p = new (void (*[3])());
    // p now points to an array of pointers to functions
    q = new void(*[3])(); // error
    // error - bound as 'q = (new void) (*[3])();'
    p[0] = f; // p[0] to point to function f
    q[2] = g; // q[2] to point to function g
    p[0](); // call f()
    q[2](); // call g()
    return (0);
}
```

However, the second use of **new** causes an erroneous binding of $q = (new \ void) \ (*[3])()$.

The type of the object being created cannot contain class declarations, enumeration declarations, or **const** or **volatile** types. It can contain pointers to **const** or **volatile** objects.

For example, const char* is allowed, but char* const is not.

Additional arguments can be supplied to **new** by using the *argument_list*, also called the *placement syntax*. If placement arguments are used, a declaration of **operator new()** or **operator new[]()** with these arguments must exist. For example:

```
#include <new>
using namespace std;

class X
{
public:
    void* operator new(size_t,int, int){ /* ... */ }
};

// ...

int main ()
{
    X* ptr = new(1,2) X;
}
```

RELATED REFERENCES

- "C++ Scope Resolution Operator ::" on page 77
- "Free Store" on page 269
- "set_new_handler() Set Behavior for new Failure" on page 94
- "C++ delete Operator" on page 95

Unary Expressions

- "Constructors and Destructors Overview" on page 257
- "Objects" on page 24
- "Integer Constant Expressions" on page 75

Initializing Objects Created with the new Operator

You can initialize objects created with the **new** operator in several ways. For nonclass objects, or for class objects without constructors, a new initializer expression can be provided in a new expression by specifying (expression) or (). For example:

```
double* pi = new double(3.1415926);
int* score = new int(89);
float* unknown = new float();
```

If a class does not have a default constructor, the new initializer must be provided when any object of that class is allocated. The arguments of the new initializer must match the arguments of a constructor.

You cannot specify an initializer for arrays. You can initialize an array of class objects only if the class has a default constructor. The constructor is called to initialize each array element (class object).

Initialization using the new initializer is performed only if **new** successfully allocates storage.

RELATED REFERENCES

- "Free Store" on page 269
- "Constructors and Destructors Overview" on page 257

set_new_handler() — Set Behavior for new Failure

When the **new** operator creates a new object, it calls the **operator new()** or **operator new**[]() function to obtain the needed storage.

When **new** cannot allocate storage to create a new object, it calls a new handler function if one has been installed by a call to **set_new_handler()**. The std::set_new_handler() function is declared in the header <new>. Use it to call a new handler you have defined or the default new handler.

Your new handler must perform one of the following:

- obtain more storage for memory allocation, then return
- throw an exception of type std::bad alloc or a class derived from std::bad alloc
- call either abort() or exit()

The **set_new_handler()** function has the prototype:

```
typedef void(*PNH)()
PNH set new handler(PNH);
```

set_new_handler() takes as an argument a pointer to a function (the new handler), which has no arguments and returns void. It returns a pointer to the previous new handler function.

If you do not specify your own set_new_handler() function, new throws an exception of type **std::bad_alloc**.

The following program fragment shows how you could use set_new_handler() to return a message if the **new** operator cannot allocate storage:

```
#include <iostream>
#include <new>
#include <cstdlib>
using namespace std;

void no_storage()
{
    std::cerr << "Operator new failed: no storage is available.\n";
        std::exit(1);
}
int main(void)
{
    std::set_new_handler(&no_storage);
    // Rest of program ...
}</pre>
```

If the program fails because **new** cannot allocate storage, the program exits with the message:

Operator new failed: no storage is available.

RELATED REFERENCES

- "C++ new Operator" on page 92
- "Free Store" on page 269

C++ delete Operator

The **delete** operator destroys the object created with **new** by deallocating the memory associated with the object.

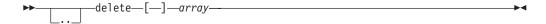
The **delete** operator has a **void** return type. It has the syntax:



The operand of **delete** must be a pointer returned by **new**, and cannot be a pointer to constant. Deleting a null pointer has no effect.

The **delete**[] operator frees storage allocated for array objects created with **new**[]. The **delete** operator frees storage allocated for individual objects created with **new**.

It has the syntax:



The result of deleting an array object with **delete** is undefined, as is deleting an individual object with **delete**[]. The array dimensions do not need to be specified with **delete**[].

The result of any attempt to access a deleted object or array is undefined.

If a destructor has been defined for a class, **delete** invokes that destructor. Whether a destructor exists or not, **delete** frees the storage pointed to by calling the function **operator delete()** of the class if one exists.

Unary Expressions

The global **::operator delete()** is used if:

- The class has no **operator delete()**.
- The object is of a nonclass type.
- The object is deleted with the **::delete** expression.

The global **::operator delete**[]() is used if:

- The class has no operator delete[]()
- The object is of a nonclass type
- The object is deleted with the ::delete[] expression.

The default global **operator delete()** only frees storage allocated by the default global operator new(). The default global operator delete[]() only frees storage allocated for arrays by the default global operator new[]().

RELATED REFERENCES

- "Free Store" on page 269
- "Constructors and Destructors Overview" on page 257
- "void Type" on page 35

Allocation and Deallocation Functions

You may define your own new operator or allocation function as a class member function or a global namespace function with the following restrictions:

- The first parameter must be of type std::size_t. It cannot have a default parameter.
- The return type must be of type **void***.
- Your allocation function may be a template function. Neither the first parameter nor the return type may depend on a template parameter.
- If you declare your allocation function with the empty exception specification throw(), your allocation function must return a null pointer your function fails. Otherwise, your function must throw an exception of type std::bad_alloc or a class derived from **std::bad_alloc** if your function fails.

You may define your own delete operator or deallocation function as a class member function or a global namespace function with the following restrictions:

- The first parameter must be of type **void***.
- The return type must be of type **void**.
- Your dellocation function may be a template function. Neither the first parameter nor the return type may depend on a template parameter.

The following example defines replacement functions for global namespace new and delete:

```
#include <cstdio>
#include <cstdlib>
using namespace std;
void* operator new(size t sz) {
 printf("operator new with %d bytes\n", sz);
 void* p = malloc(sz);
 if (p == 0) printf("Memory error\n");
 return p;
void operator delete(void* p) {
 if (p == 0) printf ("Deleting a null pointer\n");
    printf("delete object\n");
    free(p);
```

```
struct A {
  const char* data;
  A() : data("Text String") { printf("Constructor of S\n"); }
  A() { printf("Destructor of S\n"); }
};

int main() {
  A* ap1 = new A;
  delete ap1;

  printf("Array of size 2:\n");
  A* ap2 = new A[2];
  delete[] ap2;
}
```

The following is the output of the above example:

```
operator new with 40 bytes operator new with 33 bytes operator new with 4 bytes Constructor of S Destructor of S delete object Array of size 2: operator new with 16 bytes Constructor of S Constructor of S Destructor of S Destructor of S delete object
```

RELATED REFERENCES

• "Free Store" on page 269

Binary Expressions

A binary expression contains two operands separated by one operator.

Not all binary operators have the same precedence.

All binary operators have left-to-right associativity.

The order in which the operands of most binary operators are evaluated is not specified. To ensure correct results, avoid creating binary expressions that depend on the order in which the compiler evaluates the operands.

As indicated in the following descriptions, the usual arithmetic conversions are performed on the operands of most binary expressions.

The following table summarizes the operators for binary expressions:

Rank	Right Associative?	Operator Function	Usage
5		multiplication	expr * expr
5		division	expr I expr
5		modulo (remainer)	expr % expr
6		binary addition	expr + expr

Rank	Right Associative?	Operator Function	Usage
6		binary subtraction	expr - expr
7		bitwise shift left	expr << expr
7		bitwise shift right	expr >> expr
8		less than	expr < expr
8		less than or equal to	expr <= expr
8		greater than	expr > expr
8		greater than or equal to	expr >= expr
9		equal	expr == expr
9		not equal	expr != expr
10		bitwise AND	expr & expr
11		bitwise exclusive OR	expr ^ expr
12		bitwise inclusive OR	expr expr
13		logical AND	expr && expr
14		logical inclusive OR	expr expr
16	yes	simplie assignment	lvalue = expr
16	yes	multiply and assign	lvalue *= expr
16	yes	divide and assign	lvalue I= expr
16	yes	modulo and assign	lvalue %= expr
16	yes	add and assign	lvalue += expr
16	yes	subtract and assign	lvalue -= expr
16	yes	shift left and assign	lvalue <<= expr
16	yes	shift right and assign	lvalue >>= expr
16	yes	bitwise AND and assign	lvalue &= expr
16	yes	bitwise exclusive OR and assign	lvalue ^= expr
16	yes	bitwise inclusive OR and assign	lvalue I = expr
18		comma (sequencing)	expr, expr

RELATED REFERENCES

- "Operator Precedence and Associativity" on page 71
- "Arithmetic Conversions" on page 121

Multiplication *

The * (multiplication) operator yields the product of its operands. The operands must have an arithmetic or enumeration type. The result is not an Ivalue. The usual arithmetic conversions on the operands are performed.

Because the multiplication operator has both associative and commutative properties, the compiler can rearrange the operands in an expression that contains more than one multiplication operator. For example, the expression:

```
sites * number * cost
```

can be interpreted in any of the following ways:

```
(sites * number) * cost
sites * (number * cost)
(cost * sites) * number
```

RELATED REFERENCES

- "Lvalues and Rvalues" on page 75
- "Arithmetic Conversions" on page 121

Division /

The / (division) operator yields the quotient of its operands. The operands must have an arithmetic or enumeration type. The result is not an Ivalue.

If both operands are positive integers and the operation produces a remainder, the remainder is ignored. For example, expression 7 / 4 yields the value 1 (rather than 1.75 or 2).

On all IBM C and C++ compilers, if either operand is negative, the result is rounded towards zero.

The result is undefined if the second operand evaluates to 0.

The usual arithmetic conversions on the operands are performed.

RELATED REFERENCES

- "Lvalues and Rvalues" on page 75
- "Arithmetic Conversions" on page 121

Remainder %

The % (remainder) operator yields the remainder from the division of the left operand by the right operand. For example, the expression 5 % 3 yields 2. The result is not an Ivalue.

Both operands must have an integral or enumeration type. If the right operand evaluates to θ , the result is undefined. If either operand has a negative value, the result is such that the following expression always yields the value of a if b is not θ and a/b is representable:

```
(a/b)*b+a%b;
```

The sign of the remainder is the same as the sign of the quotient.

The usual arithmetic conversions on the operands are performed.

RELATED REFERENCES

"Arithmetic Conversions" on page 121

Addition +

The + (addition) operator yields the sum of its operands. Both operands must have an arithmetic type, or one operand must be a pointer to an object type and the other operand must have an integral or enumeration type.

When both operands have an arithmetic type, the usual arithmetic conversions on the operands are performed. The result has the type produced by the conversions on the operands and is not an Ivalue.

A pointer to an object in an array can be added to a value having integral type. The result is a pointer of the same type as the pointer operand. The result refers to another element in the array, offset from the original element by the amount of the integral value treated as a subscript. If the resulting pointer points to storage outside the array, other than the first location outside the array, the result is undefined. The compiler does not provide boundary checking on the pointers. For example, after the addition, ptr points to the third element of the array:

```
int array[5];
int *ptr;
ptr = array + 2;
```

RELATED REFERENCES

- "Pointer Arithmetic" on page 60
- "Pointer Conversions" on page 118

Subtraction -

The - (subtraction) operator yields the difference of its operands. Both operands must have an arithmetic or enumeration type, or the left operand must have a pointer type and the right operand must have the same pointer type or an integral or enumeration type. You cannot subtract a pointer from an integral value.

When both operands have an arithmetic type, the usual arithmetic conversions on the operands are performed. The result has the type produced by the conversions on the operands and is not an Ivalue.

When the left operand is a pointer and the right operand has an integral type, the compiler converts the value of the right to an address offset. The result is a pointer of the same type as the pointer operand.

If both operands are pointers to the same type, the compiler converts the result to an integral type that represents the number of objects separating the two addresses. Behavior is undefined if the pointers do not refer to objects in the same array.

RELATED REFERENCES

- "Pointer Arithmetic" on page 60
- "Pointer Conversions" on page 118

Bitwise Left and Right Shift << >>

The bitwise shift operators move the bit values of a binary object. The left operand specifies the value to be shifted. The right operand specifies the number of positions that the bits in the value are to be shifted. The result is not an Ivalue. Both operands have the same precedence and are left-to-right associative.

Operator	Usage
<<	Indicates the bits are to be shifted to the left.
>>	Indicates the bits are to be shifted to the right.

Each operand must have an integral or enumeration type. The compiler performs integral promotions on the operands, and then the right operand is converted to type int. The result has the same type as the left operand (after the arithmetic conversions).

The right operand should not have a negative value or a value that is greater than or equal to the width in bits of the expression being shifted. The result of bitwise shifts on such values is unpredictable.

If the right operand has the value 0, the result is the value of the left operand (after the usual arithmetic conversions).

The << operator fills vacated bits with zeros. For example, if left_op has the value 4019, the bit pattern (in 16-bit format) of left_op is: 0000111110110011

The expression left_op << 3 yields: 0111110110011000

Relational < > <= >=

The relational operators compare two operands and determine the validity of a relationship.

The type of the result is **int** and has the values 1 if the specified relationship os true, and 0 if false.

The type of the result is **bool** and has the values **true** or **false**.

The result is not an Ivalue.

The following table describes the four relational operators:

Operator	Usage
<	Indicates whether the value of the left operand is less than the value of the right operand.
>	Indicates whether the value of the left operand is greater than the value of the right operand.
<=	Indicates whether the value of the left operand is less than or equal to the value of the right operand.
>=	Indicates whether the value of the left operand is greater than or equal to the value of the right operand.

Both operands must have arithmetic or enumeration types or be pointers to the same type.

The result has type int.

The result has type **bool**.

If the operands have arithmetic types, the usual arithmetic conversions on the operands are performed.

When the operands are pointers, the result is determined by the locations of the objects to which the pointers refer. If the pointers do not refer to objects in the same array, the result is not defined.

A pointer can be compared to a constant expression that evaluates to 0. You can also compare a pointer to a pointer of type void*. The pointer is converted to a pointer of type void*.

If two pointers refer to the same object, they are considered equal. If two pointers refer to nonstatic members of the same object, the pointer to the object declared later is greater, provided that they are not separated by an access specifier; otherwise the comparison is undefined. If two pointers refer to data members of the same union, they have the same address value.

If two pointers refer to elements of the same array, or to the first element beyond the last element of an array, the pointer to the element with the higher subscript value is greater.

You can only compare members of the same object with relational operators.

Relational operators have left-to-right associativity. For example, the expression: a < b <= c

is interpreted as:

$$(a < b) <= c$$

If the value of a is less than the value of b, the first relationship yields 1 in C, or true in C++. The compiler then compares the value true (or 1) with the value of c (integral promotions are carried out if needed).

Equality == !=

The equality operators, like the relational operators, compare two operands for the validity of a relationship. The equality operators, however, have a lower precedence than the relational operators.

The type of the result is **int** and has the values 1 if the specified relationship os true, and 0 if false.

The type of the result is **bool** and has the values **true** or **false**.

The following table describes the two equality operators:

Operator	Usage
==	Indicates whether the value of the left operand is equal to the value of the right operand.
!=	Indicates whether the value of the left operand is not equal to the value of the right operand.

Both operands must have arithmetic or enumeration types or be pointers to the same type, or one operand must have a pointer type and the other operand must be a pointer to void or a null pointer. The result is type int in C or bool in C++.

If the operands have arithmetic types, the usual arithmetic conversions on the operands are performed.

If the operands are pointers, the result is determined by the locations of the objects to which the pointers refer.

If one operand is a pointer and the other operand is an integer having the value 0, the == expression is true only if the pointer operand evaluates to NULL. The != operator evaluates to true if the pointer operand does not evaluate to NULL.

You can also use the equality operators to compare pointers to members that are of the same type but do not belong to the same object. The following expressions contain examples of equality and relational operators:

```
time < max time == status < complete
letter != EOF
```

Note: The equality operator (==) should not be confused with the assignment (=) operator.

For example,

if (x == 3)evaluates to **true** (or 1) if x is equal to three. Equality tests like this should be coded with spaces between the operator and the operands to prevent unintentional assignments.

while

if (x = 3)is taken to be true because (x = 3) evaluates to a non-zero value (3). The expression also assigns the value 3 to x.

RELATED REFERENCES

• "Simple Assignment =" on page 107

Bitwise AND &

The & (bitwise AND) operator compares each bit of its first operand to the corresponding bit of the second operand. If both bits are 1's, the corresponding bit of the result is set to 1. Otherwise, it sets the corresponding result bit to 0.

Both operands must have an integral or enumeration type. The usual arithmetic conversions on each operand are performed. The result has the same type as the converted operands.

Because the bitwise AND operator has both associative and commutative properties, the compiler can rearrange the operands in an expression that contains more than one bitwise AND operator.

The following example shows the values of a, b, and the result of a & b represented as 16-bit binary numbers:

bit pattern of a 0000000001011100 bit pattern of b 0000000000101110 bit pattern of a & b 0000000000001100

Note: The bitwise AND (&) should not be confused with the logical AND. (&&) operator. For example,

```
1 & 4 evaluates to 0
while
   1 && 4 evaluates to true
```

RELATED REFERENCES

"Logical AND &&" on page 105

Bitwise Exclusive OR

The bitwise exclusive OR operator (in EBCDIC, the ^ symbol is represented by the ¬ symbol) compares each bit of its first operand to the corresponding bit of the second operand. If both bits are 1's or both bits are 0's, the corresponding bit of the result is set to 0. Otherwise, it sets the corresponding result bit to 1.

Both operands must have an integral or enumeration type. The usual arithmetic conversions on each operand are performed. The result has the same type as the converted operands and is not an Ivalue.

Because the bitwise exclusive OR operator has both associative and commutative properties, the compiler can rearrange the operands in an expression that contains more than one bitwise exclusive OR operator. Note that the ^ character can be represented by the trigraph ??'.

The following example shows the values of a, b, and the result of a b represented as 16-bit binary numbers:

```
0000000001011100
bit pattern of a
bit pattern of b
                              0000000000101110
bit pattern of a b
                              0000000001110010
```

RELATED REFERENCES

"Trigraph Sequences" on page 12

Bitwise Inclusive OR |

The (bitwise inclusive OR) operator compares the values (in binary format) of each operand and yields a value whose bit pattern shows which bits in either of the operands has the value 1. If both of the bits are 0, the result of that bit is 0; otherwise, the result is 1.

Both operands must have an integral or enumeration type. The usual arithmetic conversions on each operand are performed. The result has the same type as the converted operands and is not an Ivalue.

Because the bitwise inclusive OR operator has both associative and commutative properties, the compiler can rearrange the operands in an expression that contains more than one bitwise inclusive OR operator. Note that the | character can be represented by the trigraph ??!.

The following example shows the values of a, b, and the result of a | b represented as 16-bit binary numbers:

```
0000000001011100
bit pattern of a
bit pattern of b
                               0000000000101110
bit pattern of a | b
                               0000000001111110
```

Note: The bitwise OR (|) should not be confused with the logical OR (||) operator. For example,

```
1 | 4 evaluates to 5
1 | 4 evaluates to true
```

RELATED REFERENCES

- "Trigraph Sequences" on page 12
- "Logical OR | |"

Logical AND &&

The && (logical AND) operator indicates whether both operands are true.

If both operands have nonzero values, the result has the value 1. Otherwise, the result has the value 0. The type of the result is **int**. Both operands must have a arithmetic or pointer type. The usual arithmetic conversions on each operand are performed.

If both operands have values of **true**, the result has the value **true**. Otherwise, the result has the value false. Both operands are implicitly converted to **bool** and the result type is **bool**.

Unlike the & (bitwise AND) operator, the && operator guarantees left-to-right evaluation of the operands. If the left operand evaluates to 0 (or false), the right operand is not evaluated.

The following examples show how the expressions that contain the logical AND operator are evaluated:

Expression	Result	
1 && 0	false or 0	
1 && 4	true or 1	
0 && 0	false or 0	

The following example uses the logical AND operator to avoid division by zero: (y != 0) && (x / y)

The expression x / y is not evaluated when y != 0 evaluates to 0 (or false).

Note: The logical AND (&&) should not be confused with the bitwise AND (&) operator. For example:

```
1 && 4 evaluates to 1 (or true)
while
   1 & 4 evaluates to 0
```

RELATED REFERENCES

"Bitwise AND &" on page 103

Logical OR ||

The | (logical OR) operator indicates whether either operand is true.

If either of the operands has a nonzero value, the result has the value 1. Otherwise, the result has the value 0. The type of the result is **int**. Both operands must have a arithmetic or pointer type. The usual arithmetic conversions on each operand are performed.

If either operand has a value of **true**, the result has the value **true**. Otherwise, the result has the value false. Both operands are implicitly converted to bool and the result type is bool.

Unlike the | (bitwise inclusive OR) operator, the | | operator guarantees left-to-right evaluation of the operands. If the left operand has a nonzero (or true) value, the right operand is not evaluated.

The following examples show how expressions that contain the logical OR operator are evaluated:

Expression	Result	
1 0	true or 1	
1 4	true or 1	
0 0	false or 0	

The following example uses the logical OR operator to conditionally increment y: ++x || ++y;

The expression ++y is not evaluated when the expression ++x evaluates to a nonzero (or true) quantity.

Note: The logical OR (||) should not be confused with the bitwise OR (|) operator. For example:

```
1 | 4 evaluates to 1 (or true)
while
   1 | 4 evaluates to 5
```

RELATED REFERENCES

• "Bitwise Inclusive OR |" on page 104

C++ Pointer to Member Operators .* ->*

There are two pointer to member operators: .* and ->*.

The .* operator is used to dereference pointers to class members. The first operand must be of class type. If the type of the first operand is class type T, or is a class that has been derived from class type T, the second operand must be a pointer to a member of a class type T.

The ->* operator is also used to dereference pointers to class members. The first operand must be a pointer to a class type. If the type of the first operand is a pointer to class type T, or is a pointer to a class derived from class type T, the second operand must be a pointer to a member of class type T.

The .* and ->* operators bind the second operand to the first, resulting in an object or function of the type specified by the second operand.

If the result of .* or ->* is a function, you can only use the result as the operand for the () (function call) operator. If the second operand is an Ivalue, the result of .* or ->* is an lvalue.

RELATED REFERENCES

- "Class Member Lists" on page 209
- "Lvalues and Rvalues" on page 75
- "Pointers to Members" on page 214

Assignment Expressions

An assignment expression stores a value in the object designated by the left operand. There are two types of assignment operators: simple assignment and compound assignment.

The left operand in all assignment expressions must be a modifiable lyalue. The type of the expression is the type of the left operand. The value of the expression is the value of the left operand after the assignment has completed.

The result of an assignment expression is not an Ivalue.

All assignment operators have the same precedence and have right-to-left associativity.

Simple Assignment =

The simple assignment operator has the following form:

lvalue = expr

The operator stores the value of the right operand *expr* in the object designated by the left operand lvalue.

The left operand must be a modifiable lvalue. The type of an assignment operation is the type of the left operand.

If the left operand is not a class type, the right operand is implicitly converted to the type of the left operand. This converted type will not be qualified by **const** or volatile.

If the left operand is a class type, that type must be complete. The copy assignment operator of the left operand will be called.

If the left operand is an object of reference type, the compiler will assign the value of the right operand to the object denoted by the reference.

RELATED REFERENCES

- "Pointers" on page 58
- "volatile and const Qualifiers" on page 50
- "Pointers to Members" on page 214
- "References" on page 68
- "Structures" on page 36
- "Unions" on page 42
- "Equality == !=" on page 102
- "Initializing Pointers" on page 59

Compound Assignment

The compound assignment operators consist of a binary operator and the simple assignment operator. They perform the operation of the binary operator on both operands and give the result of that operation to the left operand.

The following table shows the operand types of compound assignment expressions:

Assignment Expressions

Operator	Left Operand	Right Operand
+= or -=	Arithmetic	Arithmetic
+= or -=	Pointer	Integral type
*=, /=, and %=	Arithmetic	Arithmetic
<=, >>=, &=, ^=, and =	Integral type	Integral type

Note that the expression

a *= b + c

is equivalent to

a = a * (b + c)

and not

a = a * b + c

The following table lists the compound assignment operators and shows an expression using each operator:

Operator	Example	Equivalent Expression
+=	index += 2	index = index + 2
_=	*(pointer++) -= 1	*pointer = *(pointer++) - 1
*=	bonus *= increase	bonus = bonus * increase
/=	time /= hours	time = time / hours
%=	allowance %= 1000	allowance = allowance % 1000
<<=	result <<= num	result = result << num
>>=	form >>= 1	form = form >> 1
&=	mask &= 2	mask = mask & 2
^ =	test ^= pre_test	test = test ^ pre_test
=	flag = ON	flag = flag ON

Although the equivalent expression column shows the left operands (from the example column) twice, it is in effect evaluated only once.

Cast Expressions

The cast operator is used for *explicit type conversions*. This operator has the following form, where *T* is a type, and *expr* is an expression:

(T) expr

It converts the value of *expr* to the type *T*. The result of this operation is an Ivalue if *T* is a reference. In all other cases, the result is an rvalue.

You can also use the following function-style notation:

expr(T)

This form also converts the value of *expr* to the type *T*. A function-style cast with no arguments, such as X() is equivalent to the declaration X t(), where t is a temporary object. Similarily, a function-style cast with more than one argument, such as X(a, b), is equivalent to the declaration X t(a, b).

For C++, the operand can have class type. If the operand has class type, it can be cast to any type for which the class has a user-defined conversion function. Casts can invoke a constructor, if the target type is a class, or they can invoke a conversion function, if the source type is a class. They can be ambiguous if both conditions hold.

An explicit type conversion can also be expressed by using the C++ type conversion operator **static_cast**.

The following demonstrates the use of the cast operator. The example dynamically creates an integer array of size 10:

```
#include <stdlib.h>
int main(void) {
   int* myArray = (int*) malloc(10 * sizeof(int));
   free(myArray);
   return 0;
}
```

The malloc() library function returns a **void** pointer that points to memory that will hold an object of the size of its argument. The statement int* myArray = (int*) malloc(10 * sizeof(int)) does the following

- Creates a void pointer that points to memory that can hold ten integers.
- Converts that **void** pointer into an integer pointer with the use of the cast operator.
- Assigns that integer pointer to myArray. Because a name of an array is the same as a pointer to the inital element of the array, myArray is an array of ten integers stored in the memory created by the call to malloc().

RELATED REFERENCES

- "Conversion Functions" on page 277
- "Conversion by Constructor" on page 276
- "Standard Type Conversions" on page 116
- "Lvalues and Rvalues" on page 75
- "References" on page 68
- "Temporary Objects" on page 273

C++ throw Expressions

A throw expression is used to throw exceptions to C++ exception handlers. A throw expression is of type **void**.

RELATED REFERENCES

- "Chapter 17. Exception Handling" on page 317
- "void Type" on page 35

Conditional Expressions

A *conditional expression* is a compound expression that contains a condition implicitly converted to bool ($operand_1$), an expression to be evaluated if the condition evaluates to true ($operand_2$), and an expression to be evaluated if the condition has the value false ($operand_3$).

Conditional expressions have right-to-left associativity. The left most operand is evaluated first, and then only one of the remaining two operands is evaluated.

Conditional Expressions

The conditional expression contains one two-part operator. The ? symbol follows the condition, and the: symbol appears between the two action expressions. All expressions that occur between the ? and : are treated as one expression.

The first operand must have a scalar type. The type of the second and third operands must be one of the following:

- An arithmetic type
- A compatible pointer, structure, or union type
- void

The second and third operands can also be a pointer or a null pointer constant.

Two objects are compatible when they have the same type but not necessarily the same type qualifiers (volatile orconst). Pointer objects are compatible if they have the same type or are pointers to void.

The first operand is evaluated, and its value determines whether the second or third operand is evaluated:

- If the value is true, the second operand is evaluated.
- If the value is false, the third operand is evaluated.

The result is the value of the second or third operand.

If the second and third expressions evaluate to arithmetic types, the usual arithmetic conversions are performed on the values. The types of the second and third operands determine the type of the result as shown in the following tables.

Type of Conditional C Expressions

Type of One Operand	Type of Other Operand	Type of Result
Arithmetic	Arithmetic	Arithmetic type after usual arithmetic conversions
Structure or union type	Compatible structure or union type	Structure or union type with all the qualifiers on both operands
void	void	void
Pointer to compatible type	Pointer to compatible type	Pointer to type with all the qualifiers specified for the type
Pointer to type	NULL pointer (the constant 0)	Pointer to type
Pointer to object or incomplete type	Pointer to void	Pointer to void with all the qualifiers specified for the type

Type of Conditional C++ Expressions

Type of One Operand	Type of Other Operand	Type of Result
Reference to type	Reference to type	Reference after usual reference conversions
Class T	Class T	Class T

Conditional Expressions

Type of One Operand	Type of Other Operand	Type of Result
Class T	Class X	Class type for which a conversion exists. If more than one possible conversion exists, the result is ambiguous.
throw expression	Other (type, pointer, reference)	Type of the expression that is not a throw expression

Examples of Conditional Expressions

The following expression determines which variable has the greater value, y or z, and assigns the greater value to the variable x:

```
x = (y > z) ? y : z;
```

The following is an equivalent statement:

```
if (y > z)
    x = y;
else
    x = z;
```

The following expression calls the function printf, which receives the value of the variable c, if c evaluates to a digit. Otherwise, printf receives the character constant 'x'.

```
printf(" c = %c\n", isdigit(c) ? c : 'x');
```

If the last operand of a conditional expression contains an assignment operator, use parentheses to ensure the expression evaluates properly. For example, the = operator has higher precedence than the ?: operator in the following expression:

```
int i,j,k;
(i == 7) ? j ++ : k = j;
```

The compiler will interpret this expression as if it were parenthesized this way:

```
int i,j,k;
((i == 7) ? j ++ : k) = j;
```

That is, k is treated as the third operand, not the entire assignment expression k = 0.1.

To assign the value of j to k i == 7 is false, enclose the last operand in parentheses:

```
int i,j,k;
(i == 7) ? j ++ : (k = j);
```

The typeid operator

The *typeid* operator provides program variable and expression type information. The operator has the following form:

```
typeid (expr)
```

The **typeid** operator returns an lvalue of type **const std::type_info** that represents the type of expression *expr*.

Conditional Expressions

You must include the standard template library header **<typeinfo>** to use the typeid operator.

If expr is a reference or a dereferenced pointer to a polymorphic class, typeid will return a type_info object that represents the object that the reference or pointer denotes at runtime. If it is not a polymorphic class, typeid will return a type_info object that represents the type of the reference or dereferenced pointer. The following example demonstrates this:

```
#include <iostream>
#include <typeinfo>
using namespace std;
struct A { virtual ~A() { } };
struct B : A { };
struct C { };
struct D : C { };
int main() {
  B bobj;
  A* ap = \&bobj;
  A& ar = bobj;
  cout << "ap: " << typeid(*ap).name() << endl;
cout << "ar: " << typeid(ar).name() << endl;</pre>
  D dobj;
  C* cp = &dobj;
  C& cr = dobj;
cout << "cp: " << typeid(*cp).name() << endl;</pre>
  cout << "cr: " << typeid(cr).name() << endl;</pre>
```

The following is the output of the above example:

ap: B ar: B cp: C cr: C

Classes A and B are polymorphic; classes C and D are not. Although cp and cr refer to an object of type D, typeid(*cp) and typeid(cr) return objects that represent class C.

Lvalue-to-rvalue, array-to-pointer, and function-to-pointer conversions will not be applied to expr. For example, the output of the following example will be int [10], not int *:

```
#include <iostream>
#include <typeinfo>
using namespace std;
int main() {
 int myArray[10];
 cout << typeid(myArray).name() << endl;</pre>
```

If *expr* is a class type, that class must be completely defined.

The **typeid** operator ignores top-level **const** or **volatile** qualifiers.

Comma Expression,

A *comma expression* contains two operands separated by a comma. Although the compiler evaluates both operands, the value of the right operand is the value of the expression. The left operand is evaluated, possibly producing side effects, and the value is discarded. The result of a comma expression is not an lyalue.

Both operands of a comma expression can have any type. All comma expressions have left-to-right associativity. The left operand is fully evaluated before the right operand.

In the following example, if omega has the value 11, the expression increments delta and assigns the value 3 to alpha:

```
alpha = (delta++, omega % 4);
```

Any number of expressions separated by commas can form a single expression. The compiler evaluates the left-most expression first. The value of the rightmost expression becomes the value of the entire expression.

For example, the value of the expression:

```
intensity++, shade * increment, rotate(direction);
```

is the value of the expression:

rotate(direction)

The primary use of the comma operator is to produce side effects in the following situations:

- Calling a function
- Entering or repeating an iteration loop
- Testing a condition
- Other situations where a side effect is required but the result of the expression is not immediately needed

To use the comma operator in a context where the comma has other meanings, such as in a list of function arguments or a list of initializers, you must enclose the comma operator in parentheses. For example, the function

$$f(a, (t = 3, t + 2), c);$$

has only three arguments: the value of a, the value 5, and the value of c. The value of the second argument is the result of the comma expression in parentheses:

$$t = 3, t + 2$$

which has the value 5.

The following table gives some examples of the uses of the comma operator:

Statement	Effects
for (i=0; i<2; ++i, f());	A for statement in which i is incremented and f() is called at each iteration.

Comma Expression

Statement	Effects
if (f(), ++i, i>1) { /* */ }	An if statement in which function f() is called, variable i is incremented, and variable i is tested against a value. The first two expressions within this comma expression are evaluated before the expression i>1. Regardless of the results of the first two expressions, the third is evaluated and its result determines whether the if statement is processed.
func((++a, f(a)));	A function call to func() in which a is incremented, the resulting value is passed to a function f(), and the return value of f() is passed to func(). The function func() is passed only a single argument, because the comma expression is enclosed in parentheses within the function argument list.

Chapter 6. Implicit Type Conversions

An expression e of a given type is *implicitly converted* if used in one of the following situations:

- Expression e is used as an operand of an arithmetic or logical operation.
- Expression e is used as a condition in an **if** statement or an iteration statement (such as a **for** loop). Expression e will be converted to **bool** (or **int** in C).
- Expression e is used in a **switch** statement. Expression e will be converted to an integral type.
- Expression e is used in an initialization. This includes the following:
 - An assignment is made to an Ivalue that has a different type than e.
 - A function is provided an argument value of e that has a different type than the parameter.
 - Expression e is specified in the return statement of a function, and e has a different type from the defined return type for the function.

The compiler will allow an implicit conversion of an expression e to a type T if and only if the compiler would allow the following statement:

```
T var = e;
```

For example when you add values having different data types, both values are first converted to the same type: when a **short int** value and an **int** value are added together, the **short int** value is converted to the **int** type.

You can perform explicit type conversions using one of the cast operators, the function style cast, or the C style cast.

RELATED REFERENCES

- "Chapter 5. Expressions and Operators" on page 71
- "static_cast Operator" on page 82
- "reinterpret_cast Operator" on page 83
- "const_cast Operator" on page 84
- "dynamic_cast Operator" on page 85
- "Cast Expressions" on page 108

Integral and Floating-Point Promotions

An *integral promotion* is the conversion of one integral type to another where the second type can hold all possible values of the first type. Certain fundamental types can be used wherever an integer can be used. The following fundamental types can be converted through integral promotion are:

- char
- ▶ C++ bool
- wchar t
- · short int
- enumerators
- objects of enumeration type
- integer bit fields (both signed and unsigned)

Except for wchar_t, if the value cannot be represented by an int, the value is converted to an unsigned int. For wchar_t, if an int can represent all the values of

Integral Promotions

the original type, the value is converted to the type that can best represent all the values of the original type. For example, if a long can represent all the values, the value is converted to a long.

Floating-Point Promotions You can convert an rvalue of type float to an rvalue of type double. The value of the expression is unchanged. This conversion is a floating-point promotion.

RELATED REFERENCES

- "char and wchar_t Type Specifiers" on page 32
- "Boolean Variables" on page 33
- "Integer Variables" on page 34
- "Enumerations" on page 46
- "Declaring and Using Bit Fields in Structures" on page 39

Standard Type Conversions

Many C and C++ operators cause *implicit type conversions*, which change the type of an expression. When you add values having different data types, both values are first converted to the same type. For example, when a **short int** value and an int value are added together, the short int value is converted to the int type. It can result in loss of data if the value of the original object is outside the range representable by the shorter type.

Implicit type conversions can occur when:

- An operand is prepared for an arithmetic or logical operation.
- An assignment is made to an Ivalue that has a different type than the assigned
- A function is provided an argument value that has a different type than the parameter.
- The value specified in the **return** statement of a function has a different type from the defined return type for the function.

You can perform explicit type conversions using the C-style cast, the C++ function-style cast, or one of the C++ cast operators.

```
#include <iostream>
using namespace std;
int main() {
 float num = 98.76;
 int x1 = (int) num;
 int x2 = int(num);
 int x3 = static cast<int>(num);
 cout << "x1 = " << x1 << endl;
 cout << "x2 = " << x2 << endl;
 cout << "x3 = " << x3 << endl;
```

The following is the output of the above example:

```
x1 = 98
x2 = 98
x3 = 98
```

The integer x1 is assigned a value in which num has heen explicitly converted to an int with the C-style cast. The integer x2 is assigned a value that has been converted with the function-style cast. The integer x3 is assigned a value that has been converted with the static_cast operator.

RELATED REFERENCES

- "return Statement" on page 158
- "Cast Expressions" on page 108

Lvalue-to-Rvalue Conversions

If an Ivalue appears in a situation in which the compiler expects an rvalue, the compiler converts the lvalue to an rvalue.

An Ivalue e of a type T can be converted to an rvalue if T is not a function or array type. The type of e after conversion will be T. The following table lists exceptions to this:

Situation before conversion	Resulting behavior
T is an incomplete type	compile-time error
e refers to an uninitialized object	undefined behavior
e refers to an object not of type T, nor a type derived from T	undefined behavior
T is a non-class type	the type of e after conversion is T, not qualified by either const or volatile

RELATED REFERENCES

• "volatile and const Qualifiers" on page 50

Boolean Conversions

You can convert integral, floating-point, arithmetic, enumeration, pointer, and pointer to member rvalue types to an rvalue of type bool. Any value other than a zero, null pointer, or null member pointer value is converted to true; A zero, null pointer, or null member pointer value is converted to false.

The following is an example of boolean conversions:

```
void f(int* a, int b)
 bool d = a; // false if a == NULL
 bool e = b; // false if b == 0
```

The variable d will be false if a is equal to a null pointer. Otherwise, d will be true. The variable e will be **false** if b is equal to zero. Otherwise,e will be **true**.

RELATED REFERENCES

• "Boolean Variables" on page 33

Integral Conversions

You can convert the following:

- An rvalue of integer type (including signed and unsigned integer types) to another rvalue of integer type
- An rvalue of enumeration type to an rvalue of integer type

If you are converting a integer a to an unsigned type, the resulting value x is the least unsigned integer such that a and x are congruent modulo 2^n , where n is the number of bits used to represent an unsigned type. If two numbers a and x are

congruent modulo 2^n , the following expression is true, where the function pow(m, n) returns the value of m to the power of n:

```
a \% pow(2, n) == x \% pow(2, n)
```

If you are converting an integer a to a signed type, the compiler does not change the resulting value if the new type is large enough to hold the a. If the new type is not large enough, the behavior is defined by the compiler.

If you are converting a bool to an integer, values of false are converted to 0; values of **true** are converted to 1.

Integer promotions belong to a different category of conversions; they are not integeral conversions.

RELATED REFERENCES

• "Integer Variables" on page 34

Floating-Point Conversions

You can convert an rvalue of floating-point type to another rvalue of floating-point type.

Floating-point promotions (converting from float to double) belong to a different category of conversions; they are not floating-point conversions.

RELATED REFERENCES

- "Floating-Point Variables" on page 33
- "Integral and Floating-Point Promotions" on page 115

Pointer Conversions

Pointer conversions are performed when pointers are used, including pointer assignment, initialization, and comparison.

Conversion to void*

Any pointer to an object of a type T, optionally qualified with const, volatile, or const volatile, can be converted to void*, keeping the same const or volatile qualifications. You can also convert any pointer to a function to a void*, provided that a void* has sufficient bits to hold it.

Derived-to-Base Conversions

You can convert an rvalue pointer of type B* to an rvalue pointer of class A* where A is an accessible base class of B as long as the conversion is not ambiguous. The conversion is ambiguous if the expression for the accessible base class can refer to more than one distinct class. The resulting value points to the base class subobject of the derived class object. If the pointer of type B* is null, it will be converted to a null pointer of type A*. Note that you cannot convert a pointer to a class into a pointer to its base class if the base class is a virtual base class of the derived class.

Null Pointer Constants

A constant expression that evaluates to zero is a *null pointer constant*. This expression can be converted to a pointer. This pointer will be a null pointer (pointer with a zero value), and is guaranteed not to point to any object.

Array-to-Pointer Conversions

You can convert an Ivalue or rvalue with type "array of N," where N is the type of a single element of the array, to N*. The result is a pointer to the initial element of the array. You cannot perform the conversion if the expression is used as the operand of the & (address) operator or the **sizeof** operator.

Function-to-Pointer Conversions

You can convert an Ivalue that is a function of type T to an rvalue that is a pointer to a function of type T, except when the expression is used as the operand of the & (address) operator, the () (function call) operator, or the **sizeof** operator.

RELATED REFERENCES

- "void Type" on page 35
- "Pointers" on page 58
- "Integer Constant Expressions" on page 75
- "Arrays" on page 62
- "Pointers to Functions" on page 141

Reference Conversions

A reference conversion can be performed wherever a reference initialization occurs, including reference initialization done in argument passing and function return values. A reference to a class can be converted to a reference to an accessible base class of that class as long as the conversion is not ambiguous. The result of the conversion is a reference to the base class subobject of the derived class object.

Reference conversion is allowed if the corresponding pointer conversion is allowed.

RELATED REFERENCES

- "References" on page 68
- "Initializing References" on page 69
- "Calling Functions and Passing Arguments" on page 134
- "Function Return Values" on page 140

Pointer-to-Member Conversions

Pointer-to-member conversion can occur when pointers to members are initialized, assigned, or compared.

A constant expression that evaluates to zero can be converted to the null pointer to a member.

Note: A pointer to a member is not the same as a pointer to an object or a pointer to a function.

A pointer to a member of a base class can be converted to a pointer to a member of a derived class if the following conditions are true:

- The conversion is not ambiguous. The conversion is ambiguous if multiple instances of the base class are in the derived class.
- A pointer to the derived class can be converted to a pointer to the base class. If this is the case, the base class is said to be accessible.
- Member types must match. For example suppose class A is a base class of class B. You cannot convert a pointer to member of A of type int to a pointer to member of type B of type float.

The base class cannot be virtual.

RELATED REFERENCES

- "Integer Constant Expressions" on page 75
- "Access Control of Base Class Members" on page 237
- "Pointers to Members" on page 214
- "C++ Pointer to Member Operators .* ->*" on page 106

Qualification Conversions

You can convert an rvalue of type cv1 T* where cv1 is any combination of zero or more const or volatile qualifications, to an rvalue of type cv2 T* if cv2 T* is more const or volatile qualified than cv1 T*.

You can convert an rvalue of type pointer to member of a class X of cv1 T, to an rvalue of type pointer to member of a class X of cv2 T if cv2 T is more const or **volatile** qualified than *cv1* T.

RELATED REFERENCES

• "volatile and const Qualifiers" on page 50

Function Argument Conversions



If a function declaration is present and includes declared argument types, the compiler performs type checking. If no function declaration is visible when a function is called, the compiler can perform default argument promotions, which consist of the following:

- Integral promotions
- Arguments with type float are converted to type double.

SEC++ Function declarations in C++ must always specify their parameter types. Also, functions may not be called if it has not already been declared.

RELATED REFERENCES

- "Integral and Floating-Point Promotions" on page 115
- "Function Declarations" on page 124

Other Conversions

By definition, the void type has no value. Therefore, it cannot be converted to any other type, and no other value can be converted to void by assignment. However, a value can be explicitly cast to void.

No conversions between structure or union types are allowed.

There are no standard conversions between class types.

You can write your own conversion operators for class types.

In C, when you define a value using the **enum** type specifier, the value is treated as an int. Conversions to and from an enum value proceed as for the int type.

You can convert from an **enum** to any integral type but not from an integral type to an enum.

RELATED REFERENCES

- "void Type" on page 35
- "User-Defined Conversions" on page 274
- "Enumerations" on page 46

Arithmetic Conversions

The conversions depend on the specific operator and the type of the operand or operands. However, many operators perform similar conversions on operands of integer and floating-point types. These standard conversions are known as the arithmetic conversions because they apply to the types of values ordinarily used in arithmetic.

Arithmetic conversions are used for matching operands of arithmetic operators.

Arithmetic conversion proceeds in the following order:

Operand Type	Conversion
One operand has long long type.	The other operand is converted to long long .
One operand has long double type	The other operand is converted to long double .
One operand has double type	The other operand is converted to double .
One operand has float type	The other operand is converted to float .
One operand has unsigned long int type	The other operand is converted to unsigned long int .
One operand has unsigned int type and the other operand has long int type and the value of the unsigned int can be represented in a long int	The operand with unsigned int type is converted to long int .
One operand has unsigned int type and the other operand has long int type and the value of the unsigned int cannot be represented in a long int	Both operands are converted to unsigned long int
One operand has long int type	The other operand is converted to long int .
One operand has unsigned int type	The other operand is converted to unsigned int .
Both operands have int type	The result is type int .

RELATED REFERENCES

- "Chapter 5. Expressions and Operators" on page 71
- "Integer Variables" on page 34
- "Floating-Point Variables" on page 33

The explicit Keyword

The explicit keyword controls unwanted implicit type conversions. It can only be used in declarations of constructors within a class declaration.

Arithmetic Conversions

A constructor declared with only one argument and without the explicit keyword is a converting constructor. You can construct objects with a converting constructor using the assignment operator. Declaring a constructor of this type with the explicit keyword prevents this behavior. For example, except for the default constructor, the constructors in the following class are converting constructors.

```
class A
  public:
  A();
  A(int);
  A(const char*, int = 0);
```

The following declarations are legal.

```
A c = 1;
A d = "Venditti";
```

The first declaration is equivalent to A c = A(1).

If you declare the constructor of the class with the explicit keyword the previous declarations would be illegal.

For example, if you declare the class as:

```
class A
{ public:
   explicit A();
   explicit A(int);
   explicit A(const char*, int = 0);
};
```

You can only assign values that match the values of the class type.

For example, the following statements will be legal:

```
A a1;
A a2 = A(1);
A a3(1);
A a4 = A("Venditti");
A* p = new A(1);
A a5 = (A)1;
A a6 = static cast<A>(1);
```

Chapter 7. Functions

A function *declaration* consists of a return type, a name, and an argument list. It is used to declare the format and existence of a function prior to its use.

A function *definition* contains a function declaration and the body of the function. A function can only have one definition.

The declaration is used by the compiler for argument type checking and argument conversions. Declarations can appear several times in a program, provided the declarations are compatible. They allow the compiler to check for mismatches between the parameters of a function call and those in the function declaration.

Declarations are typically placed in header files, while function definitions appear in source files.

RELATED REFERENCES

- "Function Declarations" on page 124
- "Function Definitions" on page 128

C++ Enhancements to C Functions

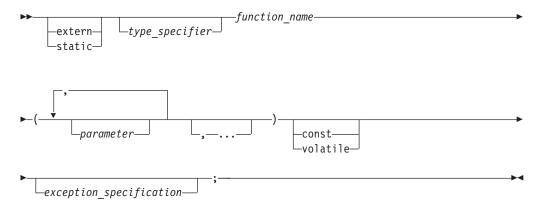
The C++ language provides many enhancements to C functions. These are:

- Reference arguments
- Default arguments
- Reference return types
- Inline functions
- · Member functions
- Overloaded functions
- Operator functions
- · Constructor and destructor functions
- Conversion functions
- Virtual functions
- Function templates
- · Exception specifications
- Constructor initializers

RELATED REFERENCES

- "Passing Arguments by Reference" on page 136
- "Default Arguments in C++ Functions" on page 137
- "Using References as Return Types" on page 141
- "Inline Functions" on page 141
- "Member Functions" on page 211
- "Overloading Functions" on page 185
- "Overloading Operators" on page 187
- "Constructors and Destructors Overview" on page 257
- "Conversion Functions" on page 277
- "Virtual Functions" on page 249

A function declaration establishes the name and the number and types of the parameters of the function.



A function argument is an expression that you use within the parenthesis of a function call. A function parameter is an object or reference declared within the parenthesis of a function declaration or definition. When you call a function, the arguments are evaluated, and each parameter is initialized with the value of the corresponding argument. The semantics of argument passing are identical to those of assignment.

Implicit declaration of functions is not allowed.

- The default return type of a function is **int**.
- There is no default return type.

To indicate that the function does not return a value, declare it with a return type of void.

Only member functions may have **const** or **volatile** specifiers after the parenthesized parameter list.

A function cannot be declared as returning a data object having a volatile or const type but it can return a pointer to a volatile or const object. A function may return a pointer to function, or a pointer to the first element of an array, but it may not return a value that has a type of array or function.

The exception_specification limits the function from throwing only a specified list of exceptions.

Some declarations do not name the parameters within the parameter lists; the declarations simply specify the types of parameters and the return values. This is called *prototyping*. The following example demonstrates this:

int func(int,long);

The ellipsis (...) may be the only argument in C++. In this case, the comma is not required. In C, you cannot have the ellipsis as the only argument.

Types cannot be defined in return or argument types. For example, the C++ compiler will allow the following declaration of print():

```
struct X { int i; };
void print(X x);
```

Similarly, the C compiler will allow the following declaration:

```
struct X { int i; };
void print(struct X x);
```

The C and C++ compilers will not allow the following declaration of the same function:

```
void print(struct X { int i; } x);
                                             //error
```

This example attempts to declare a function print() that takes an object x of class X as its argument. However, the class definition is not allowed within the argument list.

In another example, the C++ compiler will allow the following declaration of counter():

```
enum count {one, two, three};
count counter();
```

Similarly the C compiler will allow the following declaration:

```
enum count {one, two, three};
enum count counter();
```

The C and C++ compilers will not allow the following declaration of the same

```
enum count{one, two, three} counter();
                                             //error
```

In the attempt to declare counter(), the enumeration type definition cannot appear in the return type of the function declaration.

RELATED REFERENCES

- "volatile and const Qualifiers" on page 50
- "Exception Specifications" on page 329

C++ Function Declarations

C++ In C++, you can specify the qualifiers **volatile** and **const** in member function declarations. You can also specify exception specifications in function declarations. All C++ functions must be declared before they can be called.

VAC++ The unordered incremental compiler does not require functions to be declared before they are called.

RELATED REFERENCES

- "volatile and const Qualifiers" on page 50
- "const and volatile Member Functions" on page 212
- "Exception Specifications" on page 329

Multiple Function Declarations

C++ All function declarations for one particular function must have the same number and type of parameters, and must have the same return type.

These return and parameter types are part of the function type, although the default arguments and exception specifications are not.

If a previous declaration of an object or function is visible in an enclosing scope, the identifier has the same linkage as the first declaration. For example, a variable or function that is first declared with the keyword static and later declared with the keyword extern has internal linkage. However, a variable or function that has no linkage and later declared with a linkage specifier will have the linkage you have specified.

For the purposes of argument matching, ellipsis and linkage keywords are considered a part of the function type. They must be used consistently in all declarations of a function. If the only difference between the parameter types in two declarations is in the use of typedef names or unspecified argument array bounds, the declarations are the same. A const or volatile specifier is also part of the function type, but can only be part of a declaration or definition of a nonstatic member function.

You may overload function names. An overloaded function declaration is a declaration that had been declared with the same name as a previously declared declaration in the same scope, except that both declarations have different types. If you call an overloaded function name, the compiler determines the most appropriate definition to use by comparing the argument types you used to call the function or operator with the parameter types specified in the definitions.

Declaring two functions differing only in return type is not valid function overloading, and is flagged as an error. For example:

```
void f();
              // error, two definitions differ only in
int f();
              // return type
int q()
   return f();
```

RELATED REFERENCES

"Overloading Functions" on page 185

Argument Names in Function Declarations

You can supply parameter names in a function declaration, but the compiler ignores them except in the following two situations:

- 1. If two parameter names have the same name within a single declaration. This is an error.
- 2. If a parameter name is the same as a name outside the function. In this case the name outside the function is hidden and cannot be used in the parameter declaration.

In the following example, the third parameter name intersects is meant to have enumeration type subway line, but this name is hidden by the name of the first parameter. The declaration of the function subway() causes a compile-time error because subway_line is not a valid type name because the first parameter name subway_line hides the namespace scope enum type and cannot be used again in the second parameter.

```
enum subway line {yonge, university, spadina, bloor};
int subway (char * subway line, int stations,
                  subway line intersects);
```

RELATED REFERENCES

"Function Declarations" on page 124

Examples of Function Declarations

The following code fragments show several function declarations. The first declares a function f that takes two integer arguments and has a return type of void:

```
void f(int, int);
```

The following code fragment declares a pointer p1 to a function that takes a pointer to a constant character and returns an integer:

```
int (*p1) (const char*);
```

The following code fragment declares a function f1 that takes an integer argument, and returns a pointer to a function that takes an integer argument and returns an integer:

```
int (*f1(int))(int);
```

Alternatively, a **typedef** can be used for the complicated return type of function f1:

```
typedef int f1 return type(int);
fl_return_type* fl(int);
```

The following declaration is of an external function f2 that takes a constant integer as its first argument, can have a variable number and variable types of other arguments, and returns type int.

```
int extern f2(const int ...);
```

In C, a comma is required before the ellipsis:

```
int extern f2(const int, ...);
```

Function f3 has a return type int, and takes a int argument with a default value that is the value returned from function f2:

```
const int j = 5;
int f3( int x = f2(j));
```

EXAMPLE 2 Function f6 is a **const** class member function of class X, takes no arguments, and has a return type of int:

```
class X
public:
      int f6() const;
```

C++ Function f4 takes no arguments, has return type **void**, and can throw class objects of types X and Y.

```
class X;
class Y;
// ...
void f4() throw(X,Y);
```

Function f5 takes no arguments, has return type void, and will call unexpected() if it throws an exception of any type.

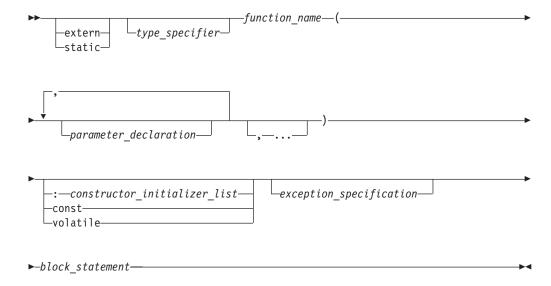
```
void f5() throw();
```

RELATED REFERENCES

- "Default Arguments in C++ Functions" on page 137
- "const and volatile Member Functions" on page 212
- "Exception Specifications" on page 329
- "extern Storage Class Specifier" on page 25
- "Chapter 4. Declarators" on page 55

Function Definitions

A function definition contains a function declaration and the body of a function.



A function definition contains the following:

• An optional storage class specifier extern or static, which determines the scope of the function. If a storage class specifier is not given, the function has external linkage.

A type specifier, which determines the type of value that the function returns.

In C, the type specifier is optional. If a type specifier is not given, the function has type int.

The type specifier is not optional in C++.

A function declarator, which is the function name followed by a parenthesized list of parameter types and names. It can further describe the type of the value that the function returns, and lists the type and name of each parameter that the function expects. In the following function dfinition, f(int a, int b) is the function declarator:

```
int f(int a, int b) {
 return a + b;
```

Optional **const** or **volatile** specifiers after the function declarator. Only member functions may have these.

- An optional exception specification, which limits the function from throwing only a specified list of exceptions.
 - A block statement, which contains data definitions and code.
 - **Solution** You can also have a function try block instead of a block statement. If the function definition is a constructor, you can have a constructor initializer list before the block statement. In the following class definition, x(0), y('c') is a constructor initializer list:

```
class A {
  int x;
  char y;
public:
  A() : x(0), y('c') \{ \}
```

A function can be called by itself or by other functions. By default, function definitions have external linkage, and can be called by functions defined in other files. A storage class specifier of static means that the function name has global scope only, and can be directly invoked only from within the same translation unit.

- This use of static is deprecated in C++. Instead, place the function in the unnamed namespace.
- In C only, if a function definition has external linkage and a return type of int, calls to the function can be made before it is visible because an implicit declaration of extern int func(); is assumed.

All declarations for a given function must be compatible; that is, the return type is the same and the parameters have the same type. Note that overloaded functions have the same name.

The default type for the return value and parameters of a function is int, and the default storage class specifier is extern. If the function does not return a value, use the keyword void as the type specifier. You can use the keyword void as a parameter declaration to indicate the function is not passed any arguments.

A function cannot return a function, array, or object with a volatile or const type, but it can return a pointer to these or any other types.

- In C, you cannot declare a function as a struct or union member.
- In C, a function cannot return any type having the volatile or const qualifier.

You cannot define an array of functions. You can, however, define an array of pointers to functions.

The following example is a definition of the function sum:

```
int sum(int x, int y)
   return(x + y);
```

Function Definitions

The function sum has external linkage, returns an object that has type int, and has two parameters of type int declared as x and y. The function body contains a single statement that returns the sum of x and y.

In the following example, ary is an array of two function pointers. Type casting is performed to the values assigned to any for compatibility:

```
#include <stdio.h>
typedef void (*ARYTYPE)();
int func1(void);
void func2(double a);
int main(void)
 double num = 333.3333;
 int retnum;
 ARYTYPE ary[2];
 ary[0] = (ARYTYPE) func1;
 ary[1] = (ARYTYPE) func2;
 retnum=((int (*)())ary[0])(); /* calls func1 */
 printf("number returned = %i\n", retnum);
  ((void (*)(double))ary[1])(num); /* calls func2 */
 return(0);
int func1(void)
 int number=3:
 return number;
void func2(double a)
 printf("result of func2 = %f\n", a);
```

The following is the output of the above example:

```
number returned = 3
result of func2 = 333.333300
```

RELATED REFERENCES

- "extern Storage Class Specifier" on page 25
- "static Storage Class Specifier" on page 28
- "Block Statement" on page 145
- "Pointers" on page 58
- "References" on page 68
- "Structures" on page 36
- "Unions" on page 42
- "volatile and const Qualifiers" on page 50

Ellipsis and void

An ellipsis at the end of the parameter specifications is used to specify that a function has a variable number of parameters. The number of parameters is equal to, or greater than, the number of parameter specifications. At least one parameter declaration must come before the ellipsis.

```
int f(int, ...);
```

The comma before the ellipsis is optional. In addition, a parameter declaration is not required before the ellipsis.

The comma before the ellipsis as well as a parameter declaration before the ellipsis are both required in C.

Parameter promotions are performed as needed, but no type checking is done on the variable arguments.

You can declare a function with no arguments in two ways:

```
int f(void);
int f();
```

C++ An empty argument declaration list or the argument declaration list of (void) indicates a function that takes no arguments.

An empty argument declaration list means that the function may take any number or type of parameters.

The type **void** cannot be used as an argument type, although types derived from **void** (such as pointers to **void**) can be used.

In the following example, the function f() takes one integer argument and returns no value, while g() expects no arguments and returns an integer.

```
void f(int);
int g(void);
```

RELATED REFERENCES

• "void Type" on page 35

Examples of Function Definitions

The following example contains a function declarator i_sort with table declared as a pointer to int and length declared as type int. Note that arrays as parameters are implicitly converted to a pointer to the element type.

```
** This example illustrates function definitions.
** Note that arrays as parameters are implicitly
** converted to a pointer to the type.
#include <stdio.h>
void i sort(int table[], int length);
int main(void)
  int table[]=\{1,5,8,4\};
  int length=4;
  printf("length is %d\n",length);
  i sort(table,length);
void i_sort(int table[], int length)
 int i, j, temp;
 for (i = 0; i < length -1; i++)
```

Function Definitions

```
for (j = i + 1; j < length; j++)
      if (table[i] > table[j])
        temp = table[i];
        table[i] = table[j];
        table[j] = temp;
}
```

The following are examples of function declarations (also called function prototypes):

```
double square(float x);
int area(int x,int y);
static char *search(char);
```

The following example illustrates how a **typedef** identifier can be used in a function declarator:

```
typedef struct tm fmt { int minutes;
                      int hours;
                      char am pm;
                    } struct t;
long time seconds(struct t arrival)
```

The following function set_date declares a pointer to a structure of type date as a parameter. date_ptr has the storage class specifier register.

```
void set date(register struct date *date ptr)
 date ptr->mon = 12;
 date ptr->day = 25;
 date ptr->year = 87;
```

RELATED REFERENCES

- "Block Statement" on page 145
- "Function Definitions" on page 128
- "Function Declarations" on page 124

The main() Function

When a program begins running, the system calls the function main, which marks the entry point of the program. Every program must have one function named main. No other function in the program can be called main. A main function has one of two forms:

- **int main (void)** block statement
- **C++ int main ()** block_statement
- int main (int argc, char ** argv) block_statement

The argument argc is the number of command-line arguments passed to the program. The argument *argv* is a pointer to an array of strings, where argv [0] is the name you used to run your program from the command-line, argv[1] the first argument that you passed to your program, argv[2] the second argument, and so on.

By default, main has the storage class extern.

C++ You cannot declare **main** as **inline** or **static**. You cannot call **main** from within a program or take the address of main. You cannot overload this function.

RELATED REFERENCES

- "extern Storage Class Specifier" on page 25
- "Inline Functions" on page 141
- "static Storage Class Specifier" on page 28

Arguments to main

The function **main** can be declared with or without parameters.

```
int main(int argc, char *argv[])
```

Although any name can be given to these parameters, they are usually referred to as *argc* and *argv*.

The first parameter, argc (argument count), has type int and indicates how many arguments were entered on the command line.

The second parameter, argv (argument vector), has type array of pointers to char array objects. char array objects are null-terminated strings.

The value of argc indicates the number of pointers in the array argv. If a program name is available, the first element in argv points to a character array that contains the program name or the invocation name of the program that is being run. If the name cannot be determined, the first element in argv points to a null character.

This name is counted as one of the arguments to the function main. For example, if only the program name is entered on the command line, argc has a value of 1 and argv[0] points to the program name.

Regardless of the number of arguments entered on the command line, argv[argc] always contains NULL.

RELATED REFERENCES

- "Integer Variables" on page 34
- "char and wchar_t Type Specifiers" on page 32

Example of Arguments to main

The following program backward prints the arguments entered on a command line such that the last argument is printed first:

```
#include <stdio.h>
int main(int argc, char *argv[])
 while (--argc > 0)
   printf("%s ", argv[argc]);
```

Invoking this program from a command line with the following:

```
backward string1 string2
```

```
gives the following output:
  string2 string1
```

The arguments *argc* and *argv* would contain the following values:

Object	Value	
argc	3	
argv[0]	pointer to string "backward"	
argv[1]	pointer to string "string1"	
argv[2]	pointer to string "string2"	
argv[3]	NULL	

Note: Be careful when entering mixed case characters on a command line because some environments are not case sensitive. Also, the exact format of the string pointed to by argv[0] is system dependent.

RELATED REFERENCES

- · "Calling Functions and Passing Arguments"
- "Type Specifiers" on page 30
- "Identifiers" on page 14
- "Block Statement" on page 145

Calling Functions and Passing Arguments

The arguments of a function call are used to initialize the parameters of the function definition.

Integral and floating-point promotions will first be done to the values of the arguments before the function is called.

The type of an argument is checked against the type of the corresponding parameter in the function declaration. All standard and user-defined type conversions are applied as necessary.

```
For example:
#include <stdio.h>
#include <math.h>
/* Declaration */
extern double root(double, double);
/* Definition */
double root(double value, double base) {
 double temp = exp(log(value)/base);
 return temp;
int main(void) {
 int value = 144;
 int base = 2;
 printf("The root is: f\n", root(value, base));
 return 0;
```

The output is The root is: 12.000000

In the above example, because the function root is expecting arguments of type double, the two int arguments value and base are implicitly converted to type **double** when the function is called.

Calling Functions and Passing Arguments

The order in which arguments are evaluated and passed to the function is implementation-defined. For example, the following sequence of statements calls the function tester:

```
int x;
x = 1;
tester(x++, x);
```

The call to tester in the example may produce different results on different compilers. Depending on the implementation, x++ may be evaluated first or x may be evaluated first. To avoid the ambiguity and have x++ evaluated first, replace the preceding sequence of statements with the following:

```
int x, y;
x = 1;
y = \chi + + ;
tester(y, x);
```

EC++ In C++, if a nonstatic class member function is passed as an argument, the argument is converted to a pointer to member.

Solution If a class has a destructor or a copy constructor that does more than a bitwise copy, passing a class object by value results in the construction of a temporary that is actually passed by reference.

It is an error when a function argument is a class object and all of the following properties hold:

- The class needs a copy constructor.
- The class does not have a user-defined copy constructor.
- A copy constructor cannot be generated for that class.

RELATED REFERENCES

- "Function Calls ()" on page 78
- "Integral and Floating-Point Promotions" on page 115
- "Constructors" on page 259

Passing Arguments by Value

If you call a function with an argument that corresponds to a non-reference parameter, you have passed that argument by value. The parameter is initialized with the value of the argument. You can change the value of the parameter (if that parameter has not been declared **const**) within the scope of the function, but these changes will not affect the value of the argument alue of the argument in the calling function.

The following are examples of passing arguments by value:

The following statement calls the function **printf**, which receives a character string and the return value of the function sum, which receives the values of a and b: printf("sum = $%d\n$ ", sum(a,b));

The following program passes the value of count to the function increment, which increases the value of the parameter x by 1.

```
** An example of passing an argument to a function
**/
#include <stdio.h>
```

Calling Functions and Passing Arguments

```
void increment(int);
int main(void)
 int count = 5;
 /* value of count is passed to the function */
 increment(count);
 printf("count = %d\n", count);
 return(0);
void increment(int x)
 ++χ;
 printf("x = %d\n", x);
```

The output illustrates that the value of count in main remains unchanged:

```
count = 5
```

RELATED REFERENCES

• "Function Calls ()" on page 78

Passing Arguments by Reference

Passing by reference refers to a method of passing arguments where the value of an argument in the calling function can be modified in the called function.

To pass an argument by reference, you declare the corresponding parameter with a reference type.

The following example shows how arguments are passed by reference. Note that reference parameters are initialized with the actual arguments when the function is called.

```
#include <stdio.h>
void swapnum(int &i, int &j) {
 int temp = i;
 i = j;
 j = temp;
int main(void) {
 int a = 10;
 int b = 20;
 swapnum(a, b);
 printf("A is %d and B is %d\n", a, b);
 return 0;
```

When the function swapnum() is called, the actual values of the variables a and b are exchanged because they are passed by reference. The output is:

```
A is 20 and B is 10
```

You must define the parameters of swapnum() as references if you want the values of the actual arguments to be modified by the function swapnum().

Calling Functions and Passing Arguments

EC++ In order to modify a reference that is **const**-qualified, you must cast away its constness with the const_cast operator. The following example demonstrates this:

```
#include <iostream>
using namespace std;
void f(const int& x) {
 int* y = const cast<int>(&x);
  (*y)++;
int main() {
 int a = 5;
 f(a);
  cout << a << endl;</pre>
```

This example outputs 6.

You can modify the values of nonconstant objects through pointer parameters. The following example demonstrates this:

```
#include <stdio.h>
int main(void)
 void increment(int *x);
 int count = 5;
 /* address of count is passed to the function */
 increment(&count);
 printf("count = %d\n", count);
  return(0);
void increment(int *x)
 ++*X;
 printf("*x = %d\n", *x);
```

The following is the output of the above code:

```
*x = 6
count = 6
```

The example passes the address of count to increment(). Function increment() increments count through the pointer parameter x.

RELATED REFERENCES

- "References" on page 68
- "const_cast Operator" on page 84

Default Arguments in C++ Functions

You can provide default values for function parameters.

```
For example:
#include <iostream>
```

using namespace std;

Default Arguments in C++ Functions

```
int a = 1:
int f(int a) { return a; }
int g(int x = f(a)) { return x; }
int h() {
 a = 2;
   int a = 3;
   return g();
int main() {
 cout << h() << endl;
```

This example prints 2 to standard output, because the a referred to in the declaration of g() is the one at file scope, which has the value 2 when g() is called.

The default argument must be implicitly convertible to the parameter type.

A pointer to a function must have the same type as the function. Attempts to take the address of a function by reference without specifying the type of the function will produce an error. The type of a function is not affected by arguments with default values.

The following example shows that default arguments are not considered part of a function's type. The default argument allows you to call a function without specifying all of the arguments, it does not allow you to create a pointer to the function that does not specify the types of all the arguments. Function f can be called without an explicit argument, but the pointer badpointer cannot be defined without specifying the type of the argument:

```
int f(int = 0);
void g()
                           // ok
  int a = f(1);
  int b = f();
                              // ok, default argument used
                              // ok, type of f() specified (int)
int (*pointer)(int) = &f;
int (*badpointer)() = &f;
                               // error, badpointer and f have
                               // different types. badpointer must
                               // be initialized with a pointer to
                               // a function taking no arguments.
```

RELATED REFERENCES

• "Pointers to Functions" on page 141

Restrictions on Default Arguments

Of the operators, only the function call operator and the operator **new** can have default arguments when they are overloaded.

Parameters with default arguments must be the trailing parameters in the function declaration parameter list. For example:

```
void f(int a, int b = 2, int c = 3); // trailing defaults
void g(int a = 1, int b = 2, int c); // error, leading defaults
void h(int a, int b = 3, int c);
                                  // error, default in middle
```

Once a default argument has been given in a declaration or definition, you cannot redefine that argument, even to the same value. However, you can add default

Default Arguments in C++ Functions

arguments not given in previous declarations. For example, the last declaration below attempts to redefine the default values for a and b:

```
void f(int a, int b, int c=1);
void f(int a, int b=1, int c);  // valid, add another default
void f(int a=1, int b, int c);  // valid, add another default
void f(int a=1, int b=1, int c=1);  // error, redefined defaults
```

You can supply any default argument values in the function declaration or in the definition. Any parameters in the parameter list following a default argument value must have a default argument value specified in this or a previous declaration of the function.

You cannot use local variables in default argument expressions. For example, the compiler generates errors for both function q() and function h() below:

```
void f(int a)
     int b=4;
     void g(int c=a); // Local variable "a" cannot be used here
     void h(int d=b); // Local variable "b" cannot be used here
```

RELATED REFERENCES

- "Function Calls ()" on page 78
- "C++ new Operator" on page 92
- "Default Arguments in C++ Functions" on page 137

Evaluating Default Arguments

When a function defined with default arguments is called with trailing arguments missing, the default expressions are evaluated. For example:

```
void f(int a, int b = 2, int c = 3); // declaration
// ...
int a = 1;
                 // same as call f(a,2,3)
f(a);
f(a,10); // same as call f(a,10,3)
f(a,10,20); // no default arguments
```

Default arguments are checked against the function declaration and evaluated when the function is called. The order of evaluation of default arguments is undefined. Default argument expressions cannot use other parameters of the function. For example:

```
int f(int q = 3, int r = q); // error
```

The argument r cannot be initialized with the value of the argument q because the value of q may not be known when it is assigned to r. If the above function declaration is rewritten:

```
int f(int q = 3, int r = q); // error
```

The value of r in the function declaration still produces an error because the variable q defined outside of the function is hidden by the argument q declared for the function. Similarly:

```
typedef double D;
int f(int D, int z = D(5.3)); // error
```

Here the type D is interpreted within the function declaration as the name of an integer. The type D is hidden by the argument D. The cast D(5.3) is therefore not interpreted as a cast because D is the name of the argument not a type.

Default Arguments in C++ Functions

In the following example, the nonstatic member a cannot be used as an initializer because a does not exist until an object of class X is constructed. You can use the static member b as an initializer because b is created independently of any objects of class X. You can declare the member b after its use as a default argument because the default values are not analyzed until after the final bracket } of the class declaration.

```
class X
   int a;
    f(int z = a); // error
   g(int z = b); // valid
   static int b;
};
```

RELATED REFERENCES

"Default Arguments in C++ Functions" on page 137

Function Return Values

You must return a value from a function unless the function has a return type of void.

All functions must return a value, including those that have a return type of void. The return value is specified in a return statement. The following code fragment shows a function definition, including the **return** statement:

```
int add(int i, int j)
 return i + j; // return statement
```

The function add() can be called as shown in the following code fragment:

```
int a = 10,
   b = 20;
int answer = add(a, b); // answer is 30
```

In this example, the return statement initializes a variable of the returned type. The variable answer is initialized with the **int** value 30. The type of the returned expression is checked against the returned type. All standard and user-defined conversions are performed as necessary.

The following return statements show different ways of returning values to a caller:

```
// Returns no value
return;
return result;
                        // Returns the value of result
                         // Returns the value 1
return 1;
return (x * x);
                         // Returns the value of x * x
```

Each time a function is called, new copies of its variables with automatic storage are created. Because the storage for these automatic variables may be reused after the function has terminated, a pointer or reference to an automatic variable should not be returned.

If a class object is returned, a temporary object may be created if the class has copy constructors or a destructor.

RELATED REFERENCES

- "auto Storage Class Specifier" on page 25
- "Temporary Objects" on page 273
- "Destructors" on page 266

Using References as Return Types

References can also be used as return types for functions. The reference returns the lvalue of the object to which it refers. This allows you to place function calls on the left side of assignment statements.

Referenced return values are used when assignment operators and subscripting operators are overloaded so that the results of the overloaded operators can be used as actual values.

Note: Returning a reference to an automatic variable gives unpredictable results.

RELATED REFERENCES

- "Overloading Assignments" on page 190
- "Overloading Subscripting" on page 193
- "auto Storage Class Specifier" on page 25

Pointers to Functions

A pointer to a function points to the address of the function's executable code. You can use pointers to call functions and to pass functions as arguments to other functions. You cannot perform pointer arithmetic on pointers to functions.

The type of a pointer to a function is based on both the return type and parameter types of the function.

A declaration of a pointer to a function must have the pointer name in parentheses. The function call operator () has a higher precedence than the dereference operator *. Without them, the compiler interprets the statement as a function that returns a pointer to a specified return type. For example:

```
int *f(int a);
                     // function f returning an int*
int (*g)(int a);
                     // pointer g to a function returning an int
```

In the first declaration, f is interpreted as a function that takes an **int** as argument, and returns a pointer to an int. In the second declaration, g is interpreted as a pointer to a function that takes an int argument and that returns an int.

RELATED REFERENCES

- "Pointers" on page 58
- "Pointer Conversions" on page 118
- "extern Storage Class Specifier" on page 25

Inline Functions

C++ A function is declared inline by using the inline function specifier or by defining a member function within a class or structure definition.

The **inline** specifier is a suggestion to the compiler that an inline expansion can be performed. Instead of transferring control to and from the function code segment, a modified copy of the function body may be substituted directly for the function call.

Inline Functions

An inline function can be declared and defined simultaneously. If it is declared with the keyword inline, it can be declared without a definition. The following code fragment shows an inline function definition. Note that the definition includes both the declaration and body of the inline function.

```
inline int add(int i, int j) { return i + j; }
```

Both member and nonmember functions can be inline, and both have external linkage by default.

The use of the **inline** specifier does not change the meaning of the function. The inline expansion of a function may not preserve the order of evaluation of the actual arguments.

RELATED REFERENCES

• "Member Functions" on page 211

Chapter 8. Statements

A statement is the smallest independent computational unit. The following is a summary of the statements available in C and C++:

- labeled statements
 - identifier labels
 - case labels
 - default labels
- expression statements
- · blocks or comound statements
- selection statements
 - if statements
 - switch statements
- iteration statements
 - while statements
 - **do** statements
 - for statements
- jump statements
 - break statements
 - continuestatements
 - return statements
 - goto statements
- · declaration statements
- try blocks

Labels

There are three kinds of labels: identifier, case, and default.

Identifier label statements have the following form:



The label consists of the *identifier* and the colon (:) character. An identifier label is only used as the the target of a **goto** statement. Identifier labels have their own namespace; you do not have to worry about identifier labels conflicting with other identifiers.

Case statements have the following form:

▶►—case—constant expression—:—statement—

Default label statements have the following form:

▶—default—:—statement—

Case and default label statements only appear in switch statements.

Examples of Labels

Labels

```
/* null statement label */
```

RELATED REFERENCES

- "goto Statement" on page 159
- "switch Statement" on page 148

Expression Statements

An expression statement contains an expression. The expression can be null.

An expression statement has the form:



An expression statement evaluates expression, then discards the value of the expression. An expression statement without an expression is a null statement.

Examples of Expressions

```
printf("Account Number: \n");
marks = dollars * exch_rate;
                                          /* call to the printf
                                          /* assignment to marks
                                                                       */
(difference < 0) ? ++losses : ++gain; /* conditional increment */
```

RELATED REFERENCES

"Chapter 5. Expressions and Operators" on page 71

Resolving Ambiguous Statements in C++

The C++ syntax does not disambiguate between expression statements and declaration statements. The ambiguity arises when an expression statement has a function-style cast as its leftmost subexpression. (Note that, because C does not support function-style casts, this ambiguity does not occur in C programs.) If the statement can be interpreted both as a declaration and as an expression, the statement is interpreted as a declaration statement.

Note: The ambiguity is resolved only on a syntactic level. The disambiguation does not use the meaning of the names, except to assess whether or not they are type names.

The following expressions disambiguate into expression statements because the ambiguous subexpression is followed by an assignment or an operator. type spec in the expressions can be any type specifier:

```
type spec(i)++;
                          // expression statement
                         // expression statement
type_spec(i,3)<<d;</pre>
type_spec(i)->1=24;
                        // expression statement
```

In the following examples, the ambiguity cannot be resolved syntactically, and the statements are interpreted as declarations. type spec is any type specifier:

```
type spec(*i)(int);
                          // declaration
                         // declaration
type spec(j)[5];
type_spec(m) = { 1, 2 };  // declaration
type_spec(*k) (float(3)); // declaration
```

The last statement above causes a compile-time error because you cannot initialize a pointer with a float value.

Any ambiguous statement that is not resolved by the above rules is by default a declaration statement. All of the following are declaration statements:

```
type_spec(a);
                         // declaration
type_spec(*b)();
                         // declaration
                        // declaration
type_spec(c)=23;
type_spec(d),e,f,g=0;
                        // declaration
type spec(h)(e,3);
                        // declaration
```

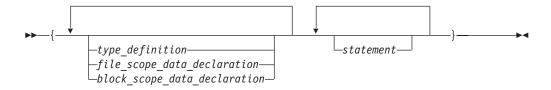
RELATED REFERENCES

- "Chapter 3. Declarations" on page 23
- "Chapter 5. Expressions and Operators" on page 71
- "Function Calls ()" on page 78

Block Statement

A block statement, or compound statement, lets you group any number of data definitions, declarations, and statements into one statement. All definitions, declarations, and statements enclosed within a single set of braces are treated as a single statement. You can use a block wherever a single statement is allowed.

A block statement has the form:



In C, Any definitions and declarations must come before the statements.

A block defines a local scope. If a data object is usable within a block and its identifier is not redefined, all nested blocks can use that data object.

Example of Blocks

The following program shows how the values of data objects change in nested blocks:

```
** This example shows how data objects change in nested blocks.
#include <stdio.h>
 int main(void)
   int x = 1;
                                   /* Initialize x to 1 */
   int y = 3;
   if (y > 0)
       int x = 2;
                                   /* Initialize x to 2 */
      printf("second x = %4d n", x);
```

Block Statement

```
printf("first x = %4d n", x);
  return(0);
}
```

The program produces the following output:

```
first x =
```

Two variables named x are defined in main. The first definition of x retains storage while main is running. However, because the second definition of x occurs within a nested block, printf("second x = $4d\n$, x); recognizes x as the variable defined on the previous line. Because printf("first $x = %4d\n$ ", x); is not part of the nested block, x is recognized as the first definition of x.

RELATED REFERENCES

- "Storage Class Specifiers" on page 24
- "Type Specifiers" on page 30

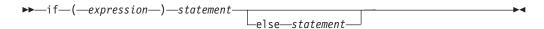
if Statement

An if statement lets you conditionally process a statement when the specified test expression, implicitly converted to bool, evaluates to true. If the implicit conversion to bool fails the program is ill-formed.

In C, an if statement lets you conditionally process a statement when the specified test expression evaluates to a nonzero value. The test expression must be of arithmetic or pointer type.

You can optionally specify an **else** clause on the **if** statement. If the test expression evaluates to false (or in C, a zero value) and an else clause exists, the statement associated with the else clause runs. If the test expression evaluates to true, the statement following the expression runs and the **else** clause is ignored.

An **if** statement has the form:



When if statements are nested and else clauses are present, a given else is associated with the closest preceding if statement within the same block.

A single statement following any selection statements (if, switch) is treated as a compound statement containing the original statement. As a result any variables declared on that statement will be out of scope after the if statement. For example:

```
if (x)
int i;
is equivalent to:
if(x)
{ int i; }
```

Variable i is visible only within the if statement. The same rule applies to the else part of the if statement.

Examples of if Statements

The following example causes grade to receive the value A if the value of score is greater than or equal to 90.

```
if (score >= 90)
  grade = 'A';
```

The following example displays Number is positive if the value of number is greater than or equal to 0. If the value of number is less than 0, it displays Number is negative.

```
if (number >= 0)
  printf("Number is positive\n");
  printf("Number is negative\n");
```

The following example shows a nested if statement:

```
if (paygrade == 7)
   if (level >= 0 && level <= 8)
      salary *= 1.05;
   else
      salary *= 1.04;
   salary *= 1.06;
cout << "salary is " << salary << endl;</pre>
```

The following example shows a nested if statement that does not have an else clause. Because an else clause always associates with the closest if statement, braces might be needed to force a particular else clause to associate with the correct if statement. In this example, omitting the braces would cause the else clause to associate with the nested if statement.

```
if (kegs > 0) {
   if (furlongs > kegs)
      fpk = furlongs/kegs;
else
  fpk = 0;
```

The following example shows an **if** statement nested within an **else** clause. This example tests multiple conditions. The tests are made in order of their appearance. If one test evaluates to a nonzero value, a statement runs and the entire if statement ends.

```
if (value > 0)
  ++increase;
else if (value == 0)
  ++break_even;
  ++decrease;
```

RELATED REFERENCES

"Boolean Variables" on page 33

switch Statement

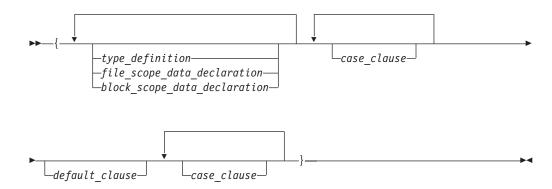
A switch statement lets you transfer control to different statements within the switch body depending on the value of the switch expression. The switch expression must evaluate to an integral or enumeration value. The body of the switch statement contains case clauses that consist of

- A case label
- An optional **default** label
- A case expression
- · A list of statements.

If the value of the switch expression equals the value of one of the case expressions, the statements following that case expression are processed. If not, the default label statements, if any, are processed.

A **switch** statement has the form:

The switch body is enclosed in braces and can contain definitions, declarations, case clauses, and a default clause. Each case clause and default clause can contain statements.

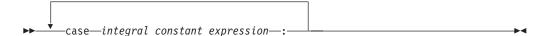


Note: An initializer within a *type_definition*, *file_scope_data_declaration* or block_scope_data_declaration is ignored.

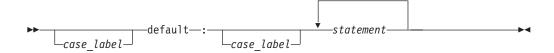
A case clause contains a case label followed by any number of statements. A case clause has the form:



A case label contains the word case followed by an integral constant expression and a colon. The value of each integral constant expression must represent a different value; you cannot have duplicate case labels. Anywhere you can put one case label, you can put multiple case labels. A case label has the form:



A default clause contains a default label followed by one or more statements. You can put a case label on either side of the default label. A switch statement can have only one **default** label. A *default_clause* has the form:



The switch statement passes control to the statement following one of the labels or to the statement following the switch body. The value of the expression that precedes the switch body determines which statement receives control. This expression is called the *switch expression*.

The value of the **switch** expression is compared with the value of the expression in each case label. If a matching value is found, control is passed to the statement following the case label that contains the matching value. If there is no matching value but there is a default label in the switch body, control passes to the default labelled statement. If no matching value is found, and there is no **default** label anywhere in the **switch** body, no part of the **switch** body is processed.

When control passes to a statement in the switch body, control only leaves the switch body when a break statement is encountered or the last statement in the switch body is processed.

If necessary, an integral promotion is performed on the controlling expression, and all expressions in the case statements are converted to the same type as the controlling expression. The switch expression can also be of class type if there is a single conversion to integral or enumeration type.

You can put data definitions at the beginning of the switch body, but the compiler does not initialize auto and register variables at the beginning of a switch body. You can have declarations in the body of the **switch** statement.

You cannot use a **switch** statement to jump over initializations.

In C++, you cannot transfer control over a declaration containing an explicit or implicit initializer unless the declaration is located in an inner block that is completely bypassed by the transfer of control. All declarations within the body of a switch statement that contain initializers must be contained in an inner block.

Examples of switch Statements

The following switch statement contains several case clauses and one default clause. Each clause contains a function call and a break statement. The break statements prevent control from passing down through each statement in the **switch** body.

switch Statement

If the **switch** expression evaluated to '/', the switch statement would call the function divide. Control would then pass to the statement following the switch body.

```
char key;
printf("Enter an arithmetic operator\n");
scanf("%c",&key);
switch (key)
  case '+':
     add();
      break;
  case '-':
      subtract();
      break;
  case '*':
     multiply();
     break;
  case '/':
     divide();
     break;
   default:
      printf("invalid key\n");
      break;
```

If the switch expression matches a case expression, the statements following the case expression are processed until a break statement is encountered or the end of the switch body is reached. In the following example, break statements are not present. If the value of text[i] is equal to 'A', all three counters are incremented. If the value of text[i] is equal to 'a', lettera and total are increased. Only total is increased if text[i] is not equal to 'A' or 'a'.

```
char text[100]:
int capa, lettera, total;
// ...
for (i=0; i<sizeof(text); i++) {</pre>
    switch (text[i])
       case 'A':
         capa++;
       case 'a':
         lettera++;
       default:
         total++;
}
```

The following switch statement performs the same statements for more than one case label:

```
/**
** This example contains a switch statement that performs
** the same statement for more than one case label.
#include <stdio.h>
```

```
int main(void)
 int month;
 /* Read in a month value */
 printf("Enter month: ");
 scanf("%d", &month);
  /* Tell what season it falls into */
 switch (month)
    case 12:
    case 1:
    case 2:
        printf("month %d is a winter month\n", month);
       break;
    case 3:
    case 4:
    case 5:
       printf("month %d is a spring month\n", month);
       break:
    case 6:
    case 7:
    case 8:
       printf("month %d is a summer month\n", month);
        break;
    case 9:
    case 10:
    case 11:
        printf("month %d is a fall month\n", month);
    case 66:
    case 99:
     default:
        printf("month %d is not a valid month\n", month);
 }
  return(0);
```

If the expression month has the value 3, control passes to the statement: printf("month %d is a spring month\n", month);

The break statement passes control to the statement following the switch body.

RELATED REFERENCES

"break Statement" on page 155

while Statement

A while statement repeatedly runs the body of a loop until the controlling expression evaluates to false (or 0 in C).

A **while** statement has the form:

```
▶►—while—(—expression—)—statement—
```

while Statement

The expression is evaluated to determine whether or not to process the body of the loop.

The expression must be convertible to **bool**.

The *expression* must be of arithmetic or pointer type.

If the expression evaluates to false, the body of the loop never runs. If the expression does not evaluate to false, the loop body is processed. After the body has run, control passes back to the expression. Further processing depends on the value of the condition.

A break, return, or goto statement can cause a while statement to end, even when the condition does not evaluate to false.

Example of while Statements

In the following program, item[index] triples and is printed out, as long as the value of the expression ++index is less than MAX INDEX. When ++index evaluates to MAX INDEX, the **while** statement ends.

```
** This example illustrates the while statement.
**/
#define MAX INDEX (sizeof(item) / sizeof(item[0]))
#include <stdio.h>
int main(void)
   static int item[] = { 12, 55, 62, 85, 102 };
  int index = 0;
  while (index < MAX_INDEX)
     item[index] *= 3;
     printf("item[%d] = %d\n", index, item[index]);
      ++index;
   }
   return(0);
```

RELATED REFERENCES

- "break Statement" on page 155
- "do Statement"
- "for Statement" on page 153
- "continue Statement" on page 156
- "return Statement" on page 158
- "goto Statement" on page 159
- "Boolean Variables" on page 33

do Statement

A do statement repeatedly runs a statement until the test expression evaluates to false (or θ in C). Because of the order of processing, the statement is run at least once.

A do statement has the form:

- The controlling *expression* must convertible to type **bool**.
- The *expression* must be of arithmetic or pointer type.

The body of the loop is run before the controlling **while** clause is evaluated. Further processing of the do statement depends on the value of the while clause. If the while clause does not evaluate to false, the statement runs again. When the while clause evaluates to false, the statement ends.

A break, return, or goto statement can cause the processing of a do statement to end, even when the while clause does not evaluate to false.

Example of do Statements

The following example keeps incrementing i while i is less than 5:

#include <stdio.h>

```
int main(void) {
 int i = 0;
 do {
   i++;
   printf("Value of i: %d\n", i);
 while (i < 5);
  return 0;
```

The following is the output of the above example:

Value of i: 1 Value of i: 2 Value of i: 3 Value of i: 4 Value of i: 5

RELATED REFERENCES

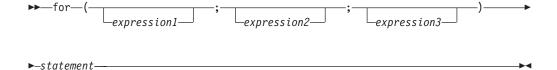
- "break Statement" on page 155
- "return Statement" on page 158
- "goto Statement" on page 159
- "Boolean Variables" on page 33

for Statement

A for statement lets you do the following:

- Evaluate an expression before the first iteration of the statement (initialization)
- · Specify an expression to determine whether or not the statement should be processed (the *condition*)
- Evaluate an expression after each iteration of the statement (often used to increment for each iteration)
- Repeatedly process the statement if the controlling part does not evaluate to false (or 0 in C).

A **for** statement has the form:



Expression1

Is the *initialization expression*. It is evaluated only before the statement is processed for the first time. You can use this expression to initialize a variable. If you do not want to evaluate an expression prior to the first iteration of the statement, you can omit this expression.

Expression2

Is the conditional expression. It is evaluated before each iteration of the statement.

It must evaluate to an arithmetic or pointer type.

If it evaluates to **false** (or 0 in C), the statement is not processed and control moves to the next statement following the for statement. If expression2 does not evaluate to false, the statement is processed. If you omit expression2, it is as if the expression had been replaced by true, and the for statement is not terminated by failure of this condition.

Expression3

Is evaluated after each iteration of the *statement*. This expression is often used for incrementing, decrementing, or assigning to a variable. This expression is optional.

A break, return, or goto statement can cause a for statement to end, even when the second expression does not evaluate to false. If you omit expression2, you must use a break, return, or goto statement to end the for statement.

Solution In C++ programs, you can also use *expression1* to declare a variable as well as initialize it. If you declare a variable in this expression, or anywhere else in statement, that variable goes out of scope at the end of the for loop.

YAC: You can set a compile option where a variable declared in the scope of a **for** statement is not local to the **for** statement.

Examples of for Statements

The following for statement prints the value of count 20 times. The for statement initially sets the value of count to 1. After each iteration of the statement, count is incremented.

```
int count;
for (count = 1; count <= 20; count++)
   printf("count = %d\n", count);
```

The following sequence of statements accomplishes the same task. Note the use of the **while** statement instead of the **for** statement.

```
int count = 1;
while (count <= 20)
   printf("count = %d\n", count);
   count++;
```

The following for statement does not contain an initialization expression:

```
for (; index > 10; --index)
  list[index] = var1 + var2;
  printf("list[%d] = %d\n", index, list[index]);
```

The following for statement will continue running until scanf receives the letter e:

```
{
  scanf("%c", &letter);
   if (letter == '\n')
      continue;
   if (letter == 'e')
      break;
  printf("You entered the letter %c\n", letter);
```

The following for statement contains multiple initializations and increments. The comma operator makes this construction possible. The first comma in the for expression is a punctuator for a declaration. It declares and initializes two integers, i and j. The second comma, a comma operator, allows both i and j to be incremented at each step through the loop.

```
for (int i = 0, j = 50; i < 10; ++i, j += 50)
   cout << "i = " << i << "and j = " << j << endl;
```

The following example shows a nested for statement. It prints the values of an array having the dimensions [5][3].

```
for (row = 0; row < 5; row++)
   for (column = 0; column < 3; column++)</pre>
      printf("%d\n", table[row][column]);
```

The outer statement is processed as long as the value of row is less than 5. Each time the outer for statement is executed, the inner for statement sets the initial value of column to zero and the statement of the inner for statement is executed 3 times. The inner statement is executed as long as the value of column is less than 3.

RELATED REFERENCES

- "break Statement"
- "return Statement" on page 158
- "goto Statement" on page 159
- "Boolean Variables" on page 33

break Statement

A break statement lets you end an iterative (do, for, or while) statement or a switch statement and exit from it at any point other than the logical end. A break may only appear on one of these statements.

A break statement has the form:

```
▶▶—break—:-
```

break Statement

In an iterative statement the **break** statement ends the loop and moves control to the next statement outside the loop. Within nested statements, the break statement ends only the smallest enclosing do, for, switch, or while statement.

In a switch statement, the break passes control out of the switch body to the next statement outside the switch statement.

RELATED REFERENCES

- "do Statement" on page 152
- "for Statement" on page 153
- "while Statement" on page 151
- "switch Statement" on page 148

continue Statement

A continue statement ends the current iteration of a loop. Program control is passed from the **continue** statement to the end of the loop body.

A **continue** statement has the form:

```
▶►—continue—;-
```

A **continue** statement can only appear within the body of an iterative statement.

The **continue** statement ends the processing of the action part of an iterative (do, for, or while) statement and moves control to the loop continuation portion of the statement. For example, if the iterative statement is a for statement, control moves to the third expression in the condition part of the statement, then to the second expression (the test) in the condition part of the statement.

Within nested statements, the **continue** statement ends only the current iteration of the do, for, or while statement immediately enclosing it.

Examples of continue Statements

The following example shows a **continue** statement in a **for** statement. The continue statement causes processing to skip over those elements of the array rates that have values less than or equal to 1.

```
** This example shows a continue statement in a for statement.
**/
#include <stdio.h>
#define SIZE 5
int main(void)
  static float rates[SIZE] = \{1.45, 0.05, 1.88, 2.00, 0.75\};
  printf("Rates over 1.00\n");
   for (i = 0; i < SIZE; i++)
      if (rates[i] <= 1.00) /* skip rates <= 1.00 */
        continue;
     printf("rate = %.2f\n", rates[i]);
```

```
}
  return(0);
The program produces the following output:
```

```
Rates over 1.00
rate = 1.45
rate = 1.88
rate = 2.00
```

The following example shows a **continue** statement in a nested loop. When the inner loop encounters a number in the array strings, that iteration of the loop ends. Processing continues with the third expression of the inner loop. The inner loop ends when the ' $\0$ ' escape sequence is encountered.

```
/**
** This program counts the characters in strings that are part
** of an array of pointers to characters. The count excludes
** the digits 0 through 9.
**/
#include <stdio.h>
#define SIZE 3
int main(void)
  static char *strings[SIZE] = { "ab", "c5d", "e5" };
  int i;
  int letter_count = 0;
  char *pointer;
  for (i = 0; i < SIZE; i++)
                                             /* for each string
                                             /* for each each character */
     for (pointer = strings[i]; *pointer != '\0'; ++pointer)
                                  /* if a number
         if (*pointer >= '0' && *pointer <= '9')
           continue;
        letter_count++;
  printf("letter count = %d\n", letter count);
  return(0);
```

The program produces the following output:

```
letter count = 5
```

RELATED REFERENCES

- "do Statement" on page 152
- "for Statement" on page 153
- "while Statement" on page 151

return Statement

A return statement ends the processing of the current function and returns control to the caller of the function.

A **return** statement has the form:



A return statement in a function is optional. The compiler issues a warning if a return statement is not found in a function declared with a return type. If the end of a function is reached without encountering a return statement, control is passed to the caller as if a **return** statement without an expression were encountered. A function can contain multiple return statements.

RELATED REFERENCES

• "Chapter 7. Functions" on page 123

Value of a return Expression and Function Value

If an expression is present on a return statement, the value of the expression is returned to the caller. If the data type of the expression is different from the function return type, conversion of the return value takes place as if the value of the expression were assigned to an object with the same function return type.

If an expression is not present on a return statement, the value of the return statement is undefined. If an expression is not given on a return statement in a function declared with a nonvoid return type, an error message is issued, and the result of calling the function is unpredictable. For example:

```
int func1()
return;
int func2()
return (4321);
void main() {
int a=func1(); // result is unpredictable!
int b=func2();
```

You cannot use a return statement with an expression when the function is declared as returning type void.

Examples of return Statements

```
/* Returns no value
return;
return result; /* Returns the value of result */
return 1; /* Returns the value 1 */
return (x * x); /* Returns the value of x * x */
```

The following function searches through an array of integers to determine if a match exists for the variable number. If a match exists, the function match returns the value of i. If a match does not exist, the function match returns the value -1 (negative one).

```
int match(int number, int array[], int n)
   int i;
  for (i = 0; i < n; i++)
     if (number == array[i])
        return (i);
  return(-1);
```

goto Statement

A goto statement causes your program to unconditionally transfer control to the statement associated with the label specified on the goto statement.

A **goto** statement has the form:

```
▶►—goto—label_identifier—;-
```

Because the **goto** statement can interfere with the normal sequence of processing, it makes a program more difficult to read and maintain. Often, a break statement, a continue statement, or a function call can eliminate the need for a goto statement.

You cannot use a **goto** statement to jump over initializations.

If an active block is exited using a **goto** statement, any local variables are destroyed when control is transferred from that block.

Example of goto Statements

The following example shows a **goto** statement that is used to jump out of a nested loop. This function could be written without using a **goto** statement.

```
** This example shows a goto statement that is used to
** jump out of a nested loop.
#include <stdio.h>
void display(int matrix[3][3]);
int main(void)
   int matrix[3][3]=\{1,2,3,4,5,2,8,9,10\};
  display(matrix);
  return(0);
void display(int matrix[3][3])
  int i, j;
  for (i = 0; i < 3; i++)
     for (j = 0; j < 3; j++)
         if ( (matrix[i][j] < 1) | (matrix[i][j] > 6) )
            goto out_of_bounds;
         printf("matrix[%d][%d] = %d\n", i, j, matrix[i][j]);
```

goto Statement

```
return;
out of bounds: printf("number must be 1 through 6\n");
```

RELATED REFERENCES

- "break Statement" on page 155
- "continue Statement" on page 156
- "Labels" on page 143

Null Statement

The *null statement* performs no operation. It has the form:



A null statement can hold the label of a labeled statement or complete the syntax of an iterative statement.

Examples of Null Statements

The following example initializes the elements of the array price. Because the initializations occur within the for expressions, a statement is only needed to finish the for syntax; no operations are required.

```
for (i = 0; i < 3; price[i++] = 0)
```

A null statement can be used when a label is needed before the end of a block statement. For example:

```
void func(void) {
 if (error_detected)
   goto depart;
  /* further processing */
 depart: ; /* null statement required */
```

Chapter 9. Preprocessor Directives

Preprocessor Overview

Preprocessing is a preliminary operation on C and C++ files before they are passed to the compiler. It allows you to do the following:

- Replace tokens in the current file with specified replacement tokens
- Imbed files within the current file
- Conditionally compile sections of the current file
- Generate diagnostic messages
- Change the line number of the next line of source and change the file name of the current file
- Apply machine-specific rules to specified sections of code

A *token* is a series of characters delimited by white space. The only white space allowed on a preprocessor directive is the space, horizontal tab, vertical tab, form feed, and comments. The new-line character can also separate preprocessor tokens.

The preprocessed source program file must be a valid C or C++ program.

The preprocessor is controlled by the following directives:

#define Defines a macro.

#undef Removes a preprocessor macro definition.

#error Defines text for a compile-time error message.

#include Inserts text from another source file.

#if Conditionally suppresses portions of source code, depending on

the result of a constant expression.

#ifdef Conditionally includes source text if a macro name is defined.

#ifndef Conditionally includes source text if a macro name is not defined.

#else Conditionally includes source text if the previous #if, #ifdef,

#ifndef, or #elif test fails.

#elif Conditionally includes source text if the previous #if, #ifdef,

#ifndef, or **#elif** test fails, depending on the value of a constant

expression.

#endif Ends conditional text.

#line Supplies a line number for compiler messages.

#pragma Specifies implementation-defined instructions to the compiler.

RELATED REFERENCES

- "Tokens" on page 9
- "Preprocessor Directive Format" on page 162

Preprocessor Directive Format

Preprocessor directives begin with the # token followed by a preprocessor keyword. The # token must appear as the first character that is not white space on a line. The # is not part of the directive name and can be separated from the name with white spaces.

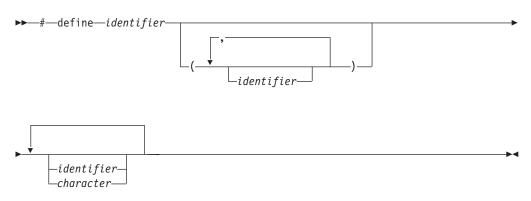
A preprocessor directive ends at the new-line character unless the last character of the line is the \ (backslash) character. If the \ character appears as the last character in the preprocessor line, the preprocessor interprets the \ and the new-line character as a continuation marker. The preprocessor deletes the \ (and the following new-line character) and splices the physical source lines into continuous logical lines.

Except for some #pragma directives, preprocessor directives can appear anywhere in a program.

Macro Definition and Expansion (#define)

A preprocessor define directive directs the preprocessor to replace all subsequent occurrences of a macro with specified replacement tokens.

A preprocessor **#define** directive has the form:



The #define directive can contain an object-like definition or a function-like definition.

Object-Like Macros

An object-like macro definition replaces a single identifier with the specified replacement tokens. The following object-like definition causes the preprocessor to replace all subsequent instances of the identifier COUNT with the constant 1000:

#define COUNT 1000

If the statement int arry[COUNT];

appears after this definition and in the same file as the definition, the preprocessor would change the statement to int arry[1000];

in the output of the preprocessor.

Other definitions can make reference to the identifier COUNT:

```
#define MAX COUNT COUNT + 100
```

The preprocessor replaces each subsequent occurrence of MAX COUNT with COUNT + 100, which the preprocessor then replaces with 1000 + 100.

If a number that is partially built by a macro expansion is produced, the preprocessor does not consider the result to be a single value. For example, the following will not result in the value 10.2 but in a syntax error.

```
#define a 10
a.2
```

Identifiers that are partially built from a macro expansion may not be produced. Therefore, the following example contains two identifiers and results in a syntax error:

```
#define d efg
abcd
```

Function-Like Macros

Function-like macro definition:

An identifier followed by a parameter list in parenthesis and the replacement tokens. The parameters are imbedded in the replacement code. White space cannot separate the identifier (which is the name of the macro) and the left parenthesis of the parameter list. A comma must separate each parameter. For portability, you should not have more than 31 parameters for a macro.

Function-like macro invocation:

An identifier followed by a list of arguments in parentheses. A comma must separate each argument. Once the preprocessor identifies a function-like macro invocation, argument substitution takes place. A parameter in the replacement code is replaced by the corresponding argument. Any macro invocations contained in the argument itself are completely replaced before the argument replaces its corresponding parameter in the replacement code.

The following line defines the macro SUM as having two parameters a and b and the replacement tokens (a + b):

```
#define SUM(a,b) (a + b)
```

This definition would cause the preprocessor to change the following statements (if the statements appear after the previous definition):

```
c = SUM(x,y);
c = d * SUM(x,y);
```

In the output of the preprocessor, these statements would appear as:

```
c = (x + y);
c = d * (x + y);
```

Use parentheses to ensure correct evaluation of replacement text. For example, the definition:

```
#define SQR(c) ((c) * (c))
```

requires parentheses around each parameter c in the definition in order to correctly evaluate an expression like:

```
y = SQR(a + b);
```

The preprocessor expands this statement to:

```
y = ((a + b) * (a + b));
```

Without parentheses in the definition, the correct order of evaluation is not preserved, and the preprocessor output is:

```
y = (a + b * a + b);
```

Arguments of the # and ## operators are converted before replacement of parameters in a function-like macro.

The number of arguments in a macro invocation must be the same as the number of parameters in the corresponding macro definition.

Commas in the macro invocation argument list do not act as argument separators when they are:

- in character constants
- in string literals
- surrounded by parentheses

Once defined, a preprocessor identifier remains defined and in scope independent of the scoping rules of the language. The scope of a macro definition begins at the definition and does not end until a corresponding **#undef** directive is encountered. If there is no corresponding **#undef** directive, the scope of the macro definition lasts until the end of the compilation unit.

A recursive macro is not fully expanded. For example, the definition

```
#define x(a,b) x(a+1,b+1) + 4
expands
   x(20,10)
to
  x(20+1.10+1) + 4
```

rather than trying to expand the macro x over and over within itself. After the macro x is expanded, it is a call to function x().

A definition is not required to specify replacement tokens. The following definition removes all instances of the token debug from subsequent lines in the current file: #define debug

You can change the definition of a defined identifier or macro with a second preprocessor #define directive only if the second preprocessor #define directive is preceded by a preprocessor #undef directive. The #undef directive nullifies the first definition so that the same identifier can be used in a redefinition.

Within the text of the program, the preprocessor does not scan character constants or string constants for macro invocations.

Example of #define Directives

The following program contains two macro definitions and a macro invocation that refers to both of the defined macros:

```
** This example illustrates #define directives.
**/
#include <stdio.h>
#define SQR(s) ((s) * (s))
#define PRNT(a,b) \
 printf("value 1 = %d\n", a); \
 printf("value 2 = %d\n", b);
int main(void)
 int x = 2;
 int y = 3;
    PRNT(SQR(x),y);
 return(0);
```

After being interpreted by the preprocessor, this program is replaced by code equivalent to the following:

```
#include <stdio.h>
int main(void)
 int x = 2;
 int y = 3;
    printf("value 1 = %d\n", ( (x) * (x) );
    printf("value 2 = %d\n", y);
 return(0);
```

This program produces the following output:

```
value 1 = 4
value 2 = 3
```

RELATED REFERENCES

- "Scope of Macro Names (#undef)"
- "Operator Precedence and Associativity" on page 71
- "Parenthesized Expressions ()" on page 76

Scope of Macro Names (#undef)

A preprocessor undef directive causes the preprocessor to end the scope of a preprocessor definition.

A preprocessor **#undef** directive has the form:

```
▶►—#—undef—identifier-
```

If the identifier is not currently defined as a macro, **#undef** is ignored.

Example of #undef Directives

The following directives define BUFFER and SQR:

#undef

```
#define BUFFER 512
#define SQR(x) ((x) * (x))
```

The following directives nullify these definitions:

```
#undef BUFFER
#undef SQR
```

Any occurrences of the identifiers BUFFER and SQR that follow these #undef directives are not replaced with any replacement tokens. Once the definition of a macro has been removed by an **#undef** directive, the identifier can be used in a new #define directive.

RELATED REFERENCES

• "Macro Definition and Expansion (#define)" on page 162

Operator

The # (single number sign) operator converts a parameter of a function-like macro into a character string literal. For example, if macro ABC is defined using the following directive:

```
#define ABC(x)
```

all subsequent invocations of the macro ABC would be expanded into a character string literal containing the argument passed to ABC. For example:

Invocation	Result of Macro Expansion
ABC(1)	"1"
ABC(Hello there)	"Hello there"

The # operator should not be confused with the null directive.

Use the # operator in a function-like macro definition according to the following rules:

- A parameter following # operator in a function-like macro is converted into a character string literal containing the argument passed to the macro.
- White-space characters that appear before or after the argument passed to the macro are deleted.
- Multiple white-space characters imbedded within the argument passed to the macro are replaced by a single space character.
- If the argument passed to the macro contains a string literal and if a \ (backslash) character appears within the literal, a second \ character is inserted before the original \ when the macro is expanded.
- If the argument passed to the macro contains a " (double quotation mark) character, a \ character is inserted before the " when the macro is expanded.
- The conversion of an argument into a string literal occurs before macro expansion on that argument.
- If more than one ## operator or # operator appears in the replacement list of a macro definition, the order of evaluation of the operators is not defined.
- If the result of the macro expansion is not a valid character string literal, the behavior is undefined.

Example of the # Operator

The following examples demonstrate the use of the # operator:

#define	STR(x)	#x
#define	XSTR(x)	STR(x)
#define	ONE	1

Invocation	Result of Macro	Expansion
------------	-----------------	-----------

STR(\n "\n" '\n')	"\n \"\\n\" '\\n'"	
STR(ONE)	"ONE"	
XSTR(ONE)	"1"	
XSTR("hello")	"\"hello\""	

RELATED REFERENCES

- "Null Directive (#)" on page 175
- "Function-Like Macros" on page 163
- "Macro Definition and Expansion (#define)" on page 162
- "Scope of Macro Names (#undef)" on page 165

Macro Concatenation with the ## Operator

The ## (double number sign) operator concatenates two tokens in a macro invocation (text and/or arguments) given in a macro definition.

If a macro XY was defined using the following directive:

```
#define XY(x,y)
```

the last token of the argument for x is concatenated with the first token of the argument for y.

Use the ## operator according to the following rules:

- The ## operator cannot be the very first or very last item in the replacement list of a macro definition.
- The last token of the item in front of the ## operator is concatenated with first token of the item following the ## operator.
- Concatenation takes place before any macros in arguments are expanded.
- · If the result of a concatenation is a valid macro name, it is available for further replacement even if it appears in a context in which it would not normally be available.
- If more than one ## operator and/or # operator appears in the replacement list of a macro definition, the order of evaluation of the operators is not defined.

Examples of the ## Operator

The following examples demonstrate the use of the ## operator:

	ArgArg(x, y)	x##y
#define	ArgText(x)	x##TEXT
#define	TextArg(x)	TEXT##x
#define	TextText	TEXT##text
#define	Jitter	1
#define	bug	2
#define	Jitterbug	3

Invocation Result of Macro Expansion

ArgArg(lady, bug)	"ladybug"
ArgText(con)	"conTEXT"
TextArg(book)	"TEXTbook"
TextText	"TEXTtext"

Invocation	Result of Macro Expansion
ArgArg(Jitter, bug)	3

RELATED REFERENCES

• "Macro Definition and Expansion (#define)" on page 162

Preprocessor Error Directive (#error)

A preprocessor error directive causes the preprocessor to generate an error message and causes the compilation to fail.

The #error directive has the form:



Use the #error directive as a safety check during compilation. For example, if your program uses preprocessor conditional compilation directives, put #error directives in the source file to prevent code generation if a section of the program is reached that should be bypassed.

For example, the directive

#error Error in TESTPGM1 - This section should not be compiled

generates the following error message:

Error in TESTPGM1 - This section should not be compiled

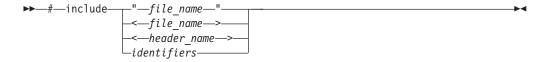
RELATED REFERENCES

• "Conditional Compilation Directives" on page 170

File Inclusion (#include)

A preprocessor include directive causes the preprocessor to replace the directive with the contents of the specified file.

A preprocessor **#include** directive has the form:



The preprocessor resolves macros contained in a **#include** directive. After macro replacement, the resulting token sequence must consist of a file name enclosed in either double quotation marks or the characters < and >.

For example:

#define MONTH <july.h> #include MONTH

If the file name is enclosed in double quotation marks, for example: #include "payroll.h"

the preprocessor treats it as a user-defined file, and searches for the file in a manner defined by the preprocessor.

If the file name is enclosed in angle brackets, for example: #include <stdio.h>

it is treated as a system-defined file, and the preprocessor searches for the file in a manner defined by the preprocessor.

The new-line and > characters cannot appear in a file name delimited by < and >. The new-line and " (double quotation marks) character cannot appear in a file name delimited by " and ", although > can.

Declarations that are used by several files can be placed in one file and included with **#include** in each file that uses them. For example, the following file defs.h contains several definitions and an inclusion of an additional file of declarations:

```
/* defs.h */
#define TRUE 1
#define FALSE 0
#define BUFFERSIZE 512
#define MAX ROW 66
#define MAX COLUMN 80
int hour;
int min;
int sec;
#include "mydefs.h"
```

You can embed the definitions that appear in defs.h with the following directive: #include "defs.h"

In the following example, a #define combines several preprocessor macros to define a macro that represents the name of the C standard I/O header file. A **#include** makes the header file available to the program.

```
#define C IO HEADER <stdio.h>
/* The following is equivalent to:
   #include <stdio.h>
 */
#include C IO HEADER
```

ANSI/ISO Standard Predefined Macro Names

Both C and C++ provide the following predefined macro names as specified in the ANSI/ISO C language standard:

```
Macro Name
            Description
```

```
__DATE__
               A character string literal containing the date when the source file
               was compiled.
               The value of __DATE__ changes as the compiler processes any
               include files that are part of your source program. The date is in
               the form:
                  "Mmm dd yyyy"
```

#include

where: Mmm Represents the month in an abbreviated form (Jan, Feb, Mar, Apr, May, Jun, Jul, Aug, Sep, Oct, Nov, or Dec). dd Represents the day. If the day is less than 10, the first d is a blank character. Represents the year. уууу FILE_ A character string literal containing the name of the source file. The value of FILE changes as the compiler processes include files that are part of your source program. It can be set with the **#line** directive. __LINE__ An integer representing the current source line number. The value of __LINE__ changes during compilation as the compiler processes subsequent lines of your source program. It can be set with the #line directive. __STDC__ The integer 0 (zero) indicates that C++ does not conform to the ANSI/ISO C language standard. **Note:** This macro is not defined if the language level is set to anything other than ANSI. This macro is only set for C compilers. __TIME__ A character string literal containing the time when the source file was compiled. The value of TIME changes as the compiler processes any include files that are part of your source program. The time is in the form: "hh:mm:ss" where: hh Represents the hour. Represents the minutes. Represents the seconds. SS cplusplus For C++ programs, this macro is set to the integer 1, indicating that the compiler is a C++ compiler. Note that this macro name has no

RELATED REFERENCES

• "Line Control (#line)" on page 174

trailing underscores.

• "Object-Like Macros" on page 162

Conditional Compilation Directives

A preprocessor conditional compilation directive causes the preprocessor to conditionally suppress the compilation of portions of source code. These directives test a constant expression or an identifier to determine which tokens the preprocessor should pass on to the compiler and which tokens should be bypassed during preprocessing. The directives are:

- #if
- #ifdef

- #else
- #ifndef
- #elif
- #endif

The preprocessor conditional compilation directive spans several lines:

- The condition specification line (beginning with **#if**, **#ifdef**, or **#ifndef**)
- · Lines containing code that the preprocessor passes on to the compiler if the condition evaluates to a nonzero value (optional)
- The **#elif** line (optional)
- Lines containing code that the preprocessor passes on to the compiler if the condition evaluates to a nonzero value (optional)
- The **#else** line (optional)
- · Lines containing code that the preprocessor passes on to the compiler if the condition evaluates to zero (optional)
- The preprocessor #endif directive

For each #if, #ifdef, and #ifndef directive, there are zero or more #elif directives, zero or one #else directive, and one matching #endif directive. All the matching directives are considered to be at the same nesting level.

You can nest conditional compilation directives. In the following directives, the first **#else** is matched with the **#if** directive.

```
#ifdef MACNAME
                /* tokens added if MACNAME is defined
                                                                      */
    if TEST <=10
                /* tokens added if MACNAME is defined and TEST <= 10 */
   else
                /* tokens added if MACNAME is defined and TEST > 10 */
   endif
#else
                /* tokens added if MACNAME is not defined
                                                                      */
#endif
```

Each directive controls the block immediately following it. A block consists of all the tokens starting on the line following the directive and ending at the next conditional compilation directive at the same nesting level.

Each directive is processed in the order in which it is encountered. If an expression evaluates to zero, the block following the directive is ignored.

When a block following a preprocessor directive is to be ignored, the tokens are examined only to identify preprocessor directives within that block so that the conditional nesting level can be determined. All tokens other than the name of the directive are ignored.

Only the first block whose expression is nonzero is processed. The remaining blocks at that nesting level are ignored. If none of the blocks at that nesting level has been processed and there is a **#else** directive, the block following the **#else** directive is processed. If none of the blocks at that nesting level has been processed and there is no **#else** directive, the entire nesting level is ignored.

RELATED REFERENCES

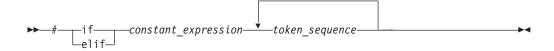
- "#if, #elif" on page 172
- "#ifdef" on page 172
- "#ifndef" on page 173
- "#else" on page 173

Conditional Compilation

"#endif" on page 174

#if, #elif

The **#if** and **#elif** directives compare the value of *constant_expression* to zero:



If the constant expression evaluates to a nonzero value, the lines of code that immediately follow the condition are passed on to the compiler.

If the expression evaluates to zero and the conditional compilation directive contains a preprocessor #elif directive, the source text located between the #elif and the next #elif or preprocessor #else directive is selected by the preprocessor to be passed on to the compiler. The **#elif** directive cannot appear after the preprocessor #else directive.

All macros are expanded, any defined() expressions are processed and all remaining identifiers are replaced with the token 0.

The constant_expression that is tested must be integer constant expressions with the following properties:

- No casts are performed.
- Arithmetic is performed using **long int** values.
- The constant_expression can contain defined macros. No other identifiers can appear in the expression.
- The constant_expression can contain the unary operator **defined**. This operator can be used only with the preprocessor keyword **#if** or **#elif**. The following expressions evaluate to 1 if the identifier is defined in the preprocessor, otherwise

```
defined identifier
defined(identifier)
For example:
#if defined(TEST1) || defined(TEST2)
```

Note: If a macro is not defined, a value of 0 (zero) is assigned to it. In the following example, TEST must be a macro identifier:

```
#if TEST >= 1
  printf("i = %d\n", i);
  printf("array[i] = %d\n", array[i]);
#elif TEST < 0
  printf("array subscript out of bounds \n");
#endif
```

#ifdef

The **#ifdef** directive checks for the existence of macro definitions.

If the identifier specified is defined as a macro, the lines of code that immediately follow the condition are passed on to the compiler.

The preprocessor **#ifdef** directive has the form:

Conditional Compilation



The following example defines MAX LEN to be 75 if EXTENDED is defined for the preprocessor. Otherwise, MAX LEN is defined to be 50.

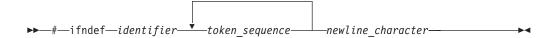
```
#ifdef EXTENDED
# define MAX LEN 75
#else
# define MAX LEN 50
#endif
```

#ifndef

The **#ifndef** directive checks whether a macro is not defined.

If the identifier specified is not defined as a macro, the lines of code immediately follow the condition are passed on to the compiler.

The preprocessor **#ifndef** directive has the form:



An identifier must follow the **#ifndef** keyword. The following example defines MAX LEN to be 50 if EXTENDED is not defined for the preprocessor. Otherwise, MAX LEN is defined to be 75.

```
#ifndef EXTENDED
# define MAX LEN 50
#else
# define MAX LEN 75
#endif
```

#else

If the condition specified in the #if, #ifdef, or #ifndef directive evaluates to 0, and the conditional compilation directive contains a preprocessor #else directive, the lines of code located between the preprocessor #else directive and the preprocessor **#endif** directive is selected by the preprocessor to be passed on to the compiler.

The preprocessor **#else** directive has the form:



RELATED REFERENCES

- "#if, #elif" on page 172
- "#ifdef" on page 172
- "#ifndef"

Conditional Compilation

#endif

The preprocessor #endif directive ends the conditional compilation directive.

It has the form:

```
▶►—#—endif—newline_character—
```

Examples of Conditional Compilation Directives

The following example shows how you can nest preprocessor conditional compilation directives:

```
#if defined(TARGET1)
   define SIZEOF_INT 16
   ifdef PHASE2
      define MAX PHASE 2
       define MAX PHASE 8
#elif defined(TARGET2)
   define SIZEOF INT 32
   define MAX PHASE 16
#else
   define SIZEOF INT 32
   define MAX PHASE 32
#endif
```

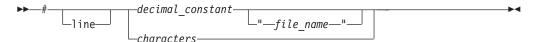
The following program contains preprocessor conditional compilation directives:

```
** This example contains preprocessor
** conditional compilation directives.
#include <stdio.h>
int main(void)
  static int array[] = \{1, 2, 3, 4, 5\};
  int i;
   for (i = 0; i \le 4; i++)
      array[i] *= 2;
#if TEST >= 1
   printf("i = %d\n", i);
  printf("array[i] = %d\n", array[i]);
#endif
   return(0);
```

Line Control (#line)

A preprocessor line control directive supplies line numbers for compiler messages. It causes the compiler to view the line number of the next source line as the specified number.

A preprocessor **#line** directive has the form:



In order for the compiler to produce meaningful references to line numbers in preprocessed source, the preprocessor inserts #line directives where necessary (for example, at the beginning and after the end of included text).

A file name specification enclosed in double quotation marks can follow the line number. If you specify a file name, the compiler views the next line as part of the specified file. If you do not specify a file name, the compiler views the next line as part of the current source file.

The token sequence on a #line directive is subject to macro replacement. After macro replacement, the resulting character sequence must consist of a decimal constant, optionally followed by a file name enclosed in double quotation marks.

Example of the #line Directive

You can use #line control directives to make the compiler provide more meaningful error messages. The following program uses #line control directives to give each function an easily recognizable line number:

```
** This example illustrates #line directives.
 **/
#include <stdio.h>
#define LINE200 200
int main(void)
   func 1();
   func 2();
#line 100
func 1()
   printf("Func_1 - the current line number is %d\n",_ _LINE_ _);
#line LINE200
func_2()
   printf("Func_2 - the current line number is %d\n",_ _LINE_ _);
This program produces the following output:
Func 1 - the current line number is 102
Func 2 - the current line number is 202
```

Null Directive (#)

The *null directive* performs no action. It consists of a single # on a line of its own.

The null directive should not be confused with the # operator or the character that starts a preprocessor directive.

(Null Directive)

In the following example, if MINVAL is a defined macro name, no action is performed. If MINVAL is not a defined identifier, it is defined 1.

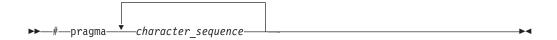
```
#ifdef MINVAL
#else
 #define MINVAL 1
#endif
```

RELATED REFERENCES

• "# Operator" on page 166

Pragma Directives (#pragma)

A pragma is an implementation-defined instruction to the compiler. It has the general form:



where character_sequence is a series of characters giving a specific compiler instruction and arguments, if any.

The character_sequence on a pragma is subject to macro substitutions. For example, #define XX_ISO_DATA isolated_call(LG_ISO_DATA) // ... #pragma XX_ISO_DATA

More than one pragma construct can be specified on a single #pragma directive. The compiler ignores unrecognized pragmas.

Chapter 10. Namespaces

A namespace is an optionally-named scope. You declare names inside a namespace as you would for a class or an enumeration. You can access names declared inside a namespace the same way you access a nested class name by using the scope resolution (::) operator. However namespaces do not have the additional features of classes or enumerations. The primary purpose of the namespace is to add an additional identifier (the name of the namespace) to a name.

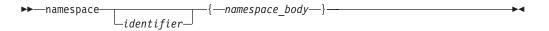
RELATED REFERENCES

• "C++ Scope Resolution Operator ::" on page 77

Defining Namespaces

In order to uniquely identify a namespace, use the **namespace** keyword.

Syntax — namespace



The *identifier* in an original namespace definition is the name of the namespace. The identifier may not be previously defined in the declarative region in which the original namespace definition appears, except in the case of extending namespace. If an identifier is not used, the namespace is an *unnamed namespace*.

RELATED REFERENCES

• "Unnamed Namespaces" on page 180

Declaring Namespaces

The identifier used for a namespace name should be unique. It should not be used previously as a global identifier.

```
namespace Raymond {
  // namespace body here...
}
```

In this example, Raymond is the identifier of the namespace. If you intend to access a namespace's elements, the namespace's identifier must be known in all translation units.

RELATED REFERENCES

"Global Scope" on page 2

Creating a Namespace Alias

An alternate name can be used in order to refer to a specific namespace identifier.

```
namespace INTERNATIONAL BUSINESS MACHINES {
 void f();
namespace IBM = INTERNATIONAL BUSINESS MACHINES;
```

In this example, the IBM identifier is an alias for INTERNATIONAL_BUSINESS_MACHINES. This is useful for referring to long namespace identifiers.

If a namespace name or alias is declared as the name of any other entity in the same declarative region, a compiler error will result. Also, if a namespace name defined at global scope is declared as the name of any other entity in any global scope of the program, a compiler error will result.

RELATED REFERENCES

• "Global Scope" on page 2

Creating an Alias for a Nested Namespace

Namespace definitions hold declarations. Since a namespace definition is a declaration itself, namespace definitions can be nested.

An alias can also be applied to a nested namespace.

```
namespace INTERNATIONAL_BUSINESS_MACHINES {
  int j;
    namespace NESTED IBM PRODUCT {
     void a() { j++; }
     int j;
     void b() { j++; }
namespace NIBM = INTERNATIONAL_BUSINESS_MACHINES::NESTED_IBM_PRODUCT
```

In this example, the NIBM identifier is an alias for the namespace NESTED IBM PRODUCT. This namespace is nested within the INTERNATIONAL BUSINESS MACHINES namespace.

Extending Namespaces

Namespaces are extensible. You can add subsequent declarations to a previously defined namespace. Extensions may appear in files separate from or attached to the original namespace definition. For example:

```
namespace X { // namespace definition
 int a;
  int b;
namespace X \{ // namespace extension
 int c;
 int d;
namespace Y \{ // equivalent to namespace X
  int a;
  int b;
 int c;
 int d;
```

In this example, namespace X is defined with a and b and later extended with c and d. namespace X now contains all four members. You may also declare all of the required members within one namespace. This method is represented by namespace Y. This namespace contains a, b, c, and d.

Namespaces and Overloading

You can overload functions across namespaces. For example:

```
// Original X.h:
  f(int);
// Original Y.h:
  f(char);
// Original program.c:
  #include "X.h"
  #include "Y.h"
void z()
  f('a'); // calls f(char) from Y.h
```

Namespaces can be introduced to the previous example without drastically changing the source code.

```
// New X.h:
namespace X {
  f(int);
// New Y.h:
namespace Y {
  f(char);
// New program.c:
  #include "X.h"
  #include "Y.h"
  using namespace X;
 using namespace Y;
void z()
  f('a'); // calls f() from Y.h
```

In program.c, function void z() calls function f(), which is a member of namespace Y. If you place the using directives in the header files, the source code for program.c remains unchanged.

RELATED REFERENCES

• "Chapter 11. Overloading" on page 185

Unnamed Namespaces

A namespace with no identifier before an opening brace produces an unnamed namespace. Each translation unit may contain its own unique unnamed namespace. The following example demonstrates how unnamed namespaces are useful.

```
#include <iostream>
using namespace std;
namespace {
  const int i = 4;
   int variable;
int main()
   cout << i << endl;</pre>
   variable = 100;
   return 0;
```

#include <iostream>

In the previous example, the unnamed namespace permits access to i and variable without using a scope resolution operator.

The following example illustrates an improper use of unnamed namespaces.

```
using namespace std;
namespace {
  const int i = 4;
int i = 2;
int main()
   cout << i << endl; // error</pre>
   return 0;
```

Inside main, i causes an error because the compiler cannot distinguish between the global name and the unnamed member with the same name. In order for the previous example to work, the namespace must be uniquely identified with an identifier and i must specify the namespace it is using.

You can extend an unnamed namespace within the same translation unit. For example:

```
#include <iostream>
using namespace std;
namespace {
   int variable;
   void funct (int);
namespace {
   void funct (int i) { cout << i << endl; }</pre>
```

```
int main()
   funct(variable);
   return 0;
```

both the prototype and definition for funct are members of the same unnamed namespace.

Note: Items defined in an unnamed namespace have internal linkage. Rather than using the keyword static to define items with internal linkage, define them in an unnamed namespace instead.

RELATED REFERENCES

- "Program Linkage" on page 4
- "Internal Linkage" on page 4

Namespace Member Definitions

A namespace can define its own members within itself or externally using explicit qualification. The following is an example of a namespace defining a member internally:

```
namespace A {
 void b() { /* definition */ }
```

Within namespace A member void b() is defined internally.

A namespace can also define its members externally using explicit qualification on the name being defined. The entity being defined must already be declared in the namespace and the definition must appear after the point of declaration in a namespace that encloses the declaration's namespace.

The following is an example of a namespace defining a member externally:

```
namespace A
 namespace B {
   void f();
  void B::f() { /* defined outside of B */ }
```

In this example, function f() is declared within namespace B and defined (outside B) in A.

Namespaces and Friends

Every name first declared in a namespace is a member of that namespace. If a friend declaration in a non-local class first declares a class or function, the friend class or function is a member of the innermost enclosing namespace.

The following is an example of this structure:

```
// f has not yet been defined
void z(int);
namespace A {
 class X {
    friend void f(X); // A::f is a friend
    };
```

```
// A::f is not visible here
 void f(X) { /* definition */} // f() is defined and known to be a friend
using A::x;
void z()
 A::f(x);
           // OK
 A::X::f(x); // error: f is not a member of A::X
```

In this example, function f() can only be called through namespace A using the call A::f(s);. Attempting to call function f() through class X using the A::X::f(x); call results in a compiler error. Since the friend declaration first occurs in a non-local class, the friend function is a member of the innermost enclosing namespace and may only be accessed through that namespace.

RELATED REFERENCES

"Friends" on page 225

Using Directive

C++ A using directive provides access to all namespace qualifiers and the scope operator. This is accomplished by applying the using keyword to a namespace identifier.

Syntax — Using directive

```
▶▶—using—namespace—name—;—
```

The *name* must be a previously defined namespace. The using directive may be applied at the global and local scope but not the class scope. Local scope takes precedence over global scope by hiding similar declarations.

If a scope contains a using directive that nominates a second namespace and that second namespace contains another using directive, the using directive from the second namespace will act as if it resides within the first scope.

```
namespace A {
 int i;
namespace B {
 int i;
 using namespace A;
void f()
 using namespace B;
 i = 7; // error
```

In this example, attempting to initialize i within function f() causes a compiler error, because function f() cannot know which i to call; i from namespace A, or i from namespace B.

RELATED REFERENCES

• "The using Declaration and Class Members" on page 238

The using Declaration and Namespaces

A using declaration provides access to a specific namespace member. This is accomplished by applying the using keyword to a namespace name with its corresponding namespace member.

Syntax — Using declaration

```
▶►—using—namespace—::—member-
```

In this syntax diagram, the qualifier name follows the using declaration and the member follows the qualifier name. For the declaration to work, the member must be declared inside the given namespace. For example:

```
namespace A {
 int i;
 int k;
 void f;
 void g;
using A::k
```

In this example, the using declaration is followed by A, the name of namespace A, which is then followed by the scope operator (::), and k. This format allows k to be accessed outside of namespace A through a using declaration. After issuing a using declaration, any extension made to that specific namespace will not be known at the point at which the using declaration occurs.

In the incremental compiler, all names in all extents of a namespace will be made visible by a using declaration regardless of their positions.

Overloaded versions of a given function must be included in the namespace prior to that given function's declaration. A using declaration may appear at namespace, block and class scope.

RELATED REFERENCES

"The using Declaration and Class Members" on page 238

Explicit Access

C++ To explicitly qualify a member of a namespace, use the namespace identifier with a :: scope resolution operator.

Syntax — Explicit access qualification

```
▶▶—namespace name—::—member-
```

For example:

```
namespace VENDITTI {
 void j()
VENDITTI::j();
```

In this example, the scope resolution operator provides access to the function j held within namespace VENDITTI. The scope resolution operator :: is used to access identifiers in both global and local namespaces. Any identifier in an application can be accessed with sufficient qualification. Explicit access cannot be applied to an unnamed namespace.

RELATED REFERENCES

• "C++ Scope Resolution Operator ::" on page 77

Chapter 11. Overloading

If you specify more than one definition for a function name or an operator in the same scope, you have *overloaded* that function name or operator.

An *overloaded declaration* is a declaration that had been declared with the same name as a previously declared declaration in the same scope, except that both declarations have different types.

If you call an overloaded function name or operator, the compiler determines the most appropriate definition to use by comparing the argument types you used to call the function or operator with the parameter types specified in the definitions. The process of selecting the most appropriate overloaded function or operator is called *overload resolution*.

RELATED REFERENCES

- "Overloading Functions"
- "Overloading Operators" on page 187
- "Overload Resolution" on page 196

Overloading Functions

You overload a function name f by declaring more than one function with the name f in the same scope. The declarations of f must differ from each other by the types and/or the number of arguments in the argument list. When you call a overloaded function named f, the correct function is selected by comparing the argument list of the function call with the parameter list of each of the overloaded candidate functions with the name f. A *candidate function* is a function that can be called based on the context of the call of the overloaded function name.

Consider a function print, which displays an **int**. As shown in the following example, you can overload the function print to display other types, for example, **double** and **char***. You can have three functions with the same name, each performing a similar operation on a different data type:

```
#include <iostream>
using namespace std;

void print(int i) {
   cout << " Here is int " << i << endl;
}

void print(double f) {
   cout << " Here is float " << f << endl;
}

void print(char* c) {
   cout << " Here is char* " << c << endl;
}

int main() {
   print(10);
   print(10.10);
   print("ten");
}</pre>
```

The following is the output of the above example:

```
Here is int 10
Here is float 10.1
Here is char* ten
```

Restrictions on Overloaded Functions

You cannot overload the following function declarations if they appear in the same scope. Note that this list applies only to explicitly declared functions and those that have been introduced through **using** declarations:

• Function declarations that differ only by return type. For example, you cannot declare the following declarations:

```
int f();
float f();
```

 Member function declarations that have the same name and the same parameter types, but one of these declarations is a static member function declaration. For example, you cannot declare the following two member function declarations of f():

```
struct A {
 static int f();
  int f();
```

- Member function template delcarations that have the same name, the same parameter types, and the same template parameter lists, but one of these declarations is a static template member function declaration.
- Function declarations that have equivalent parameter declarations. These declarations are not allowed because they would be declaring the same function.
- Function declarations with parameters that differ only by the use of typedef names that represent the same type. Note that a typedef is a synonym for another type, not a separate type. For example, the following two declarations of f() are declarations of the same function:

```
typedef int I;
void f(float, int);
void f(float I);
```

• Function declarations with parameters that differ only because one is a pointer and the other is an array. For example, the following are declarations of the same function:

```
f(char*):
f(char[10]);
```

The first array dimension is insignificant when differentiating parameters; all other array dimensions are significant. For example, the following are declarations of the same function:

```
g(char(*)[20]);
g(char[5][20]);
```

The following two declarations are *not* equivalent:

```
g(char(*)[20]);
g(char(*)[40]);
```

Function declarations with parameters that differ only because one is a function type and the other is a pointer to a function of the same type. For example, the following are declarations of the same function:

```
void f(int(float));
void f(int (*)(float));
```

• Function declarations with parameters that differ only because of const and volatile qualifiers. This only applies if you apply any of these qualifiers appear at the outermost level of an parameter type specification. For example, the following are declarations of the same function:

```
int f(int);
int f(const int);
int f(volatile int);
```

Note that you can differentiate parameters with **const** and **volatile** qualifiers if you apply these qualifiers within a parameter type specification. For example, the following declarations are *not* equivalent:

```
void g(int*);
void g(const int*);
void g(volatile int*);
```

The following declarations are also not equivalent:

```
void g(float&);
void g(const float&);
void g(volatile float&);
```

Function declarations with parameters that differ only because their default arguments differ. For example, the following are declarations of the same function:

```
void f(int);
void f(int i = 10)
```

Multiple functions with extern "C" language-linkage and the same name, regardless of whether their parameter lists are different.

RELATED REFERENCES

- "The using Declaration and Namespaces" on page 183
- "typedef" on page 29
- "volatile and const Qualifiers" on page 50
- "Linkage Specifications Linking to Non-C++ Programs" on page 6

Overloading Operators

C++= You can redefine or overload the function of most built-in operators in C++. These operators can be overloaded globally or on a class-by-class basis. Overloaded operators are implemented as functions and can be member functions or global functions.

An overloaded operator is called an operator function. You declare an operator function with the keyword operator preceding the operator. Overloaded operators are distinct from overloaded functions, but like overloaded functions, they are distinguished by the number and types of operands used with the operator.

Consider the standard + (plus) operator. When this operator is used with operands of different standard types, the operators have slightly different meanings. For example, the addition of two integers is not implemented in the same way as the addition of two floating-point numbers. C++ allows you to define your own meanings for the standard C++ operators when they are applied to class types. In the following example, a class called complx is defined to model complex numbers, and the + (plus) operator is redefined in this class to add two complex numbers.

```
// This example illustrates overloading the plus (+) operator.
```

```
#include <iostream>
using namespace std;
class complx
```

```
double real,
public:
     complx( double real = 0., double imag = 0.); // constructor
     complx operator+(const complx&) const;
                                              // operator+()
};
// define constructor
complx::complx( double r, double i )
     real = r; imag = i;
// define overloaded + (plus) operator
complx complx::operator+ (const complx& c) const
     complx result;
     result.real = (this->real + c.real);
     result.imag = (this->imag + c.imag);
     return result;
int main()
{
     complx x(4,4);
     complx y(6,6);
     complx z = x + y; // calls complx::operator+()
}
```

You can overload any of the following operators:

```
1
                                                         /=
                                                                 %=
        =
                <
                                +=
        &=
                l =
                        <<
                                >>
                                        <<=
                                                                 !=
                        \Pi
                &&
<=
        >=
                                                                 ->
()
        []
                new
                        delete
                                new[]
                                        delete[]
```

where () is the function call operator and [] is the subscript operator.

You can overload both the unary and binary forms of the following operators:

```
&
```

You cannot overload the following operators:

```
?:
::
```

You cannot overload the preprocessor symbols # and ##.

An operator function can be either a non-static member function, or a non-member function with at least one parameter that has class, reference to class, enumeration, or reference to enumeration type.

You cannot change the precedence, grouping, or the number of operands of an operator.

An overloaded operator (except for the function call operator) cannot have default arguments or an ellipsis in the argument list.

You must declare the overloaded =, [], [], and -> operators as nonstatic member functions to ensure that they receive lvalues as their first operands.

The operators new, delete, new[], and delete[] do not follow the general rules described in this section.

All operators except the = operator are inherited.

RELATED REFERENCES

• "Free Store" on page 269

Overloading Unary Operators

You overload a unary operator with either a nonstatic member function that has no parameters, or a non-member function that has one parameter. Suppose a unary operator @ is called with the statement @t, where t is an object of type T. A nonstatic member function that overloads this operator would have the following form:

```
return type operator@()
```

A nonmember function that overloads the same operator would have the following form:

```
return type operator@(T)
```

An overloaded unary operator may return any type.

The following example overloads the ! operator:

```
#include <iostream>
using namespace std;
struct X { };
void operator!(X) {
  cout << "void operator!(X)" << endl;</pre>
struct Y {
  void operator!() {
    cout << "void Y::operator!()" << endl;</pre>
};
struct Z { };
int main() {
 X ox; Y oy; Z oz;
  !ox;
 !oy;
// !oz;
```

The following is the output of the above example:

```
void operator!(X)
void Y::operator!()
```

The operator function call !ox is interpreted as operator! (x). The call !oy is interpreted as y.operator!(). (The compiler would not allow !oz because the! operator has not been defined for class Z.)

RELATED REFERENCES

"Unary Expressions" on page 86

Overloading Binary Operators

You overload a binary unary operator with either a nonstatic member function that has one parameter, or a non-member function that has two parameters. Suppose a binary operator @ is called with the statement t @ u, where t is an object of type T, and u is an object of type U. A nonstatic member function that overloads this operator would have the following form:

```
return type operator@(T)
```

A nonmember function that overloads the same operator would have the following

```
return type operator@(T, U)
```

An overloaded binary operator may return any type.

The following example overloads the * operator:

```
struct X {
  // member binary operator
 void operator*(int) { }
// non-member binary operator
void operator*(X, float) { }
int main() {
 Χх;
 int y = 10:
 float z = 10;
 x * y;
 x * Z;
```

The call x * y is interpreted as x.operator*(y). The call x * z is interpreted as operator*(x, z).

RELATED REFERENCES

• "Binary Expressions" on page 97

Overloading Assignments

Solution You overload the assignment operator, **operator**, with a nonstatic member function that has only one parameter. You cannot declare an overloaded assignment operator that is a nonmember function. The following example shows how you can overload the assignment operator for a particular class:

```
struct X {
  int data;
  X& operator=(X& a) { return a; }
  X& operator=(int a) {
    data = a;
    return *this;
};
int main() {
 X x1, x2;
 x1 = x2;  // call x1.operator=(x2)
x1 = 5;  // call x1.operator=(5)
```

The assignment x1 = x2 calls X& X::operator=(X&). The assignment x1 = 5 calls X& X::operator=(int).

Note that because the compiler implicitly declares a copy assignment operator for a class if you do not define one yourself, the copy assignment operator of a derived class will hide the copy assignment operator of its base class. However, you can declare any copy assignment operator as virtual. The following example demonstrates this:

```
#include <iostream>
using namespace std;
struct A {
  A& operator=(char) {
    cout << "A& A::operator=(char)" << endl;</pre>
    return *this;
  virtual A& operator=(const A&) {
    cout << "A& A::operator=(const A&)" << endl;</pre>
    return *this;
  }
};
struct B : A {
    B& operator=(char) {
      cout << "B& B::operator=(char)" << endl;</pre>
      return *this;
    virtual B& operator=(const A&) {
      cout << "B& B::operator=(const A&)" << endl;</pre>
      return *this;
};
struct C : B { };
int main() {
 B b1;
  B b2;
  A* ap1 = &b1;
 A* ap2 = &b1;
  *ap1 = 'z';
  *ap2 = b2;
  C c1;
// c1 = 'z';
```

The following is the output of the above example:

```
A& A::operator=(char)
B& B::operator=(const A&)
```

The assignment *ap1 = 'z' calls A& A::operator=(char). Because this operator has not been declared virtual, the compiler chooses the function based on the type of the pointer ap1. The assignment *ap2 = b2 calls B& B::operator=(const &A). Because this operator has been declared virtual, the compiler chooses the function based on the type of the object that the pointer ap1 points to. The compiler would not allow the assignment c1 = 'z' because the implicitly declared copy assignment operator declared in class C hides B& B::operator=(char).

RELATED REFERENCES

- "Member Functions" on page 211
- "Copy Assignment Operators" on page 279

"Assignment Expressions" on page 107

Overloading Function Calls

The function call operator, when overloaded, does not modify how functions are called. Rather, it modifies how the operator is to be interpreted when applied to objects of a given type.

You overload the function call operator, operator(), with a nonstatic member function that has any number of parameters. If you overload a function call operator for a class its declaration will have the following form:

```
return_type operator()(parameter_list)
```

Unlike all other overloaded operators, you can provide default arguments and ellipses in the argument list for the function call operator.

The following example demonstrates how the compiler interprets function call operators:

```
struct A {
 void operator()(int a, char b, ...) { }
 void operator()(char c, int d = 20) { }
};
int main() {
 A a;
a(5, 'z', 'a', 0);
 a('z');
// a();
```

The function call a(5, 'z', 'a', 0) is interpreted as a.operator()(5, 'z', 'a', 0). This calls void A::operator()(int a, char b, ...). The function call a('z') is interpreted as a.operator()('z'). This calls void A::operator()(char c, int d = 20). The compiler would not allow the function call a() because its argument list does not match any function call parameter list defined in class A.

The following example demonstrates an overloaded function call operator:

```
class Point {
private:
 int x, y;
public:
 Point(): x(0), y(0) { }
 Point& operator()(int dx, int dy) {
   x += dx;
   y += dy;
   return *this;
};
int main() {
 Point pt;
 // Offset this coordinate x with 3 points
 // and coordinate y with 2 points.
 pt(3, 2);
```

The above example reinterprets the function call operator for objects of class Point. If you treat an object of Point like a function and pass it two integer aguments, the function call operator will add the values of the arguments you passed to Point::x and Point::y respectively.

RELATED REFERENCES

• "Function Calls ()" on page 78

Overloading Subscripting

You overload operator[] with a nonstatic member function that has only one parameter. The following example is a simple array class that has an overloaded subscripting operator. The overloaded subscripting operator throws an exception if you try to access the array outside of its specified bounds:

```
#include <iostream>
using namespace std;
template <class T> class MyArray {
private:
 T* storage;
  int size;
public:
  MyArray(int arg = 10) {
    storage = new T[arg];
    size = arg;
  MyArray() {
    delete[] storage;
    storage = 0;
  T& operator[](const int location) throw (const char *);
};
template <class T> T& MyArray<T>::operator[](const int location)
  throw (const char *) {
    if (location < 0 | | location >= size) throw "Invalid array access";
    else return storage[location];
}
int main() {
  try {
     MyArray<int> x(13);
     x[0] = 45;
     x[1] = 2435;
     cout \ll x[0] \ll endl;
     cout \ll x[1] \ll endl;
     x[13] = 84;
  catch (const char* e) {
    cout << e << endl;</pre>
```

The following is the output of the above example:

2435 Invalid array access

The expression x[1] is interpreted as x.operator[](1) and calls int& MyArray<int>::operator[](const int).

RELATED REFERENCES

• "Array Subscript [] (Array Element Specification)" on page 80

Overloading Class Member Access

You overload operator-> with a nonstatic member function that has no parameters. The following example demonstrates how the compiler interprets overloaded class member access operators:

```
struct Y {
 void f() { };
struct X {
Y* ptr;
Y* operator->() {
  return ptr;
};
};
int main() {
 Х х;
 x->f();
```

The statement $x \rightarrow f()$ is interpreted as $(x \cdot operator \rightarrow ()) \rightarrow f()$.

The **operator-**> is used (often in conjunction with the pointer-dereference operator) to implement "smart pointers." These pointers are objects that behave like normal pointers except they perform other tasks when you access an object through them, such as automatic object deletion (either when the pointer is destroyed, or the pointer is used to point to another object), or reference counting (counting the number of smart pointers that point to the same object, then automatically deleting the object when that count reaches zero).

One example of a smart pointer is included in the C++ Standard Library called auto ptr. You can find it in the <memory> header. The auto ptr class implements automatic object deletion.

RELATED REFERENCES

• "Arrow Operator ->" on page 82

Overloading Increment and Decrement

You overload the prefix increment operator ++ with either a nonmember function operator that has one argument of class type or a reference to class type, or with a member function operator that has no arguments.

In the following example, the increment operator is overloaded in both ways:

```
class X {
public:
  // member prefix ++x
 void operator++() { }
};
class Y { };
// non-member prefix ++y
void operator++(Y&) { }
```

```
int main() {
 Хх;
 Υу;
 // calls x.operator++()
 ++x;
 // explicit call, like ++x
 x.operator++();
 // calls operator++(y)
 // explicit call, like ++y
 operator++(y);
```

The postfix increment operator ++ can be overloaded for a class type by declaring a nonmember function operator operator++() with two arguments, the first having class type and the second having type int. Alternatively, you can declare a member function operator operator++() with one argument having type int. The compiler uses the int argument to distinguish between the prefix and postfix increment operators. For implicit calls, the default value is zero.

```
For example:
class X {
public:
  // member postfix x++
  void operator++(int) { };
class Y { };
// nonmember postfix y++
void operator++(Y&, int) { };
int main() {
 X x;
  Y y;
  // calls x.operator++(0)
  // default argument of zero is supplied by compiler
  // explicit call to member postfix x++
  x.operator++(0);
  // calls operator++(y, 0)
  // explicit call to non-member postfix y++
  operator++(y, 0);
```

The prefix and postfix decrement operators follow the same rules as their increment counterparts.

RELATED REFERENCES

- "Member Functions" on page 211
- "Increment ++" on page 87
- "Decrement --" on page 88

Overload Resolution

The process of selecting the most appropriate overloaded function or operator is called overload resolution.

Suppose that f is an overloaded function name. When you call the overloaded function f(), the compiler creates a set of candidate functions. This set of functions includes all of the functions named f that can be accessed from the point where you called f(). The compiler may include as a candidate function an alternative representation of one of those accessible functions named f to facilitate overload resolution.

After creating a set of candidate functions, the compiler creates a set of viable functions. This set of functions is a subset of the candidate functions. The number of parameters of each viable function agrees with the number of arguments you used to call f().

The compiler chooses the best viable function, the function declaration that the C++ runtime will use when you call f(), from the set of viable functions. The compiler does this by *implicit conversion sequences*. An implicit conversion sequence is the sequence of conversions required to convert an argument in a function call to the type of the corresponding parameter in a function declaration. The implicit conversion sequences are ranked; some implicit conversion sequences are better than others. The compiler tries to find one viable function in which all of its parameters have either better or equal-ranked implicit conversion sequences than all of the other viable functions. The viable function that the compiler finds is the best viable function. The compiler will not allow a program in which the compiler was able to find more than one best viable function.

You can override an exact match by using an explicit cast. In the following example, the second call to f() matches with f(void*):

```
void f(int) { };
void f(void*) { };
int main() {
  f(0xaabb);
                     // matches f(int);
  f((void*) 0xaabb); // matches f(void*)
```

RELATED REFERENCES

"Implicit Conversion Sequences"

Implicit Conversion Sequences

An implicit conversion sequence is the sequence of conversions required to convert an argument in a function call to the type of the corresponding parameter in a function declaration.

The compiler will try to determine an implicit conversion sequence for each argument. It will then categorize each implicit conversion sequence in one of three categories and rank them depending on the category. The compiler will not allow any program in which it cannot find an implicit conversion sequence for an argument.

The following are the three categories of conversion sequences in order from best

- Standard conversion sequences
- User-defined conversion sequences
- · Ellipsis conversion sequences

Note: Two standard conversion sequences or two user-defined conversion sequences may have different ranks.

Standard Conversion Sequences

Standard conversion sequences are categorized in one of three ranks. The ranks are listed in order from best to worst:

- Exact match: This rank includes the following conversions:
 - Identity conversions
 - Lvalue-to-rvalue conversions
 - Array-to-pointer conversions
 - Qualification conversions
- Promotion: This rank includes integral and floating point promotions.
- Conversion: This rank includes the following conversions:
 - Integral and floating-point conversions
 - Floating-integral conversions
 - Pointer conversions
 - Pointer-to-member conversions
 - Boolean conversions

The compiler ranks a standard conversion sequence by its worst-ranked standard conversion. For example, if a standard conversion sequence has a floating-point conversion, then that sequence has conversion rank.

User-Defined Conversion Sequences

A user-defined conversion sequence consists of the following:

- A standard conversion sequence
- A user-derfined conversion
- A second standard conversion sequence

A user-defined conversion sequence A is better than a user-defined conversion sequence B if the both have the same user-defined conversion function or constructor, and the second standard conversion sequence of A is better than the second standard conversion sequence of B.

Ellipsis Conversion Sequences

An ellipsis conversion sequence occurs when the compiler matches an argument in a function call with a corresponding ellipsis parameter.

RELATED REFERENCES

- "Lvalue-to-Rvalue Conversions" on page 117
- "Pointer Conversions" on page 118
- "Qualification Conversions" on page 120
- "Integral Conversions" on page 117
- "Floating-Point Conversions" on page 118
- "Boolean Conversions" on page 117

Resolving Addresses of Overloaded Functions

Solution If you use an overloaded function name f without any arguments, that name can refer to a function, a pointer to a function, a pointer to member function, or a specialization of a function template. Because you did not provide any arguments, the compiler cannot perform overload resolution the same way it would for a function call or for the use of an operator. Instead, the compiler will try to choose the best viable function that matches the type of one of the following expressions, depending on where you have used f:

- An object or reference you are initializing
- The left side of an assignment
- A parameter of a function or a user-defined operator
- The return value of a function, operator, or conversion
- An explicit type conversion

If the compiler chose a declaration of a non-member function or a static member function when you used f, the compiler matched the declaration with an expression of type pointer-to-function or reference-to-function. If the compiler chose a declaration of a nonstatic member function, the compiler matched that declaration with an expression of type pointer-to-member function. The following example demonstrates this:

```
struct X {
 int f(int) { return 0; }
 static int f(char) { return 0; };
int main() {
 int (X::*a)(int) = &X::f;
// int (*b)(int) = &X::f;
 int (*c)(int) = &X::f;
```

The compiler will not allow the initialization of the function pointer b. No non-member function or static function of type int(int) has been declared.

If f is a template function, the compiler will perform template argument deduction to determine which template function to use. If successful, it will add that function to the list of viable functions. If there is more than one function in this set, including a non-template function, the compiler will eliminate all template functions from the set. If there are only template functions in this set, the compiler will choose the most specialized template function. The following example demonstrates this:

```
template<class T> int f(T) { return 0; }
template<> int f(int) { return 0; }
int f(int) { return 0; }
int main() {
 int (*a)(int) = f;
 a(1);
```

The function call a(1) calls int f(int).

RELATED REFERENCES

- "Chapter 7. Functions" on page 123
- "Pointers to Functions" on page 141
- "Pointers to Members" on page 214
- "Function Templates" on page 294
- "Explicit Specialization" on page 306

Chapter 12. Classes

A *class* is a mechanism for creating user-defined data types. It is similar to the C-language structure data type. In C, a structure is composed of a set of data members. In C++, a class type is like a C structure, except that a class is composed of a set of data members and a set of operations that can be performed on the class.

In C++, a class type can be declared with the keywords **union**, **struct**, or **class**. A union object can hold any one of a set of named members. Structure and class objects hold a complete set of members. Each class type represents a unique set of class members including data members, member functions, and other type names. The default access for members depends on the class key:

- The members of a class declared with the keyword **class** are private by default. A class is inherited privately by default.
- The members of a class declared with the keyword **struct** are public be default. A structure is inherited publicly by default.
- The members of a union (declared with the keyword **union**) are public by default. A union cannot be used as a base class in derivation.

Once you create a class type, you can declare one or more objects of that class type. For example:

```
class X
{
    /* define class members here */
};
void main()
{
    X xobject1;    // create an object of class type X
    X xobject2;    // create another object of class type X
}
```

You may have *polymorphic* classes in C++. Polymorphism is the ability to use a function name that appears in different classes (related by inheritance), without knowing exactly the class the function belongs to at compile time.

C++ allows you to redefine standard operators and functions through the concept of overloading. Operator overloading facilitates data abstraction by allowing you to use classes as easily as built-in types.

RELATED REFERENCES

- "Structures" on page 36
- "Chapter 13. Class Members and Friends" on page 209
- "Chapter 14. Inheritance" on page 231
- "Chapter 11. Overloading" on page 185
- "Virtual Functions" on page 249

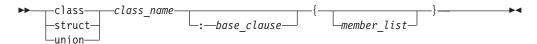
Declaring Class Types

A class declaration creates a unique type class name.

Declaring Class Objects

A class specifier is a type specifier used to declare a class. Once a class specifier has been seen and its members declared, a class is considered to be defined even if the member functions of that class are not yet defined. A class specifier has the following form:

Syntax — Class Specifier



The class_name is a unique identifier that becomes a reserved word within its scope. Once a class name is declared, it hides other declarations of the same name within the enclosing scope.

The member_list specifies the class members, both data and functions, of the class class_name. If the member_list of a class is empty, objects of that class have a nonzero size. You can use a class_name within the member_list of the class specifier itself as long as the size of the class is not required..

The base_clause specifies the base class or classes from which the class class_name inherits members. If the base_clause is not empty, the class class_name is called a derived class.

A structure is a class declared with the class_key struct. The members and base classes of a strucutre are public by default. A union is a class declared with the class_key union. The members of a union are public by default; a union holds only one data member at a time.

An aggregate class is a class that has no user-defined constructors, no private or protected non-static data members, no base classes, and no virtual functions.

RELATED REFERENCES

- "Class Member Lists" on page 209
- "Derivation" on page 233

Using Class Objects

Solution You can use a class type to create instances or *objects* of that class type. For example, you can declare a class, structure, and union with class names X, Y, and Z respectively:

```
class X {
 // members of class X
struct Y {
 // members of struct Y
union Z {
 // members of union Z
```

You can then declare objects of each of these class types. Remember that classes, structures, and unions are all types of C++ classes.

In C++, unlike C, you do not need to precede declarations of class objects with the keywords **union**, **struct**, and **class** unless the name of the class is hidden. For example:

When you declare more than one class object in a declaration, the declarators are treated as if declared individually. For example, if you declare two objects of class S in a single declaration:

```
class S { /* ... */ };
void main()
{
      S S,T; // declare two objects of class type S
}
this declaration is equivalent to:
class S { /* ... */ };
void main()
{
      SS;
      class S T;
                       // keyword class is required
                       // since variable S hides class type S
}
but is not equivalent to:
class S { /* ... */ };
void main()
{
      SS;
      S T;
                       // error, S class type is hidden
}
```

You can also declare references to classes, pointers to classes, and arrays of classes. For example:

Objects of class types that are not copy restricted can be assigned, passed as arguments to functions, and returned by functions.

Declaring Class Objects

RELATED REFERENCES

- "Chapter 12. Classes" on page 199
- "Structures" on page 36
- "Unions" on page 42
- "References" on page 68
- "Scope of Class Names" on page 203

Classes and Structures

The C++ class is an extension of the C-language structure. Because the only difference between a structure and a class is that structure members have public access by default and a class members have private access by default, you can use the keywords **class** or **struct** to define equivalent classes.

For example, in the following code fragment, the class X is equivalent to the structure Y:

```
class X {
 // private by default
 int a;
public:
  // public member function
 int f() { return a = 5; };
struct Y {
 // public by default
 int f() { return a = 5; };
private:
 // private data member
 int a;
};
```

If you define a structure and then declare an object of that structure using the keyword class, the members of the object are still public by default. In the following example, main() has access to the members of obj X even though obj X has been declared using an elaborated type specifier that uses the class key class:

```
#include <iostream>
using namespace std;
struct X {
int a;
int b;
};
class X obj X;
int main() {
  obj_X.a = 0;
  obj X.b = 1;
  cout << "Here are a and b: " << obj X.a << " " << obj X.b << endl;
```

The following is the output of the above example:

```
Here are a and b: 0 1
```

RELATED REFERENCES

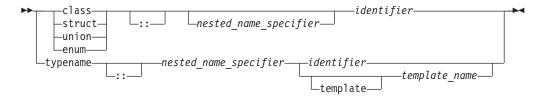
• "Structures" on page 36

Scope of Class Names

A class declaration introduces the class name into the scope where it is declared. Any class, object, function or other declaration of that name in an enclosing scope is hidden.

If a class name is declared in the same scope as a function, enumerator, or object with the same name, you must refer to that class using an elaborated type specifier:

Syntax — Elaborated Type Specifier



Syntax — Nested Name Specifier

```
-class name-
−namespace name—
                           —template—nested name specifier
                           \sqsubseteqnested_name_specifier-
```

The following example must use an elaborated type specifier to refer to class A because this class is hidden by the definition of the function A():

```
class A { };
void A (class A*) { };
int main()
      class A* x;
      A(x);
```

The declaration class A* x is an elaborated type specifier. Declaring a class with the same name of another function, enumerator, or object as demonstrated above is not recommended.

An elaborated type specifier can also be used in the incomplete declaration of a class type to reserve the name for a class type within the current scope.

RELATED REFERENCES

- "Scope" on page 1
- "Incomplete Class Declarations" on page 204

Incomplete Class Declarations

■ C++■ An incomplete class declaration is a class declaration that does not define any class members. You cannot declare any objects of the class type or refer to the members of a class until the declaration is complete. However, an incomplete declaration allows you to make specific references to a class prior to its definition as long as the size of the class is not required.

For example, you can define a pointer to the structure first in the definition of the structure second. Structure first is declared in an incomplete class declaration prior to the definition of second, and the definition of oneptr in structure second does not require the size of first:

```
struct first;
                        // incomplete declaration of struct first
struct second
                        // complete declaration of struct second
                        // pointer to struct first refers to
     first* oneptr;
                        // struct first prior to its complete
                        // declaration
     first one:
                        // error, you cannot declare an object of
                        // an incompletely declared class type
     int x, y;
};
struct first
                        // complete declaration of struct first
      second two;
                        // define an object of class type second
      int z;
};
```

If you declare a class with an empty member list, it is a complete class declaration. For example:

```
// incomplete class declaration
class X;
class Z {};
                        // empty member list
class Y
public:
     X yobj;
                        // error, cannot create an object of an
                        // incomplete class type
      Z zobj;
                       // valid
};
```

RELATED REFERENCES

"Class Member Lists" on page 209

Nested Classes

C++ A nested class is declared within the scope of another class. The name of a nested class is local to its enclosing class. Unless you use explicit pointers, references, or object names, declarations in a nested class can only use visible constructs, including type names, static members, and enumerators from the enclosing class and global variables.

Member functions of a nested class follow regular access rules and have no special access privileges to members of their enclosing classes. Member functions of the enclosing class have no special access to members of a nested class. The following example demonstrates this:

```
class A {
  int x;
  class B { };
  class C {
    // The compiler cannot allow the following
    // declaration because A::B is private:
   // B b;
    int y;
   void f(A* p, int i) {
   // The compiler cannot allow the following
    // statement because A::x is private:
        p->x = i;
  };
  void g(C* p) {
    // The compiler cannot allow the following
    // statement because C::y is private:
       int z = p->y;
};
int main() { }
```

The compiler would not allow the declaration of object b because class A::B is private. The compiler would not allow the statement p->x = i because A::x is private. The compiler would not allow the statement int z = p->y because C::y is private.

You can define member functions and static data members of a nested class in namespace scope. For example, in the following code fragment, you can access the static members x and y and member functions f() and g() of the nested class nested by using a qualified type name. Qualified type names allow you to define a **typedef** to represent a qualified class name. You can then use the **typedef** with the :: (scope resolution) operator to refer to a nested class or class member, as shown in the following example:

```
class outside
public:
      class nested
      public:
            static int x;
            static int y;
            int f();
            int g();
      };
};
int outside::nested::x = 5;
int outside::nested::f() { return 0; };
                                        // define a typedef
typedef outside::nested outnest;
                                        // use typedef with ::
int outnest::y = 10;
int outnest::g() { return 0; };
```

However, using a typedef to represent a nested class name hides information and may make the code harder to understand.

Scope of Class Names

You cannot use a typedef name in an elaborated type specifier. To illustrate, you cannot use the following declaration in the above example:

```
class outnest obj;
```

A nested class may inherit from private members of its enclosing class. The following example demonstrates this:

```
class A {
private:
 class B { };
 B *z;
 class C : private B {
 private:
      В у;
      A::B y2;
     C *x;
        A::C *x2;
//
    };
};
```

The nested class A::C inherits from A::B. The compiler does not allow the declarations A::B y2 and A::C *x2 because both A::B and A::C are private.

RELATED REFERENCES

- "Scope of Class Names" on page 203
- "Member Functions" on page 211
- "Member Access" on page 223
- "Static Members" on page 218
- "typedef" on page 29
- "C++ Scope Resolution Operator ::" on page 77

Local Classes

Letter A local class is declared within a function definition. Declarations in a local class can only use type names, enumerations, static variables from the enclosing scope, as well as external variables and functions.

```
For example:
int x;
                              // global variable
void f()
                              // function definition
     static int y;
                              // static variable y can be used by
                              // local class
                              // auto variable x cannot be used by
     int x;
                             // local class
     extern int g();
                             // extern function g can be used by
                              // local class
                             // local class
     class local
           int g() { return x; }
                                      // error, local variable x
                                      // cannot be used by g
                                     // valid,static variable y
           int h() { return y; }
           int k() { return y; } // valid, global x
           int 1() { return g(); } // valid, extern function g
     };
}
```

```
void main()
      local* z;
                               // error: the class local is not visible
      // ...}
```

Member functions of a local class have to be defined within their class definition. As a result, member functions of a local class are inline functions. Like all member functions, those defined within the scope of a local class do not need the keyword inline.

A local class cannot have static data members. In the following example, an attempt to define a static member of a local class causes an error:

```
void f()
   class local
      int f();
                            // error, local class has noninline
                            // member function
      int g() {return 0;} // valid, inline member function
                           // error, static is not allowed for
      static int a;
                           // local class
                            // valid, nonstatic variable
      int b;
   };
//
```

An enclosing function has no special access to members of the local class.

RELATED REFERENCES

- "Member Functions" on page 211
- "Inline Functions" on page 141

Local Type Names

C++ Local type names follow the same scope rules as other names. Type names defined within a class declaration have class scope and cannot be used outside their class without qualification.

If you use a class name, typedef name, or a constant name that is used in a type name, in a class declaration, you cannot redefine that name after it is used in the class declaration.

```
For example:
void main ()
{
      typedef double db;
      struct st
            typedef int db; // error
            db y;
      };
The following declarations are valid:
typedef float T;
class s {
      typedef int T;
      void f(const T);
};
```

Scope of Class Names

Here, function f() takes an argument of type s::T. However, the following declarations, where the order of the members of s has been reversed, cause an

```
typedef float T;
class s {
      void f(const T);
      typedef int T;
};
```

In a class declaration, you cannot redefine a name that is not a class name, or a typedef name to a class name or typedef name once you have used that name in the class declaration.

RELATED REFERENCES

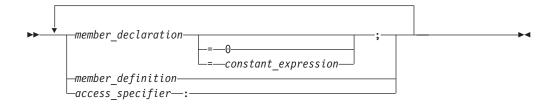
- "Scope" on page 1
- "Global Scope" on page 2
- "typedef" on page 29

Chapter 13. Class Members and Friends

Class Member Lists

An optional *member list* declares sub-objects called *class members*. Class members can be data, functions, nested types, and enumerators.

Syntax — Class Member List



The member list follows the class name and is placed between braces. The following applies to member lists, and members of member lists:

- A member_declaration or a member_definition may be a declaration or definition of a data member, member function, nested type, or enumeration. (The enumerators of a enumeration defined in a class member list are also members of the class.)
- A member list is the only place where you can declare class members.
- Friend declarations are not class members but must appear in member lists.
- The member list in a class definition declares all the members of a class; you cannot add members elsewhere.
- You cannot declare a member twice in a member list.
- You may declare a data member or member function as **static** but not **auto**, **extern**, or **register**.
- You may delcare a nested class, a member class template, or a member function, and then define it later outside the class.
- You must define static data members later outside the class.
- Nonstatic members that are class objects must be objects of previously defined classes; a class A cannot contain an object of class A, but it can contain a pointer or reference to an object of class A.
- You must specify all dimensions of a nonstatic array member.

A *constant initializer* (= *constant_expression*) may only appear in a class member of integral or enumeration type that has been declared as **static**.

A *pure specifier* (= 0) indicates that a function has no definition. It is only used with member functions declared as **virtual** and replaces the function definition of a member function in the member list.

An access specifier is one of **public**, **private**, or **protected**.

A *member declaration* declares a class member for the class containing the declaration.

Class Member Lists

EVAC+-- The order of allocation of nonstatic class members separated by an access_specifier is implementation dependent. The VisualAge C++ compiler allocates class members in the order that they are declared.

Suppose A is a name of a class. The following class members of A must have a name different from A:

- · All data members
- All type members
- All enumerators of enumerated type members
- · All members of all anonymous union members

RELATED REFERENCES

- "Chapter 3. Declarations" on page 23
- "Declaring Class Types" on page 199
- "Member Access" on page 223
- "Virtual Functions" on page 249
- "Static Members" on page 218

Data Members

C++ Data members include members that are declared with any of the fundamental types, as well as other types, including pointer, reference, array types, bit fields, and user-defined types. You can declare a data member the same way as a variable, except that explicit initializers are not allowed inside the class definition. However, a const static data member of integral or enumeration type may have an explicit initializer.

If an array is declared as a nonstatic class member, you must specify all of the dimensions of the array.

A class can have members that are of a class type or are pointers or references to a class type. Members that are of a class type must be of a class type that is previously declared. An incomplete class type can be used in a member declaration as long as the size of the class is not needed. For example, a member can be declared that is a pointer to an incomplete class type.

A class X cannot have a member that is of type X, but it can contain pointers to X, references to X, and static objects of X. Member functions of X can take arguments of type X and have a return type of X. For example:

```
class X
      X();
      X *xptr;
      X &xref;
      static X xcount;
      X xfunc(X);
};
```

RELATED REFERENCES

- "Pointers" on page 58
- "References" on page 68
- "Arrays" on page 62
- "Pointers" on page 58
- "References" on page 68
- "Incomplete Class Declarations" on page 204

Member Functions

Member functions are operators and functions that are declared as members of a class. Member functions do not include operators and functions declared with the **friend** specifier. These are called *friends* of a class. You can declare a member function as **static**; this is called a *static member function*. A member function that is not declared as **static** is called a *nonstatic member function*.

Suppose that you create an object named x of class A, and class A has a nonstatic member function f(). If you call the function x.f(), the keyword **this** in the body of f() is the address of x.

The definition of a member function is within the scope of its enclosing class. The body of a member function is analyzed after the class declaration so that members of that class can be used in the member function body, even if the member function definition appears before the declaration of that member in the class member list. When the function add() is called in the following example, the data variables a, b, and c can be used in the body of add().

Inline Member Functions

You may either define a member function inside its class definition, or you may it outside if you have alrealy declared (but not defined) the member function in the class definition.

A member function that is defined inside its class member list is called an *inline member function*. Member functions containing a few lines of code are usually declared inline. In the above example, add() is an inline member function. If you define a member function outside of its class definition, it must appear in a namespace scope enclosing the class definition. You must also qualify the member function name using the scope resolution (::) operator.

An equivalent way to declare an inline member function is to either declare it in the class with the **inline** keyword (and define the function outside of its class) or to define it outside of the class declaration using the **inline** keyword.

In the following example, member function Y::f() is an inline member function:

```
struct Y {
private:
   char a*;
public:
   char* f() { return a; }
};
```

The following example is equivalent to the previous example; Y::f() is an inline member function:

```
struct Y {
private:
   char a*;
```

Member Functions

```
public:
 char* f();
inline char* Z::f() { return a; }
```

Inline member functions have internal linkage. Noninline member functions have external linkage.

Member Functions of Local Classes

Member functions of a local class have to be defined within their class definition. As a result, member functions of a local class are inline functions. Like all member functions, those defined within the scope of a local class do not need the keyword inline.

RELATED REFERENCES

- "Friends" on page 225
- "Static Member Functions" on page 221
- "Chapter 7. Functions" on page 123
- "Inline Functions" on page 141
- "Local Classes" on page 206

const and volatile Member Functions

A member function declared with the **const** qualifier can be called for constant and nonconstant objects. A nonconstant member function can only be called for a nonconstant object. Similarly, a member function declared with the volatile qualifier can be called for volatile and nonvolatile objects. A nonvolatile member function can only be called for a nonvolatile object.

RELATED REFERENCES

• "volatile and const Qualifiers" on page 50

Virtual Member Functions

■ C++■ Virtual member functions are declared with the keyword virtual. They allow dynamic binding of member functions. Because all virtual functions must be member functions, virtual member functions are simply called virtual functions.

If the definition of a virtual function is replaced by a pure specifier in the declaration of the function, the function is said to be declared pure. A class that has at least one pure virtual function is called an abstract class.

RELATED REFERENCES

- "Virtual Functions" on page 249
- "Abstract Classes" on page 255

Special Member Functions

Special member functions are used to create, destroy, initialize, convert, and copy class objects. These include the following:

- Constructors
- Destructors
- Conversion constructors

- Conversion functions
- Copy constructors

RELATED REFERENCES

- "Constructors" on page 259
- "Destructors" on page 266
- "Conversion by Constructor" on page 276
- "Conversion Functions" on page 277
- "Copy Constructors" on page 278

Member Scope

C++ Member functions and static members can be defined outside their class declaration if they have already been declared, but not defined, in the class member list. Nonstatic data members are defined when an object of their class is created. The declaration of a static data member is not a definition. The declaration of a member function is a definition if the body of the function is also given.

Whenever the definition of a class member appears outside of the class declaration, the member name must be qualified by the class name using the :: (scope resolution) operator.

The following example defines a member function outside of its class declaration.

```
#include <iostream>
using namespace std;
struct X {
 int a, b;
 // member function declaration only
 int add();
// global variable
int a = 10;
// define member function outside its class declaration
int X::add() { return a + b; }
int main() {
 int answer;
 X xobject;
 xobject.a = 1;
 xobject.b = 2;
 answer = xobject.add();
  cout << xobject.a << " + " << xobject.b << " = " << answer << endl;</pre>
```

The output for this example is: 1 + 2 = 3

All member functions are in class scope even if they are defined outside their class declaration. In the above example, the member function add() returns the data member a, not the global variable a.

The name of a class member is local to its class. Unless you use one of the class access operators, . (dot), or -> (arrow), or :: (scope resolution) operator, you can

Member Scope

only use a class member in a member function of its class and in nested classes. You can only use types, enumerations and static members in a nested class without qualification with the :: operator.

The order of search for a name in a member function body is:

- 1. Within the member function body itself
- 2. Within all the enclosing classes, including inherited members of those classes
- 3. Within the lexical scope of the body declaration

The search of the enclosing classes, including inherited members, is demonstrated in the following example:

```
class A { /* ... */ };
class B { /* ... */ };
class C { /* ... */ };
class Z : A {
     class Y : B {
           class X : C { int f(); /* ... */ };
};
int Z::Y::X f()
      char j;
      return 0;
```

In this example, the search for the name j in the definition of the function f follows this order:

- 1. In the body of the function f
- 2. In X and in its base class C
- 3. In Y and in its base class B
- 4. In Z and in its base class A
- 5. In the lexical scope of the body of f. In this case, this is global scope.

Note that when the containing classes are being searched, only the definitions of the containing classes and their base classes are searched. The scope containing the base class definitions (global scope, in this example) is not searched.

RELATED REFERENCES

- "Class Member Lists" on page 209
- "C++ Scope Resolution Operator ::" on page 77
- "Class Scope" on page 3

Pointers to Members

Pointers to members allow you to refer to nonstatic members of class objects. You cannot use a pointer to member to point to a static class member because the address of a static member is not associated with any particular object. To point to a static class member, you must use a normal pointer.

You can use pointers to member functions in the same manner as pointers to functions. You can compare pointers to member functions, assign values to them, and use them to call member functions. Note that a member function does not have the same type as a nonmember function that has the same number and type of arguments and the same return type.

Pointers to members can be declared and used as shown in the following example:

```
#include <iostream>
using namespace std;
class X {
public:
 int a;
 void f(int b) {
   cout << "The value of b is "<< b << endl;</pre>
};
int main() {
  // declare pointer to data member
  int X::*ptiptr = &X::a;
  \ensuremath{//} declare a pointer to member function
  void (X::* ptfptr) (int) = &X::f;
  // create an object of class type X
  X xobject;
  // initialize data member
  xobject.*ptiptr = 10;
  cout << "The value of a is " << xobject.*ptiptr << endl;</pre>
  // call member function
  (xobject.*ptfptr) (20);
The output for this example is:
The value of a is 10
The value of b is 20
```

To reduce complex syntax, you can declare a **typedef** to be a pointer to a member. A pointer to a member can be declared and used as shown in the following code fragment:

```
typedef int X::*my_pointer_to_member;
typedef void (X::*my_pointer_to_function) (int);
int main() {
    my_pointer_to_member ptiptr = &X::a;
    my_pointer_to_function ptfptr = &X::f;
    X xobject;
    xobject.*ptiptr = 10;
    cout << "The value of a is " << xobject.*ptiptr << endl;
    (xobject.*ptfptr) (20);
}</pre>
```

The pointer to member operators .* and ->* are used to bind a pointer to a member of a specific class object. Because the precedence of () (function call operator) is higher than .* and ->*, you must use parentheses to call the function pointed to by ptf.

For more information, see "C++ Pointer to Member Operators .* ->*" on page 106.

RELATED REFERENCES

- "Static Members" on page 218
- "typedef" on page 29
- "C++ Pointer to Member Operators .* ->*" on page 106

The this Pointer

The keyword this identifies a special type of pointer. Suppose that you create an object named x of class A, and class A has a nonstatic member function f(). If you call the function x.f(), the keyword this in the body of f() is the address of x. You cannot declare the **this** pointer or make assignments to it.

A static member function does not have a **this** pointer.

The type of the **this** pointer for a member function of a class type X, is X* const. If the member function is declared with the **const** qualifier, the type of the **this** pointer for that member function for class X, is const X* const. If the member function is declared with the volatile qualifier, the type of the this pointer for that member function for class X is volatile X* const. For example, the compiler will not allow the following:

```
struct A {
 int a:
 int f() const { return a++; }
```

The compiler will not allow the statement a++ in the body of function f(). In the function f(), the this pointer is of type A* const. The function f() is trying to modify part of the object to which **this** points.

The this pointer is passed as a hidden argument to all nonstatic member function calls and is available as a local variable within the body of all nonstatic functions.

For example, you can refer to the particular class object that a member function is called for by using the this pointer in the body of the member function. The following code example produces the output a = 5:

```
#include <iostream>
using namespace std;
struct X {
private:
 int a:
public:
 void Set a(int a) {
    // The 'this' pointer is used to retrieve 'xobj.a'
    // hidden by the automatic variable 'a'
    this->a = a:
   void Print a() { cout << "a = " << a << endl; }</pre>
};
int main() {
 X xobj;
 int a = 5;
 xobj.Set_a(a);
 xobj.Print_a();
```

In the member function Set_a(), the statment this->a = a uses the this pointer to retrieve xobj.a hidden by the automatic variable a.

Unless a class member name is hidden, using the class member name is equivalent to using the class member name with the this pointer and the class member access operator (->).

The this Pointer

The example in the first column of the following table shows code that uses class members without the **this** pointer. The code in the second column uses the variable THIS to simulate the first column's hidden use of the **this** pointer:

```
Code without using this pointer
                                           Equivalent code, the THIS variable
                                           simulating the hidden use of the this pointer
#include <string>
                                           #include <string>
#include <iostream>
                                           #include <iostream>
using namespace std;
                                           using namespace std;
struct X {
                                           struct X {
private:
                                          private:
 int len:
                                            int len;
 char *ptr;
                                            char *ptr;
public:
                                           public:
                                            int GetLen (X* const THIS) {
 int GetLen() {
   return len;
                                              return THIS->len;
 char * GetPtr() {
                                            char * GetPtr (X* const THIS) {
   return ptr;
                                              return THIS->ptr;
 X& Set(char *);
                                            X& Set(X* const, char *);
 X& Cat(char *);
                                            X& Cat(X* const, char *);
 X& Copy(X&);
                                            X& Copy(X* const, X&);
 void Print();
                                            void Print(X* const);
X& X::Set(char *pc) {
                                          X& X::Set(X* const THIS, char *pc) {
                                            THIS->len = strlen(pc);
THIS->ptr = new char[THIS->len];
 len = strlen(pc);
 ptr = new char[len];
                                            strcpy(THIS->ptr, pc);
 strcpy(ptr, pc);
 return *this;
                                            return *THIS;
X& X::Cat(char *pc) {
                                          X& X::Cat(X* const THIS, char *pc) {
 len += strlen(pc);
                                            THIS->len += strlen(pc);
                                            strcat(THIS->ptr, pc);
 strcat(ptr,pc);
 return *this;
                                            return *THIS;
X& X::Copy(X& x) {
                                          X\& X::Copy(X* const THIS, X\& x) {
 Set(x.GetPtr());
                                            THIS->Set(THIS, x.GetPtr(&x));
 return *this;
                                            return *THIS;
void X::Print() {
                                           void X::Print(X* const THIS) {
 cout << ptr << endl;</pre>
                                            cout << THIS->ptr << endl;</pre>
int main() {
                                           int main() {
 X xobj1;
                                            X xobj1;
                                            xobj1.Set("abcd")
       .Cat("efgh");
 xobj1.Print();
                                            xobj1.Print(&xobj1);
 X xobj2;
                                            X xobj2;
                                            xobj2.Copy(xobj1)
       .Cat("ijkl");
  xobj2.Print();
                                            xobj2.Print(&xobj2);
```

Both examples produces the following output: abcdefgh abcdefghijkl

RELATED REFERENCES

• "volatile and const Qualifiers" on page 50

Static Members

Class members can be declared using the storage-class specifier static in the class member list. Only one copy of the static member is shared by all objects of a class in a program. When you declare an object of a class having a static member, the static member is not part of the class object.

A typical use of static members is for recording data common to all objects of a class. For example, you can use a static data member as a counter to store the number of objects of a particular class type that are created. Each time a new object is created, this static data member can be incremented to keep track of the total number of objects.

You access a static member by qualifying the class name using the :: (scope resolution) operator. In the following example, you can refer to the static member f() of class type X as X::f() even if no object of type X is ever declared:

```
struct X {
 static int f();
};
int main() {
 X::f();
```

For more information on the storage-class specifier static, see "static Storage Class Specifier" on page 28

RELATED REFERENCES

- "static Storage Class Specifier" on page 28
- "Class Member Lists" on page 209

Using the Class Access Operators with Static Members

Soundo not have to use the class member access syntax to refer to a static member; to access a static member s of class X, you could use the expression X::s. The following example demonstrates accessing a static member:

```
#include <iostream>
using namespace std;
struct A {
 static void f() { cout << "In static function A::f()" << endl; }</pre>
int main() {
 // no object required for static member
 A::f();
 A a;
 A* ap = &a;
 a.f();
 ap->f();
```

The three statements A::f(), a.f(), and ap->f() all call the same static member function A::f().

You can directly refer to a static member in the same scope of its class, or in the scope of a class derived from the static member's class. The following example demonstrates the latter case (directly referring to a static member in the scope of a class derived from the static member's class):

```
#include <iostream>
using namespace std;
int g() {
   cout << "In function g()" << endl;</pre>
   return 0;
class X {
   public:
      static int g() {
         cout << "In static member function X::g()" << endl;</pre>
         return 1;
};
class Y: public X {
   public:
      static int i;
};
int Y::i = g();
int main() { }
```

The following is the output of the above code:

In static member function X::g()

The initialization int Y::i = g() calls X::g(), not the function g() declared in the global namespace.

A static member can be referred to independently of any association with a class object because there is only one static member shared by all objects of a class. A static member will exist even if no objects of its class have been declared.

RELATED REFERENCES

- "static Storage Class Specifier" on page 28
- "C++ Scope Resolution Operator ::" on page 77
- "Dot Operator ." on page 81
- "Arrow Operator ->" on page 82

Static Data Members

Only one copy of a static data member of a class exists; it is shared with all objects of that class.

Static data members of a class in namespace scope have external linkage. Static data members follow the usual class access rules, except that they can be initialized in file scope. Static data members and their initializers can access other static private and protected members of their class. The initializer for a static data member is in the scope of the class declaring the member.

Static Members

A static data member can be of any type except for void or void qualified with const or volatile.

The declaration of a static data member in the member list of a class is not a definition. The definition of a static data member is equivalent to an external variable definition. You must define the static member outside of the class declaration in namespace scope.

For example:

```
class X
public:
     static int i;
int X::i = 0; // definition outside class declaration
```

Once you define a static data member, it exists even though no objects of the static data member's class exist. In the above example, no objects of class X exist even though the static data member X::i has been defined.

The following example shows how you can initialize static members using other static members, even though these members are private:

```
class C {
     static int i;
     static int j;
     static int k;
     static int 1;
     static int m;
     static int n;
     static int p;
     static int q;
     static int r;
     static int s;
     static int f() { return 0; }
public:
     C() \{ a = 0; \}
     };
С с;
int C::r = 1;
                  // initialize with a constant value
class Y : private C {} y;
int C::m = Y::f();
int C::n = Y::r;
                  // error
int C::p = y.r;
int C::q = y.f();
                 // error
```

The initializations of C::p and C::x cause errors because y is an object of a class that is derived privately from C, and its members are not accessible to members of C.

If a static data member is of const integral or const enumeration type, you may specify a constant initializer in the static data member's declaration. This constant initializer must be an integral constant expression. Note that the constant initializer is not a definition. You still need to define the static member in an enclosing namespace is still required. The following example demonstrates this:

```
#include <iostream>
using namespace std;

struct X {
   static const int a = 76;
};

const int X::a;

int main() {
   cout << X::a << endl;
}</pre>
```

The tokens = 76 at the end of the declaration of static data member a is a constant initializer.

You can only have one definition of a static member in a program. Unnamed classes and classes contained within unnamed classes cannot have static data members.

You cannot declare a static data member as mutable.

Local classes cannot have static data members.

RELATED REFERENCES

- "External Linkage" on page 5
- "Member Access" on page 223
- "Local Classes" on page 206

Static Member Functions

You cannot have static and nonstatic member functions with the same names and the same number and type of arguments.

Like static data members, you may access a static member function f() of a class A without using an object of class A.

A static member function does not have a **this** pointer. The following example demonstrates this:

```
#include <iosteam>
using namespace std;
struct X {
private:
 int i;
  static int si;
public:
 void set i(int arg) { i = arg; }
 static void set si(int arg) { si = arg; }
  void print i() {
    cout << "Value of i = " << i << endl;</pre>
    cout << "Again, value of i = " << this->i << endl;</pre>
  static void print_si() {
    cout << "Value of si = " << si << endl;
     cout << "Again, value of si = " << this->si << endl;</pre>
};
```

Static Members

```
// Initialize static data member
int X::si = 77;
int main() {
 X xobj;
 xobj.set i(11);
 xobj.print i();
 // static data members and functions belong to the class and
 // can be accessed without using an object of class X
 X::print si();
 X::set si(22);
 X::print_si();
```

The following is the output of the above example:

```
Value of i = 11
Again, value of i = 11
Value of si = 77
Value of si = 22
```

The compiler would not allow the member access operation this->si in function A::print si() because this member function has been declared as static, and therefore does not have a **this** pointer.

You can call a static member function using the this pointer of a nonstatic member function. In the following example, the nonstatic member function printall() calls the static member function f() using the **this** pointer:

```
#include <iostream>
using namespace std;
class C {
 static void f() {
   cout << "Here is i: " << i << endl;</pre>
 static int i;
 int j;
public:
 C(int firstj): j(firstj) { }
 void printall();
void C::printall() {
 cout << "Here is j: " << this->j << endl;</pre>
 this->f();
int C::i = 3;
int main() {
 C obj C(0);
 obj_C.printall();
```

The following is the output of the above example:

```
Here is j: 0
Here is i: 3
```

A static member function cannot be declared with the keywords virtual, const, volatile, orconst volatile.

A static member function can access only the names of static members, enumerators, and nested types of the class in which it is declared. Suppose a static member function f() is a member of class X. The static member function f() cannot access the nonstatic members X or the nonstatic members of a base class of X.

RELATED REFERENCES

• "The this Pointer" on page 216

Member Access

Member access determines if a class member is accessible in an expression or declaration. Suppose x is a member of class A. Class member x can be one of the following:

- **public**: x can be used anywhere without the access restrictions defined by private or protected.
- private: x can be used only by the members and friends of class A.
- **protected**: x can be used only by the members and friends of class A, and the members and friends of classes derived from class A.

Members of classes declared with the keyword **class** are private by default. Members of classes declared with the keyword **struct** or **union** are public by default.

To control the access of a class member, you use one of the *access specifiers* **public**, **private**, or **protected** as a label in a class member list. The following example demonstrates these access specifiers:

```
struct A {
  friend class C;
private:
 int a;
public:
 int b;
protected:
 int c;
};
struct B : A {
 void f() {
   // a = 1;
   b = 2;
    c = 3;
};
struct C {
  void f(A x) {
   x.a = 4;
   x.b = 5;
   x.c = 6;
};
int main() {
 Ау;
// y.a = 7;
 y.b = 8;
// y.c = 9;
  Bz;
```

```
// z.a = 10;
 z.b = 11;
// z.c = 12;
```

The following table lists the access of data members A::a A::b, and A::c in various scopes of the above example:

Scope	A::a	A::b	A::c
function B::f()	No access. Member A::a is private.	Access. Member A::b is public.	Access. Class B inherits from A.
function C::f()	Access. Class C is a friend of A.	Access. Member A::b is public.	Access. Class C is a friend of A.
object y in main()	No access. Member y.a is private.	Access. Member y.a is public.	No access. Member y.c is protected.
object z in main()	No access. Member z.a is pivate.	Access. Member z.a is public.	No access. Member z.c is protected.

An access specifier specifies the accesibility of members that follow it until the next access specifier or until the end of the class definition. You can use any number of access specifiers in any order. If you later define a class member within its class definition, its access specification must be the same as its declaration. The following example demonstrates this:

```
class A {
    class B;
 public:
    class B { };
};
```

The compiler will not allow the definition of class B because this class has already been declared as private.

A class member has the same access control regardless whether it has been defined within its class or outside its class.

Access control applies to names. In particular, if you add access control to a typedef name, it affects only the typedef name. The following example demonstrates this:

```
class A {
   class B { };
 public:
   typedef B C;
int main() {
 A::C x;
// A::B y;
```

The compiler will allow the declaration A::C x because the typedef name A::C is public. The compiler would not allow the declaration A::B y because A::B is private.

Note that accessibility and visibility are independent. Visibility is based on the scoping rules of C++. A class member can be visible and inaccessible at the same time.

RELATED REFERENCES

- "Scope" on page 1
- "Class Member Lists" on page 209
- "Inherited Member Access" on page 236

Friends

A friend of a class X is a function or class that is not a member of X, but is granted the same access to X as the members of X. Functions declared with the **friend** specifier in a class member list are called *friend functions* of that class. Classes declared with the **friend** specifier in the member list of another class are called *friend classes* of that class.

A class Y must be defined before any member of Y can be declared a friend of another class.

In the following example, the friend function print is a member of class Y and accesses the private data members a and b of class X.

```
#include <iostream>
using namespace std;
class X;
class Y {
public:
 void print(X& x);
};
class X {
 int a, b;
 friend void Y::print(X& x);
public:
 X() : a(1), b(2) { }
void Y::print(X& x) {
  cout << "a is " << x.a << endl:</pre>
  cout << "b is " << x.b << endl;
int main() {
 X xob.j;
  Y yobj;
  yobj.print(xobj);
```

The following is the output of the above example:

```
a is 1 b is 2
```

You can declare an entire class as a friend. Suppose class F is a friend of class A. This means that every member function and static data member definition of class F has access to class A.

In the following example, the friend class F has a member function print that accesses the private data members a and b of class X and performs the same task as the friend function print in the above example. Any other members declared in class F also have access to all members of class X:

Friends

```
#include <iostream>
using namespace std;
class X {
  int a, b;
  friend class F;
public:
 X() : a(1), b(2) \{ \}
class F {
public:
 void print(X& x) {
  cout << "a is " << x.a << endl;</pre>
    cout << "b is " << x.b << endl;
};
int main() {
 X xobj;
  F fobj;
  fobj.print(xobj);
```

The following is the output of the above example:

```
b is 2
```

You must use an elaborated type specifier when you declare a class as a friend. The following example demonstrates this:

```
class F;
class G;
class X {
  friend class F;
  friend G;
};
```

VAC++ The VisualAge C++ compiler will warn you that the friend declaration of G must be an elaborated class name.

You cannot define a class in a friend declaration. For example, the compiler will not allow the following:

```
class F;
class X {
  friend class F { };
```

However, you can define a function in a friend declaration. The class must be a non-local class, function, the function name must be unqualified, and the function has namespace scope. The following example demonstrates this:

```
class A {
 void g();
};
void z() {
 class B {
// friend void f() { };
 };
```

```
class C {
// friend void A::g() { }
  friend void h() { }
};
```

The compiler would not allow the function definition of f() or g(). The compiler will allow the definition of h().

You cannot declare a friend with a storage class specifier.

RELATED REFERENCES

• "Member Access" on page 223

Friend Scope

The name of a friend function or class first introduced in a friend declaration is not in the scope of the class granting friendship (also called the *enclosing class*) and is not a member of the class granting friendship.

The name of a function first introduced in a friend declaration is in the scope of the first nonclass scope that contains the enclosing class. The body of a function provided in a friend declaration is handled in the same way as a member function defined within a class. Processing of the definition does not start until the end of the outermost enclosing class. In addition, unqualified names in the body of the function definition are searched for starting from the class containing the function definition.

A class that is first declared in a friend declaration is equivalent to an **extern** declaration. For example:

```
class B {};
class A
{
     friend class B; // global class B is a friend of A
};
```

If the name of a friend class has been introduced before the friend declaration, the compiler searches for a class name that matches the name of the friend class beginning at the scope of the friend declaration. If the declaration of a nested class is followed by the declaration of a friend class with the same name, the nested class is a friend of the enclosing class.

The scope of a friend class name is the first nonclass enclosing scope. For example:

```
class A {
    class B { // arbitrary nested class definitions
        friend class C;
    };
};

is equivalent to:
class C;
class A {
    class B { // arbitrary nested class definitions
        friend class C;
    };
};
```

If the friend function is a member of another class, you need to use the scope resolution operator (::). For example:

Friends

```
class A {
public:
 int f() { }
class B {
  friend int A::f();
```

Friends of a base class are not inherited by any classes derived from that base class. The following example demonstrates this:

```
class A {
 friend class B;
 int a;
};
class B { };
class C : public B {
 void f(A* p) {
// p->a = 2;
};
```

The compiler would not allow the statement p->a = 2 because class C is not a friend of class A, although C inherits from a friend of A.

Friendship is not transitive. The following example demonstrates this:

```
friend class B;
  int a;
};
class B {
 friend class C;
class C {
  void f(A* p) {
      p->a = 2;
 }
};
```

The compiler would not allow the statement p-a = 2 because class C is not a friend of class A, although C is a friend of a friend of A.

If you declare a friend in a local class, and the friend's name is unqualified, the compiler will look for the name only within the innermost enclosing nonclass scope. You must declare a function before declaring it as a friend of a local scope. You do not have to do so with classes. However, a declaration of a friend class will hide a class in an enclosing scope with the same name. The following example demonstrates this:

```
class X { };
void a();
void f() {
 class Y { };
 void b();
 class A {
    friend class X;
    friend class Y;
    friend class Z;
```

```
// friend void a();
   friend void b();
// friend void c();
};
::X moocow;
// X moocow2;
}
```

In the above example, the compiler will allow the following statements:

- friend class X: This statement does not declare :: X as a friend of A, but the local class X as a friend, even though this class is not otherwise declared.
- friend class Y: Local class Y has been declared in the scope of f().
- friend class Z: This statement declares the local class Z as a friend of A even though Z is not otherwise declared.
- friend void b(): Function b() has been declared in the scope of f().
- ::X moocow: This declaration creates an object of the nonlocal class ::X.

The compiler would not allow the following statements:

- friend void a(): This statement does not consider function a() declared in namespace scope. Since function a() has not been declared in the scope of f(), the compiler would not allow this statement.
- friend void c(): Since function c() has not been declared in the scope of f(), the compiler would not allow this statement.
- X moocow2: This declaration tries to create an object of the local class X, not the nonlocal class::X. Since local class X has not been defined, the compiler would not allow this statement.

RELATED REFERENCES

- "Scope of Class Names" on page 203
- "Nested Classes" on page 204
- "Dot Operator ." on page 81
- "Derivation" on page 233
- "External Linkage" on page 5

Friend Access

A friend of a class can access the private and protected members of that class. Normally, you can only access the private members of a class through member functions of that class, and you can only access the protected members of a class through member functions of a class or classes derived from that class.

Friend declarations are not affected by access specifiers.

RELATED REFERENCES

- "Friends" on page 225
- "Member Access" on page 223

Friends

Chapter 14. Inheritance

Inheritance is a mechanism of reusing and extending existing classes without modifying them.

Inheritance is almost like embedding an object into a class. Suppose that you declare an object x of class A in the class definition of B. As a result, class B will have access to all the public data members and member functions of class A. However, in class B, you have to access the data members and member functions of class A through object x. The following example demonstrates this:

```
#include <iostream>
using namespace std;

class A {
    int data;
public:
    void f(int arg) { data = arg; }
    int g() { return data; }
};

class B {
public:
    A x;
};

int main() {
    B obj;
    obj.x.f(20);
    cout << obj.x.g() << endl;
}// cout << obj.g() << endl;
}</pre>
```

In the main function, object obj accesses function A::f() through its data member B::x with the statement obj.x.f(20). Object obj accesses A::g() in a similar manner with the statement obj.x.g(). The compiler would not allow the statement obj.g() because g() is a member function of class A, not class B.

The inheritance mechanism lets you use a statement like obj.g() in the above example. In order for that statement to be legal, g() must be a member function of class B.

Inheritance lets you include the names and definitions of another class's members as part of a new class. The class whose members you want to include in your new class is called a *base class*. Your new class is *derived* from the base class. You new class will contain a *subobject* of the type of the base class. The following example is the same as the previous example except it uses the inheritance mechanism to give class B access to the members of class A:

```
#include <iostream>
using namespace std;

class A {
   int data;
public:
   void f(int arg) { data = arg; }
   int g() { return data; }
};
```

```
class B : public A { };
int main() {
   B obj;
   obj.f(20);
   cout << obj.g() << endl;</pre>
```

Class A is a base class of class B. The names and definitions of the members of class A are included in the definition of class B; class B inherits the members of class A. Class B is derived from class A. Class B contains a subobject of type A.

You can also add new data members and member functions to the derived class. You can modify the implementation of existing member functions or data by overriding base class member functions or data in the newly derived class.

You may derive classes from other derived classes, thereby creating another level of inheritance. The following example demonstrates this:

```
struct A { };
struct B : A {
struct C : B { };
```

Class B is a derived class of A, but is also a base class of C. The number of levels of inheritance is only limited by resources.

Multiple inheritance allows you to create a derived class that inherits properties from more than one base class. Because a derived class inherits members from all its base classes, ambiguities can result. For example, if two base classes have a member with the same name, the derived class cannot implicitly differentiate between the two members. Note that, when you are using multiple inheritance, the access to names of base classes may be ambiguous.

A direct base class is a base class that appears directly as a base specifier in the declaration of its derived class.

An indirect base class is a base class that does not appear directly in the declaration of the derived class but is available to the derived class through one of its base classes. For a given class, all base classes that are not direct base classes are indirect base classes. The following example demonstrates direct and indirect base classes:

```
class A {
 public:
   int x;
class B : public A {
 public:
   int y;
class C : public B { };
```

Class B is a direct base class of C. Class A is a direct base class of B. Class A is an indirect base class of C. (Class C has x and y as its data members.)

Polymorphic functions are functions that can be applied to objects of more than one type. In C++, polymorphic functions are implemented in two ways:

Overloaded functions are statically bound at compile time.

• C++ provides virtual functions. A virtual function is a function that can be called for a number of different user-defined types that are related through derivation. Virtual functions are bound dynamically at run time.

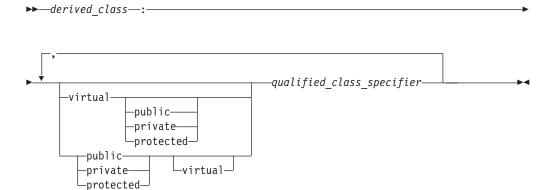
RELATED REFERENCES

- "Multiple Inheritance" on page 243
- "Overloading Functions" on page 185
- "Virtual Functions" on page 249
- "Chapter 7. Functions" on page 123
- "Chapter 12. Classes" on page 199
- "Chapter 13. Class Members and Friends" on page 209

Derivation

C++ Inheritance is implemented in C++ through the mechanism of derivation. Derivation allows you to derive a class, called a derived class, from another class, called a base class.

Syntax — Derived Class Derivation



In the declaration of a derived class, you list the base classes of the derived class. The derived class inherits its members from these base classes.

The qualified_class_specifier must be a class that has been previously declared in a class declaration.

An access specifier is one of **public**, **private**, or **protected**.

The **virtual** keyword can be used to declare virtual base classes.

The following example shows the declaration of the derived class D and the base classes V, B1, and B2. The class B1 is both a base class and a derived class because it is derived from class V and is a base class for D:

```
class V { /* ... */ };
class B1 : virtual public V { /* ... */ };
class B2 { /* ... */ };
class D : public B1, private B2 { /* ... */ };
```

Classes that are declared but not defined are not allowed in base lists.

For example:

Derivation

```
class X;
// error
class Y: public X { };
```

The compiler will not allow the declaration of class Y because X has not been defined.

When you derive a class, the derived class inherits class members of the base class. You can refer to inherited members (base class members) as if they were members of the derived class.

```
For example:
class Base {
public:
 int a,b;
class Derived : public Base {
public:
 int c;
int main() {
 Derived d;
 d.a = 1; // Base::a
 d.b = 2;
           // Base::b
 d.c = 3;  // Derived::c
```

The derived class can also add new class members and redefine existing base class members. In the above example, the two inherited members, a and b, of the derived class d, in addition to the derived class member c, are assigned values. If you redefine base class members in the derived class, you can still refer to the base class members by using the :: (scope resolution) operator.

For example:

```
#include <iostream>
using namespace std;
class Base {
public:
  char* name;
  void display() {
   cout << name << endl;</pre>
};
class Derived: public Base {
public:
  char* name;
  void display() {
    cout << name << ", " << Base::name << endl;</pre>
};
int main() {
  Derived d;
  d.name = "Derived Class";
  d.Base::name = "Base Class";
  // call Derived::display()
  d.display();
```

```
// call Base::display()
  d.Base::display();
}
```

The following is the output of the above example:

```
Derived Class, Base Class
Base Class
```

You can manipulate a derived class object as if it were a base class object. You can use a pointer or a reference to a derived class object in place of a pointer or reference to its base class. For example, you can pass a pointer or reference to a derived class object D to a function expecting a pointer or reference to the base class of D. You do not need to use an explicit cast to achieve this; a standard conversion is performed. You can implicitly convert a pointer to a derived class to point to an accessible unambiguous base class. You can also implicitly convert a reference to a derived class to a reference to a base class.

The following example demonstrates a standard conversion from a pointer to a derived class to a pointer to a base class:

```
#include <iostream>
using namespace std;
class Base {
public:
  char* name;
  void display() {
    cout << name << endl;</pre>
};
class Derived: public Base {
public:
 char* name;
 void display() {
    cout << name << ", " << Base::name << endl;</pre>
};
int main() {
  Derived d;
  d.name = "Derived Class";
  d.Base::name = "Base Class";
  Derived* dptr = &d;
  // standard conversion from Derived* to Base*
  Base* bptr = dptr;
  // call Base::display()
  bptr->display();
```

The following is the output of the above example: Base Class

The statement Base* bptr = dptr converts a pointer of type Derived to a pointer of type Base.

The reverse case is not allowed. You cannot implicitly convert a pointer or a reference to a base class object to a pointer or reference to a derived class. For

Derivation

example, the compiler will not allow the following code if the classes Base and Class are defined as in the above example:

```
int main() {
 Base b;
 b.name = "Base class";
 Derived* dptr = &b;
```

The compiler will not allow the statement Derived* dptr = &b because the statement is trying to implicitly convert a pointer of type Base to a pointer of type Derived.

If a member of a derived class and a member of a base class have the same name, the base class member is hidden in the derived class. If a member of a derived class has the same name as a base class, the base class name is hidden in the derived class.

RELATED REFERENCES

- "Virtual Base Classes" on page 244
- "Incomplete Class Declarations" on page 204
- "C++ Scope Resolution Operator ::" on page 77
- "Member Access" on page 223
- "References" on page 68

Inherited Member Access

Protected Members

A protected nonstatic base class member can be accessed by members and friends of any classes derived from that base class by using one of the following:

- · A pointer to a directly or indirectly derived class
- A reference to a directly or indirectly derived class
- An object of a directly or indirectly derived class

If a class is derived privately from a base class, all protected base class members become private members of the derived class.

If you reference a protected nonstatic member x of a base class A in a friend or a member function of a derived class B, you must access x through a pointer to, reference to, or object of a class derived from A. However, if you are accessing x to create a pointer to member, you must qualify x with a nested name specifier that names the derived class B. The following example demonstrates this:

```
class A {
public:
protected:
 int i;
class B : public A {
 friend void f(A*, B*);
 void g(A*);
void f(A* pa, B* pb) {
// pa->i = 1;
 pb->i = 2;
```

```
// int A::* point i = &A::i;
 int A::* point i2 = &B::i;
void B::g(A* pa) {
// pa->i = 1;
 i = 2;
// int A::* point_i = &A::i;
 int A::* point i2 = &B::i;
void h(A* pa, B* pb) {
// pa->i = 1;
// pb->i = 2;
int main() { }
```

Class A contains one protected data member, an integer i. Because B derives from A, the members of B have access to the protected member of A. Function f() is a friend of class B:

- The compiler would not allow pa->i = 1 because pa is not a pointer to the derived class B.
- The compiler would not allow int A::* point i = &A::i because i has not been qualified with the name of the derived class B.

Function g() is a member function of class B. The previous list of remarks about which statements the compiler would and would not allow apply for g() except for the following:

• The compiler allows i = 2 because it is equivalent to this->i = 2.

Function h() cannot access any of the protected members of A because h() is neither a friend or a member of a derived class of A.

RELATED REFERENCES

- "References" on page 68
- "Objects" on page 24

Access Control of Base Class Members

■ C++■ When you declare a derived class, an access specifier can precede each base class in the base list of the derived class. This does not alter the access attributes of the individual members of a base class as seen by the base class, but allows the derived class to restrict the access control of the members of a base class.

You can derive classes using any of the three access specifiers:

- In a public base class, public and protected members of the base class remain public and protected members of the derived class.
- In a protected base class, public and protected members of the base class are protected members of the derived class.
- In a private base class, public and protected members of the base class become private members of the derived class.

In all cases, private members of the base class remain private. Private members of the base class cannot be used by the derived class unless friend declarations within the base class explicitly grant access to them.

Inherited Member Access

In the following example, class d is derived publicly from class b. Class b is declared a public base class by this declaration.

```
class b { };
class d : public b // public derivation
```

You can use both a structure and a class as base classes in the base list of a derived class declaration:

- If the derived class is declared with the keyword class, the default access specifier in its base list specifiers is private.
- If the derived class is declared with the keyword **struct**, the default access specifier in its base list specifiers is **public**.

In the following example, private derivation is used by default because no access specifier is used in the base list and the derived class is declared with the keyword class:

```
struct B
{ };
class D : B // private derivation
{ };
```

Members and friends of a class can implicitly convert a pointer to an object of that class to a pointer to either:

- A direct private base class
- A protected base class (either direct or indirect)

RELATED REFERENCES

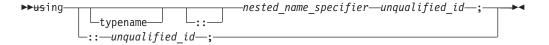
- "Member Access" on page 223
- "Structures" on page 36
- "Chapter 12. Classes" on page 199
- "Chapter 13. Class Members and Friends" on page 209
- "Friends" on page 225

The using Declaration and Class Members

■ C++■ A using declaration in a definition of a class A allows you to introduce a name of a data member or member function from a base class of A into the scope of A.

You would need a using declaration in a class definition if you want to create a set of overload a member functions from base and derived classes, or you want to change the access of a class member.

Syntax — using Declaration



A using declaration in a class A may name one of the following:

- · A member of a base class of A
- A member of an anonymous union that is a member of a base class of A
- An enumerator for an enumeration type that is a member of a base class of A

The following example demonstrates this:

```
struct Z {
   int g();
};

struct A {
   void f();
   enum E { e };
   union { int u; };
};

struct B : A {
   using A::f;
   using A::e;
   using A::u;
// using Z::g;
};
```

The compiler would not allow the using declaration using Z::g because Z is not a base class of A.

A using declaration cannot name a template. For example, the compiler will not allow the following:

```
struct A {
  template<class T> void f(T);
};

struct B : A {
  using A::f<int>;
};
```

Every instance of the name mentioned in a using declaration must be accessible. The following example demonstrates this:

```
struct A {
private:
   void f(int);
public:
   int f();
protected:
   void g();
};
struct B : A {
// using A::f;
   using A::g;
};
```

The compiler would not allow the using declaration using A::f because void A::f(int) is not accessible from B even though int A::f() is accessible.

RELATED REFERENCES

- "Scope of Class Names" on page 203
- "Overloading Member Functions from Base and Derived Classes" on page 240
- "Changing the Access of a Class Member" on page 241
- "The using Declaration and Namespaces" on page 183

Overloading Member Functions from Base and Derived Classes

A member function named f in a class A will hide all other members named f in the base classes of A, regardless of return types or arguments. The following example demonstrates this:

```
struct A {
 void f() { }
struct B : A {
 void f(int) { }
int main() {
 B obj B;
 obj_B.f(3);
// obj_B.f();
```

The compiler would not allow the function call obj_B.f() because the declaration of void B::f(int) has hidden A::f().

To overload, rather than hide, a function of a base class A in a derived class B, you introduce the name of the function into the scope of B with a using declaration. The following example is the same as the previous example except for the using declaration using A::f:

```
struct A {
 void f() { }
struct B : A {
 using A::f;
  void f(int) { }
int main() {
 B obj_B;
 obj_B_{\cdot}f(3);
  obj B.f();
```

Because of the using declaration in class B, the name f is overloaded with two functions. The compiler will now allow the function call obj B.f().

You can overload virtual functions in the same way. The following example demonstrates this:

```
#include <iostream>
using namespace std;
struct A {
  virtual void f() { cout << "void A::f()" << endl; }</pre>
  virtual void f(int) { cout << "void A::f(int)" << endl; }</pre>
};
struct B : A {
  using A::f;
  void f(int) { cout << "void B::f(int)" << endl; }</pre>
int main() {
  B obj B;
```

```
B* pb = &obj_B;
pb->f(3);
pb->f();
}
```

The following is the output of the above example:

```
void B::f(int)
void A::f()
```

Suppose that you introduce a function f from a base class A a derived class B with a using declaration, and there exists a function named B::f that has the same parameter types as A::f. Function B::f will hide, rather than conflict with, function A::f. The following example demonstrates this:

```
#include <iostream>
using namespace std;

struct A {
   void f() { }
   void f(int) { cout << "void A::f(int)" << endl; }
};

struct B : A {
   using A::f;
   void f(int) { cout << "void B::f(int)" << endl; }
};

int main() {
   B obj_B;
   obj_B.f(3);
}</pre>
```

The following is the output of the above example:

```
void B::f(int)
```

RELATED REFERENCES

- "Chapter 11. Overloading" on page 185
- "Name Hiding" on page 3
- "The using Declaration and Class Members" on page 238

Changing the Access of a Class Member

Suppose class B is a direct base class of class A. To restrict access of class B to the members of class A, derive B from A using either the access specifiers **protected** or **private**.

To increase the access of a member x of class A inherited from class B, use a using declaration. You cannot restrict the access to x with a using declaration. You may increase the access of the following members:

- A member inherited as private. (You cannot increase the access of a member declared as private because a using declaration must have access to the member's name.)
- · A member either inherited or declared as protected

The following example demonstrates this:

```
struct A {
protected:
  int y;
public:
  int z;
};
```

Inherited Member Access

```
struct B : private A { };
struct C : private A {
public:
 using A::y;
  using A::z;
struct D : private A {
protected:
 using A::y;
 using A::z;
struct E : D {
  void f() {
   y = 1;
   z = 2;
struct F : A {
public:
 using A::y;
private:
 using A::z;
int main() {
 B obj_B;
// obj_B.y = 3;
// obj_B.z = 4;
  C obj C;
  obj_C.y = 5;
 obj_C.z = 6;
 D obj D;
// obj D.y = 7;
// obj_D.z = 8;
  F obj F;
  obj_F.y = 9;
 obj F.z = 10;
```

The compiler would not allow the following assignments from the above example:

- obj B.y = 3 and obj B.z = 4: Members y and z have been inherited as private.
- obj_D.y = 7 and obj_D.z = 8: Members y and z have been inherited as **private**, but their access have been changed to protected.

The compiler allows the following statements from the above example:

- y = 1 and z = 2 in D::f(): Members y and z have been inherited as private, but their access have been changed to protected.
- obj_C.y = 5 and obj_C.z = 6: Members y and z have been inherited as **private**, but their access have been changed to public.
- obj_F.y = 9: The access of member y has been changed from **protected** to
- obj_F.z = 10: The access of member z is still public. The private using declaration using A::z has no effect on the access of z.

RELATED REFERENCES

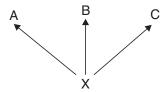
- "Member Access" on page 223
- "Protected Members" on page 236
- "Access Control of Base Class Members" on page 237
- "The using Declaration and Class Members" on page 238

You can derive a class from any number of base classes. Deriving a class from more than one direct base class is called *multiple inheritance*.

In the following example, classes A, B, and C are direct base classes for the derived class X:

```
class A { /* ... */ };
class B { /* ... */ };
class C { /* ... */ };
class X : public A, private B, public C { /* ... */ };
```

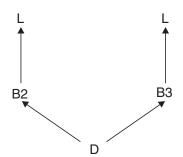
The following inheritance graph describes the inheritance relationships of the above example. An arrow points to the direct base class of the class at the tail of the arrow:



The order of derivation is relevant only to determine the order of default initialization by constructors and cleanup by destructors.

A direct base class cannot appear in the base list of a derived class more than once: class B1 { /* ... */ }; // direct base class class D : public B1, private B1 $\{ /* ... */ \}; // error$

However, a derived class can inherit an indirect base class more than once, as shown in the following example:



```
class L { /* \dots */ }; class B2 : public L { /* \dots */ };
                                               // indirect base class
class B3 : public L { /* ... */ };
class D : public B2, public B3 { /* ... */ }; // valid
```

In the above example, class D inherits the indirect base class L once through class B2 and once through class B3. However, this may lead to ambiguities because two

subobjects of class L exist, and both are accessible through class D. You can avoid this ambiguity by referring to class L using a qualified class name. For example: B2::L

or B3::L.

You can also avoid this ambiguity by using the base specifier virtual to declare a base class.

RELATED REFERENCES

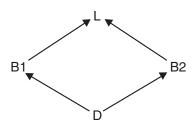
• "Virtual Base Classes"

Virtual Base Classes

EXAMPLE 2 Suppose you have two derived classes B and C that have a common base class A, and you also have another class D that inherits from B and C. You can declare the base class A as virtual to ensure that B and C share the same subobject of

In the following example, an object of class D has two distinct subobjects of class L, one through class B1 and another through class B2. You can use the keyword virtual in front of the base class specifiers in the base lists of classes B1 and B2 to indicate that only one sobobject of type L, shared by class B1 and class B2, exists.

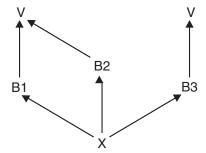
For example:



```
class L \{ /* ... */ \}; //  indirect base class
class B1 : virtual public L { /* ... */ };
class B2 : virtual public L { /* ... */ };
class D : public B1, public B2 { /* ... */ }; // valid
```

Using the keyword virtual in this example ensures that an object of class D inherits only one subobject of class L.

A derived class can have both virtual and nonvirtual base classes. For example:



```
class V { /* ... */ };
class B1 : virtual public V { /* ... */ };
class B2 : virtual public V { /* ... */ };
class B3 : public V { /* ... */ };
class D : public B1, public B2, public B3 { /* ... */ };
```

In the above example, class D has two subobjects of class V, one that is shared by classes B1 and B2 and one through class B3.

RELATED REFERENCES

• "Derivation" on page 233

Multiple Access

In an inheritance graph containing virtual base classes, a name that can be reached through more than one path is accessed through the path that gives the most access.

```
For example:
class L {
public:
    void f();
};
class B1 : private virtual L { };
class B2 : public virtual L { };
class D : public B1, public B2 {
public:
    void f() {
        // L::f() is accessed through B2
        // and is public
        L::f();
    }
};
```

In the above example, the function f() is accessed through class B2. Because class B2 is inherited publicly and class B1 is inherited privately, class B2 offers more access.

RELATED REFERENCES

- "Member Access" on page 223
- "Protected Members" on page 236
- "Access Control of Base Class Members" on page 237

Ambiguous Base Classes

When you derive classes, ambiguities can result if base and derived classes have members with the same names. Access to a base class member is ambiguous if you use a name or qualified name that does not refer to a unique function or object. The declaration of a member with an ambiguous name in a derived class is not an error. The ambiguity is only flagged as an error if you use the ambiguous member name.

For example, suppose that two classes named A and B both have a member named x, and a class named C inherits from both A and B. An attempt to access x from

class C would be ambiguous. You can resolve ambiguity by qualifying a member with its class name using the scope resolution (::) operator.

```
class B1 {
public:
 int i;
 int j;
 void g(int) { }
class B2 {
public:
 int j;
 void g() { }
class D : public B1, public B2 {
public:
 int i;
};
int main() {
 D dobj;
 D *dptr = &dobj;
 dptr->i = 5;
// dptr->j = 10;
 dptr->B1::j = 10;
// dobj.g();
 dobj.B2::g();
```

The statement dptr->j = 10 is ambiguous because the name j appears both in B1 and B2. The statement dobj.g() is ambiguous because the name g appears both in B1 and B2, even though B1::g(int) and B2::g() have different parameters.

The compiler checks for ambiguities at compile time. Because ambiguity checking occurs before access control or type checking, ambiguities may result even if only one of several members with the same name is accessible from the derived class.

Name Hiding

Suppose two subobjects named A and B both have a member name x. The member name x of subobject B hides the member name x of subobject A if A is a base class of B. The following example demonstrates this:

```
struct A {
   int x;
struct B: A {
  int x;
struct C: A, B {
   void f() \{ x = 0; \}
};
int main() {
  C i;
   i.f();
```

The assignment x = 0 in function C::f() is not ambiguous because the declaration B::x has hidden A::x. However, the compiler will warn you that deriving C from A is redundant because you already have access to the subobject A through B.

A base class declaration can be hidden along one path in the inheritance graph and not hidden along another path. The following example demonstrates this:

```
struct A { int x; };
struct B { int y; };
struct C: A, virtual B { };
struct D: A, virtual B {
   int x;
   int y;
};
struct E: C, D { };
int main() {
   E e;
// e.x = 1;
   e.y = 2;
};
```

The assignment e.x = 1 is ambiguous. The declaration D::x hides A::x along the path D::A::x, but it does not hide A::x along the path D::A::x. Therefore the variable x could refer to either D::x or A::x. The assignment e.y = 2 is not ambiguous. The declaration D::y hides B::y along both paths D::B::y and C::B::y because B is a virtual base class.

Ambiguity and using Declarations

Suppose you have a class named C that inherits from a class named A, and x is a member name of A. If you use a using declaration to declare A::x in C, then x is also a member of C; C::x does not hide A::x. Therefore using declarations cannot resolve ambiguities due to inherited members. The following example demonstrates this:

```
struct A {
   int x;
};

struct B: A { };

struct C: A {
   using A::x;
};

struct D: B, C {
   void f() { x = 0; }
};

int main() {
   D i;
   i.f();
}
```

The compiler will not allow the assignment x = 0 in function D::f() because it is ambiguous. The compiler can find x in two ways: as B::x or as C::x.

Unambiguous Class Members

The compiler can unambigously find static members, nested types, and enumerators defined in a base class A regardless of the number of subobjects of type A an object has. The following example demonstrates this:

```
struct A {
  int x;
  static int s;
  typedef A* Pointer_A;
  enum { e };
};
int A::s;
struct B: A { };
struct C: A { };
struct D: B, C {
  void f() {
     s = 1;
     Pointer_A pa;
     int i = e;
       x = 1;
//
};
int main() {
  Di;
  i.f();
```

The compiler allows the assignment s = 1, the declaration Pointer A pa, and the statement int i = e. There is only one static variable s, only one typedef Pointer A, and only one enumerator e. The compiler would not allow the assignment x = 1 because x can be reached either from from class B or class C.

Pointer Conversions

Conversions (either implicit or explicit) from a derived class pointer or reference to a base class pointer or reference must refer unambiguously to the same accessible base class object. (An accessible base class is a publicly derived base class that is neither hidden nor ambiguous in the inheritance hierarchy.) For example:

```
class W { /* ... */ };
class X : public W { /* ... */ };
class Y : public W { /* ... */ };
class Z : public X, public Y { /* ... */ };
void main ()
      Zz;
                            // valid
      X* xptr = &z;
      Y* yptr = &z;
                           // valid
      W* wptr = &z;
                             // error, ambiguous reference to class W
                             // X's W or Y's W ?
```

You can use virtual base classes to avoid ambiguous reference. For example:

```
class W { /* ... */ };
class X : public virtual W \{ /* ... */ \}; class Y : public virtual W \{ /* ... */ \};
class Z : public X, public Y \{ /* ... */ \};
void main ()
       Zz;
                          // valid
       X* xptr = &z;
       Y* yptr = &z;
                           // valid
                           // valid, W is virtual therefore only one
       W* wptr = &z;
                             // W subobject exists
```

Overload Resolution

Overload resolution takes place *after* the compiler unambigously finds a given function name. The following example demonstrates this:

```
struct A {
   int f() { return 1; }
};

struct B {
   int f(int arg) { return arg; }
};

struct C: A, B {
   int g() { return f(); }
};
```

The compiler will not allow the function call to f() in C::g() because the name f has been declared both in A and B. The compiler detects the ambiguity error before overload resolution can select the baset match A::f().

RELATED REFERENCES

- "C++ Scope Resolution Operator ::" on page 77
- "Virtual Base Classes" on page 244

Virtual Functions

By default, C++ matches a function call with the correct function definition at compile time. This is called *static binding*. You can specify that the compiler match a function call with the correct function definition at runtime; this is called *dynamic binding*. You declare a function with the keyword **virtual** if you want the compiler to use dynamic binding for that specific function.

The following examples demonstrate the differences between static and dynamic binding. The first example demonstrates static binding:

```
#include <iostream>
using namespace std;

struct A {
   void f() { cout << "Class A" << endl; }
};

struct B: A {
   void f() { cout << "Class B" << endl; }
};

void g(A& arg) {
   arg.f();
}

int main() {
   B x;
   g(x);
}</pre>
```

The following is the output of the above example: Class A

When function g() is called, function A::f() is called, although the argument refers to an object of type B. At compile time, the compiler knows only that the

Virtual Functions

argument of function g() will be a reference to an object derived from A; it cannot determine whether the argument will be a reference to an object of type A or type B. However, this can be determined at runtime. The following example is the same as the previous example, except that A::f() is declared with the virtual keyword:

```
#include <iostream>
using namespace std;
struct A {
  virtual void f() { cout << "Class A" << endl; }</pre>
struct B: A {
  void f() { cout << "Class B" << endl; }</pre>
void g(A& arg) {
   arg.f();
int main() {
   B x;
   g(x);
```

The following is the output of the above example: Class B

The virtual keyword indicates to the compiler that it should choose the appropriate definition of f() not by the type of reference, but by the type of object that the reference refers to.

Therefore, a virtual function is a member function you may redefine for other derived classes, and can ensure that the compiler will call the redefined virtual function for an object of the corresponding derived class, even if you call that function with a pointer or reference to a base class of the object.

A class that declares or inherits a virtual function is called a polymorphic class

You redefine a virtual member function, like any member function, in any derived class. Suppose you declare a virtual function named f in a class A, and you derive directly or indirectly from A a class named B. If you declare a function named f in class B with the same name and same parameter list as A::f, then B::f is also virtual (regardless whether or not you declare B::f with the **virtual** keyword) and it overrides A::f. However, if the parameter lists of A::f and B::f are different, A::f and B::f are considered different, B::f does not override A::f, and B::f is not virtual (unless you have declared it with the virtual keyword). Instead B::f hides A::f. The following example demonstrates this:

```
#include <iostream>
using namespace std;
  virtual void f() { cout << "Class A" << endl; }</pre>
struct B: A {
  void f(int) { cout << "Class B" << endl; }</pre>
struct C: B {
   void f() { cout << "Class C" << endl; }</pre>
```

```
int main() {
   B b; C c;
   A* pa1 = &b;
   A* pa2 = &c;
// b.f();
   pa1->f();
   pa2->f();
}
```

The following is the output of the above example:

```
Class A
Class C
```

The function B::f is not virtual. It hides A::f. Thus the compiler will not allow the function call b.f(). The function C::f is virtual; it overrides A::f even though A::f is not visible in C.

If you declare a base class destructor as virtual, a derived class destructor will override that base class destructor, even though destructors are not inherited.

The return type of an overriding virtual function may differ from the return type of the overridden virtual function. This overriding function would then be called a *covariant virtual function*. Suppose that B::f overrides the virtual function A::f. The return types of A::f and B::f may differ if all the following conditions are met:

- The function B::f returns a reference or pointer to a class of type T, and A::f
 returns a pointer or a reference to an unambigouous direct or indirect base class
 of T.
- The const or volatile qualification of the pointer or reference returned by B::f has the same or less const or volatile qualification of the pointer or reference returned by A::f.
- The return type of B::f must be complete at the point of declaration of B::f, or it can be of type B.

The following example demonstrates this:

#include <iostream>

```
using namespace std;
struct A { };
class B : private A {
   friend class D;
   friend class F;
};
A global A;
B global_B;
struct C {
   virtual A* f() {
      cout << "A* C::f()" << endl;</pre>
      return &global A;
};
struct D : C {
   B* f() {
      cout << "B* D::f()" << endl;</pre>
      return &global B;
};
```

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```
struct E:
struct F : C {
   Error:
   E is incomplete
    E* f();
struct G : C {
   Error:
   A is an inaccessible base class of B
//
    B* f();
};
int main() {
  D d;
  C* cp = &d;
  D* dp = &d;
  A* ap = cp->f();
  B* bp = dp->f();
};
```

The following is the output of the above example:

```
B* D::f()
B* D::f()
```

The statement A* ap = cp-f() calls D::f() and converts the pointer returned to type A*. The statement B* bp = dp->f() calls D::f() as well but does not convert the pointer returned; the type returned is B*. The compiler would not allow the declaration of the virtual function F::f() because E is not a complete class. The compiler would not allow the declaration of the virtual function G::f() because class A is not an accessible base class of B (unlike friend classes D and F, the definition of B does not give access to its members for class G).

A virtual function cannot be global or static because, by definition, a virtual function is a member function of a base class and relies on a specific object to determine which implementation of the function is called. You can declare a virtual function to be a friend of another class.

If a function is declared virtual in its base class, you can still access it directly using the scope resolution (::) operator. In this case, the virtual function call mechanism is suppressed and the function implementation defined in the base class is used. In addition, if you do not override a virtual member function in a derived class, a call to that function uses the function implementation defined in the base class.

A virtual function must be one of the following:

- Defined
- Declared pure
- · Defined and declared pure

A base class containing one or more pure virtual member functions is called an abstract class.

RELATED REFERENCES

- "Function Return Values" on page 140
- "Abstract Classes" on page 255

- "Friends" on page 225
- "C++ Scope Resolution Operator ::" on page 77

Ambiguous Virtual Function Calls

You cannot override one virtual function with two or more ambiguous virtual functions. This can happen in a derived class that inherits from two nonvirtual bases that are derived from a virtual base class.

```
For example:
class V {
public:
virtual void f() { }
class A : virtual public V {
  void f() { }
class B : virtual public V {
void f() { }
};
// Error:
// Both A::f() and B::f() try to override V::f()
class D : public A, public B { };
int main() {
  D d;
  V* vptr = &d;
 // which f(), A::f() or B::f()?
vptr->f();
```

The compiler will not allow the definition of class D. In class A, only A::f() will override V::f(). Similarly, in class B, only B::f() will override V::f(). However, in class D, both A::f() and B::f() will try to override V::f(). This attempt is not allowed because it is not possible to decide which function to call if a D object is referenced with a pointer to class V, as shown in the above example. Only one function can override a virtual function.

A special case occurs when the ambiguous overriding virtual functions come from separate instances of the same class type. In the following example, class D has two separate subobjects of class A:

```
#include <iostream>
using namespace std;

struct A {
    virtual void f() { cout << "A::f()" << endl; };
};

struct B : A {
    void f() { cout << "B::f()" << endl;};
};

struct C : A {
    void f() { cout << "C::f()" << endl;};
};

struct D : B, C { };</pre>
```

Virtual Functions

```
int main() {
  D d;
  B* bp = &d;
  A* ap = bp;
  D* dp = &d;
  ap->f();
// dp->f();
```

Class D has two occurences of class A, one inherited from B, and another inherited from C. Therefore there are also two occurences of the virtual function A::f. The statement ap->f() calls D::B::f. However the compiler would not allow the statement dp->f() because it could either call D::B::f or D::C::f.

Virtual Function Access

The access for a virtual function is specified when it is declared. The access rules for a virtual function are not affected by the access rules for the function that later overrides the virtual function. In general, the access of the overriding member function is not known.

If a virtual function is called with a pointer or reference to a class object, the type of the class object is not used to determine the access of the virtual function. Instead, the type of the pointer or reference to the class object is used.

In the following example, when the function f() is called using a pointer having type B*, bptr is used to determine the access to the function f(). Although the definition of f() defined in class D is exeuted, the access of the member function f() in class B is used. When the function f() is called using a pointer having type D*, dptr is used to determine the access to the function f(). This call produces an error because f() is declared private in class D.

```
class B {
public:
 virtual void f();
class D : public B {
private:
 void f();
};
int main() {
 D dobj;
 B* bptr = &dobj;
 D* dptr = &dobj;
 // valid, virtual B::f() is public,
  // D::f() is called
 bptr->f();
 // error, D::f() is private
 dptr->f();
```

Abstract Classes

An abstract class is a class that is designed to be specifically used as a base class. An abstract class contains at least one *pure virtual function*. You declare a pure virtual function by using a *pure specifier* (= 0) in the declaration of a virtual member function in the class declaration.

The following is an example of an abstract class::

```
class AB {
public:
   virtual void f() = 0;
};
```

Function AB::f is a pure virtual function. A function declaration cannot have both a pure specifier and a definition. For example, the compiler will not allow the following:

```
struct A {
   virtual void g() { } = 0;
};
```

You cannot use an abstract class as a parameter type, a function return type, or the type of an explicit conversion, nor can you declare an object of an abstract class. You can, however, declare pointers and references to an abstract class. The following example demonstrates this:

```
struct A {
 virtual void f() = 0;
struct B : A {
 virtual void f() { }
// Error:
// Class A is an abstract class
// A g();
// Error:
// Class A is an abstract class
// void h(A);
A& i(A&);
int main() {
// Error:
// Class A is an abstract class
// A a;
   A* pa;
  B b;
// Error:
// Class A is an abstract class
   static cast<A>(b);
```

Class A is an abstract class. The compiler would not allow the function declarations A g() or void h(A), declaration of object a, nor the static cast of b to type A.

Virtual member functions are inherited. A class derived from an abstract base class will also be abstract unless you override each pure virtual function in the derived class.

Abstract Classes

For example: class AB { public: virtual void f() = 0; class D2 : public AB { void g(); int main() { D2 d;

The compiler will not allow the declaration of object d because D2 is an abstract class; it inherited the pure virtual function f() from AB. The compiler will allow the declation of object d if you define function D2::g().

Note that you can derive an abstract class from a non-abstract class, and you can override a non-pure virtual function with a pure virtual function.

You can call member functions from a constructor or destructor of an abstract class. However, the results of calling (directly or indirectly) a pure virtual function from its constructor are undefined. The following example demonstrates this:

```
struct A {
 A() {
   direct();
   indirect();
 virtual void direct() = 0;
 virtual void indirect() { direct(); }
```

The default constructor of A calls the pure virtual function direct() both directly and indirectly (through indirect()).

The VisualAge C++ compiler issues a warning for the direct call to the pure virtual function. The compiler does not issue a warning for the indirect call to the pure virtual function.

Chapter 15. Special Member Functions

The default constructor, destructor, copy constructor, and copy assignment operator are *special member functions*. These functions create, destroy, convert, initialize, and copy class objects.

RELATED REFERENCES

- "Constructors" on page 259
- "Destructors" on page 266
- "Conversion by Constructor" on page 276
- "Conversion Functions" on page 277
- "Copy Constructors" on page 278

Constructors and Destructors Overview

Because classes have complicated internal structures, including data and functions, object initialization and cleanup for classes is much more complicated than it is for simple data structures. Constructors and destructors are special member functions of classes that are used to construct and destroy class objects. Construction may involve memory allocation and initialization for objects. Destruction may involve cleanup and deallocation of memory for objects.

Like other member functions, constructors and destructors are declared within a class declaration. They can be defined inline or external to the class declaration. Constructors can have default arguments. Unlike other member functions, constructors can have member initialization lists. The following restrictions apply to constructors and destructors:

- Constructors and destructors do not have return types nor can they return values.
- References and pointers cannot be used on constructors and destructors because their addresses cannot be taken.
- Constructors cannot be declared with the keyword virtual.
- Constructors and destructors cannot be declared static, const, or volatile.
- Unions cannot contain class objects that have constructors or destructors.

Constructors and destructors obey the same access rules as member functions. For example, if you declare a constructor with protected access, only derived classes and friends can use it to create class objects.

The compiler automatically calls constructors when defining class objects and calls destructors when class objects go out of scope. A constructor does not allocate memory for the class object its **this** pointer refers to, but may allocate storage for more objects than its class object refers to. If memory allocation is required for objects, constructors can explicitly call the **new** operator. During cleanup, a destructor may release objects allocated by the corresponding constructor. To release objects, use the **delete** operator.

Derived classes do not inherit constructors or destructors from their base classes, but they do call the constructor and destructor of base classes. Destructors can be declared with the keyword **virtual**.

Constructors are also called when local or temporary class objects are created, and destructors are called when local or temporary objects go out of scope.

You can call member functions from constructors or destructors. You can call a virtual function, either directly or indirectly, from a constructor or destructor of a class A. In this case, the function called is the one defined in A or a base class of A, but not a function overridden in any class derived from A. This avoids the possibility of accessing an unconstructed object from a constructor or destructor. The following example demonstrates this:

```
#include <iostream>
using namespace std;
struct A {
 virtual void f() { cout << "void A::f()" << endl; }</pre>
  virtual void g() { cout << "void A::g()" << endl; }</pre>
  virtual void h() { cout << "void A::h()" << endl; }</pre>
struct B : A {
  virtual void f() { cout << "void B::f()" << endl; }</pre>
    f();
    g();
    h();
  }
};
struct C : B {
  virtual void f() { cout << "void C::f()" << endl; }</pre>
  virtual void g() { cout << "void C::g()" << endl; }</pre>
  virtual void h() { cout << "void C::h()" << endl; }</pre>
};
int main() {
 C obj;
```

The following is the output of the above example:

void B::f() void A::g() void A::h()

The constructor of B does not call any of the functions overridden in C because C has been derived from B, although the example creates an object of type C named obj.

You can use the **typeid** or the **dynamic_cast** operator in constructors or destructors, as well as member initializers of constructors.

RELATED REFERENCES

- "volatile and const Qualifiers" on page 50
- "static Storage Class Specifier" on page 28
- "C++ new Operator" on page 92
- "C++ delete Operator" on page 95
- "Free Store" on page 269
- "The typeid operator" on page 111
- "dynamic_cast Operator" on page 85

Constructors

A *constructor* is a member function with the same name as its class. For example:

```
class X {
public:
   X();    // constructor for class X
};
```

Constructors are used to create, and can initialize, objects of their class type.

You cannot declare a constructor as **virtual** or **static**, nor can you declare a constructor as **const**, **volatile**, or **const volatile**.

You do not specify a return type for a constructor. A return statement in the body of a constructor cannot have a return value.

Default Constructors

A default constructor is a constructor that either has no parameters, or if it has parameters, all the parameters have default values.

If no user-defined constructor exists for a class A and one is needed, the compiler implicitly *declares* a constructor A::A(). This constructor is an inline public member of its class. The compiler will implicitly *define* A::A() when the compiler uses this constructor to create an object of type A. The constructor will have no constructor initializer and a null body.

The compiler first implictly defines the implicitly declared constructors of the base classes and nonstatic data members of a class A before defining the implicitly declared constructor of A. No default constructor is created for a class that has any constant or reference type members.

A constructor of a class A is *trivial* if all the following are true:

- It is implicitly defined
- · A has no virtual functions and no virtual base classes
- All the direct base classes of A have trivial constructors
- The classes of all the nonstatic data members of A have trivial constructors

If any of the above are false, then the constructor is *nontrivial*.

A union member cannot be of a class type that has a nontrivial constructor.

Like all functions, a constructor can have default arguments. They are used to initialize member objects. If default values are supplied, the trailing arguments can be omitted in the expression list of the constructor. Note that if a constructor has any arguments that do not have default values, it is not a default constructor.

A copy constructor for a class A is a constructor whose first parameter is of type A&, const A&, volatile A&, or const volatile A&. Copy constructors are used to make a copy of one class object from another class object of the same class type. You cannot use a copy constructor with an argument of the same type as its class; you must use a reference. You can provide copy constructors with additional parameters as long as they all have default arguments. If a user-defined copy constructor does not exist for a class and one is needed, the compiler implicitly

creates a copy constructor, with public access, for that class. A copy constructor is not created for a class if any of its members or base classes have an inaccessible copy constructor.

The following code fragment shows two classes with constructors, default constructors, and copy constructors:

```
class X {
public:
  // default constructor, no arguments
 X();
  // constructor
 X(int, int, int = 0);
  // copy constructor
 X(const X&);
  // error, incorrect argument type
};
class Y {
public:
  // default constructor with one
  // default argument
 Y(int = 0);
 // default argument
 // copy constructor
 Y(const Y\&, int = 0);
```

RELATED REFERENCES

• "Copy Constructors" on page 278

Explicit Initialization with Constructors

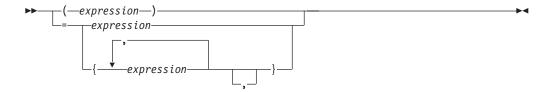
C++= A class object with a constructor must be explicitly initialized or have a default constructor. Except for aggregate initialization, explicit initialization using a constructor is the only way to initialize nonstatic constant and reference class members.

A class object that has no constructors, no virtual functions, no private or protected members, and no base classes is called an aggregate. Examples of aggregates are C-style structures and unions.

You explicitly initialize a class object when you create that object. There are two ways to initialize a class object:

- Using a parenthesized expression list. The compiler calls the constructor of the class using this list as the constructor's argument list.
- Using a single initialization value and the = operator. Because this type of expression is an initialization, not an assignment, the assignment operator function, if one exists, is not called. The type of the single argument must match the type of the first argument to the constructor. If the constructor has remaining arguments, these arguments must have default values.

The syntax for an initializer that explicitly initializes a class object with a constructor is:



The following example shows the declaration and use of several constructors that explicitly initialize class objects:

```
// This example illustrates explicit initialization
// by constructor.
#include <iostream>
using namespace std;
class complx {
  double re, im;
public:
  // default constructor
  complx() : re(0), im(0) { }
  // copy constructor
  complx(const complx& c) { re = c.re; im = c.im; }
  // constructor with default trailing argument
  complx( double r, double i = 0.0) { re = r; im = i; }
  void display() {
    cout << "re = "<< re << " im = " << im << endl;
};
int main() {
  // initialize with complx(double, double)
  complx one(1);
  // initialize with a copy of one
  // using complx::complx(const complx&)
  complx two = one;
  // construct complx(3,4)
  // directly into three
  complx three = complx(3,4);
  // initialize with default constructor
  complx four;
  // complx(double, double) and construct
  // directly into five
  complx five = 5;
  one.display();
  two.display();
  three.display();
  four.display();
  five.display();
```

The above example produces the following output:

```
re = 1 im = 0
re = 1 im = 0
re = 3 im = 4
re = 0 im = 0
re = 5 im = 0
```

RELATED REFERENCES

- "Structures" on page 36
- "Unions" on page 42

Initializing Base Classes and Members

Constructors can initialize their members in two different ways. A constructor can use the arguments passed to it to initialize member variables in the constructor definition:

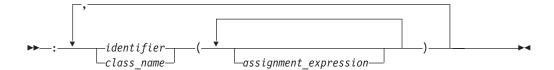
```
complx(double r, double i = 0.0) { re = r; im = i; }
```

Or a constructor can have an *initializer list* within the definition but prior to the function body:

```
complx(double r, double i = 0) : re(r), im(i) { /* ... */ }
```

Both methods assign the argument values to the appropriate data members of the class.

The syntax for a constructor initializer list is:



Include the initialization list as part of the function definition, not as part of the constructor declaration. For example:

```
#include <iostream>
using namespace std;
class B1 {
 int b;
public:
 B1() { cout << "B1::B1()" << endl; };
 // inline constructor
 B1(int i) : b(i) { cout << "B1::B1(int)" << endl; }
};
class B2 {
 int b;
protected:
 B2() { cout << "B1::B1()" << endl; }
  // noninline constructor
 B2(int i);
};
// B2 constructor definition including initialization list
B2::B2(int i) : b(i) { cout << "B2::B2(int)" << endl; }
```

```
class D : public B1, public B2 {
  int d1, d2;
public:
  D(int i, int j) : B1(i+1), B2(), d1(i) {
    cout << "D1::D1(int, int)" << end1;
    d2 = j;}
};
int main() {
  D obj(1, 2);
}</pre>
```

The following is the output of the above example:

```
B1::B1(int)
B1::B1()
D1::D1(int, int)
```

If you do not explicitly initialize a base class or member that has constructors by calling a constructor, the compiler automatically initializes the base class or member with a default constructor. In the above example, if you leave out the call B2() in the constructor of class D (as shown below), a constructor initializer with an empty expression list is automatically created to initialize B2. The constructors for class D, shown above and below, result in the same construction of an object of class D:

```
class D : public B1, public B2 {
  int d1, d2;
public:

  // call B2() generated by compiler
  D(int i, int j) : B1(i+1), d1(i) {
    cout << "D1::D1(int, int)" << endl;
    d2 = j;}
};</pre>
```

In the above example, the compiler will automatically call the default constructor for B2().

Note that you must declare constructors as public or protected to enable a derived class to call them. For example:

```
class B {
   B() { }
};

class D : public B {
   // error: implicit call to private B() not allowed
   D() { }
};
```

The compiler would not allow the definition of D::D() because this constructor cannot access the private constructor B::B().

You must initialize the following with an initializer list: base classes with no default constructors, reference data members, non-static const data members, or a class type which contains a constant data member. The following example demonstrates this:

```
class A {
public:
   A(int) { }
};
```

```
class B : public A {
 static const int i;
 const int j;
 int &k;
public:
 B(int\& arg) : A(0), j(1), k(arg) { }
int main() {
 int x = 0;
 B obj(x);
};
```

The data members j and k, as well as the base class A must be initialized in the initializer list of the constructor of B.

You can use data members when initializing members of a class. The following example demonstrate this:

```
struct A {
 int k;
 A(int i) : k(i) { }
struct B: A {
 int x;
 int i;
 int j;
 int& r;
 B(int i): r(x), A(i), j(this->i), i(i) { }
```

The constructor B(int i) initializes the following:

- B::x to refer to A::x
- Class A with the value of the argument to B(int i)
- B::j with the value of B::i
- B::i with the value of the argument to B(int i)

You can also call member functions (including virtual member functions) or use the operators typeid or dynamic_cast when initializing members of a class. However if you perform any of these operations in a member initialization list before all base classes have been initialized, the behavior is undefined. The following example demonstrates this:

```
#include <iostream>
using namespace std;
struct A {
 int i;
 A(int arg) : i(arg) {
   cout << "Value of i: " << i << endl;</pre>
};
struct B : A {
 int j;
 int f() { return i; }
 B();
};
B::B():A(f()), j(1234) {
   cout << "Value of j: " << j << endl;</pre>
```

```
int main() {
   B obj;
}
```

The output of the above example would be similar to the following:

```
Value of i: 8
Value of j: 1234
```

The behavior of the initializer A(f()) in the constructor of B is undefined. The runtime will call B::f() and try to access A::i even though the base A has not been initialized.

The following example is the same as the previous example except that the initializers of B::B() have different arguments:

```
#include <iostream>
using namespace std;
struct A {
 int i;
 A(int arg) : i(arg) {
   cout << "Value of i: " << i << endl;</pre>
};
struct B : A {
  int j;
  int f() { return i; }
 B();
};
B::B():A(5678), j(f()) {
   cout << "Value of j: " << j << endl;</pre>
int main() {
 B obj;
```

The following is the output of the above example:

```
Value of i: 5678
Value of j: 5678
```

The behavior of the initializer j(f()) in the constructor of B is well-defined. The base class A is already initialized when B::j is initialized.

RELATED REFERENCES

- "Default Constructors" on page 259
- "The typeid operator" on page 111
- "dynamic_cast Operator" on page 85

Construction Order of Derived Class Objects

When a derived class object is created using constructors, it is created in the following order:

- 1. Virtual base classes are initialized, in the order they appear in the base list.
- 2. Nonvirtual base classes are initialized, in declaration order.
- 3. Class members are initialized in declaration order (regardless of their order in the initialization list).

4. The body of the constructor is executed.

The following example demonstrates this:

```
#include <iostream>
using namespace std;
struct V {
 V() { cout << "V()" << endl; }</pre>
struct V2 {
 V2() { cout << "V2()" << end1; }</pre>
struct A {
 A() { cout << "A()" << endl; }
struct B : virtual V {
 B() { cout << "B()" << end1; }
struct C : B, virtual V2 {
 C() { cout << "C()" << end1; }
struct D : C, virtual V {
  A obj A;
 D() { cout << "D()" << endl; }
int main() {
 D c;
```

The following is the output of the above example:

V2() B() C() A() D()

> The above output lists the order in which the C++ runtime calls the constructors to create an object of type D.

RELATED REFERENCES

"Virtual Base Classes" on page 244

Destructors

Destructors are usually used to deallocate memory and do other cleanup for a class object and its class members when the object is destroyed. A destructor is called for a class object when that object passes out of scope or is explicitly deleted.

A destructor is a member function with the same name as its class prefixed by a (tilde). For example:

```
class X {
public:
  \//\ Constructor for class X
  X();
  // Destructor for class X
   X();
};
```

A destructor takes no arguments and has no return type. Its address cannot be taken. Destructors cannot be declared const, volatile, const volatile or static. A destructor can be declared virtual or pure virtual.

If no user-defined destructor exists for a class and one is needed, the compiler implicitly declares a destructor. This implicitly declared constructor is an inline public member of its class.

The compiler will implicitly define an implicitly declared destructor when the compiler uses the destructor to destroy an object of the destructor's class type. Suppose a class A has an implicitly declared destructor. The following is equivalent to the function the compiler would implicitly define for A:

```
~A::A() { }
```

The compiler first implictly defines the implicitly declared destructors of the base classes and nonstatic data members of a class A before defining the implicitly declared destructor of A

A destructor of a class A is *trivial* if all the following are true:

- It is implicitly defined
- All the direct base classes of A have trivial destructors
- The classes of all the nonstatic data members of A have trivial destructors

If any of the above are false, then the destructor is *nontrivial*.

A union member cannot be of a class type that has a nontrivial destructor.

Class members that are class types can have their own destructors. Both base and derived classes can have destructors, although destructors are not inherited. If a base class A or a member of A has a destructor, and a class derived from A does not declare a destructor, a default destructor is generated.

The default destructor calls the destructors of the base class and members of the derived class.

The destructors of base classes and members are called in the reverse order of the completion of their constructor:

- 1. The destructor for a class object is called before destructors for members and bases are called.
- 2. Destructors for nonstatic members are called before destructors for base classes are called.
- 3. Destructors for nonvirtual base classes are called before destructors for virtual base classes are called.

When an exception is thrown for a class object with a destructor, the destructor for the temporary object thrown is not called until control passes out of the catch block.

Destructors are implicitly called when an automatic object (a local object that has been declared auto or register, or not declared as static or extern) or temporary object passes out of scope. They are implicitly called at program termination for constructed external and static objects. Destructors are invoked when you use the **delete** operator for objects created with the **new** operator.

For example:

```
#include <string>
class Y {
private:
 char * string;
 int number;
public:
 // Constructor
 Y(const char*, int);
 // Destructor
  Y() { delete[] string; }
// Define class Y constructor
Y::Y(const char* n, int a) {
 string = strcpy(new char[strlen(n) + 1], n);
 number = a;
int main () {
 // Create and initialize
 // object of class Y
 Y yobj = Y("somestring", 10);
 // ...
  // Destructor Y is called before
 // control returns from main()
```

You can use a destructor explicitly to destroy objects, although this practice is not recommended. However to destroy an object created with the placement new operator, you can explicitly call the object's destructor. The following example demonstrates this:

```
#include <new>
#include <iostream>
using namespace std;
class A {
  public:
    A() { cout << "A::A()" << endl; }
     A() { cout << "A:: A()" << end1; }
int main () {
  char* p = new char[sizeof(A)];
 A* ap = new (p) A;
ap->A:: A();
  delete [] p;
```

The statement A* ap = new (p) A dynamically creates a new object of type A not in the free store but in the memory allocated by p. The statement delete [] p will delete the storage allocated by p, but the runtime will still believe that the object pointed to by ap still exists until you explicitly call the destructor of A (with the statement ap->A:: A()).

Non-class types have a pseudo destructor. The following example calls the pseudo destructor for an integer type:

```
typedef int I;
int main() {
 I x = 10;
 x.I:: I();
 x = 20;
```

The call to the pseudo destructor, x.I:: I(), has no effect at all. Object x has not been destroyed; the assignment x = 20 is still valid. Because pseudo destructors require the syntax for explicitly calling a destructor for a non-class type to be valid, you can write code without having to know whether or not a destructor exists for a given type.

RELATED REFERENCES

- "C++ new Operator" on page 92
- "C++ delete Operator" on page 95
- "Temporary Objects" on page 273

Free Store

Free store is a pool of memory available for you to allocate (and deallocate) storage for objects during the execution of your program. The **new** and **delete** operators are used to allocate and deallocate free store, respectively.

You can define your own versions of **new** and **delete** for a class by overloading them. You can declare the **new** and **delete** operators with additional parameters. When **new** and **delete** operate on class objects, the class member operator functions **new** and **delete** are called, if they have been declared.

If you create a class object with the **new** operator, one of the operator functions **operator new()** or **operator new(]()** (if they have been declared) is called to create the object. An **operator new()** or **operator new(]()** for a class is always a static class member, even if it is not declared with the keyword **static**. It has a return type **void*** and its first parameter must be the size of the object type and have type **std::size t**. It cannot be virtual.

Type **std::size_t** is an implementation-dependent unsigned integral type defined in the standard library header <cstddef>.

When you overload the **new** operator, you must declare it as a class member, returning type **void***, with its first parameter of type **std::size_t**, as described above. You can declare additional parameters in the declaration of **operator new()** or **operator new[]()**. Use the placement syntax to specify values for these parameters in an allocation expression.

The following example overloads two operator **new** functions:

X::operator new(size_t sz): This overloads the default **new** operator by allocating memory with the C function **malloc()**, and throwing a string (instead of **std::bad_alloc()** fails.

X::operator new(size_t sz, int location): This function takes an additional integer parameter, location. This function implements a very simplistic "memory manager" that manages the storage of up to three X objects.

Static array X::buffer holds three Node objects. Each Node object contains a pointer to an X object named data and a boolean variable named filled. Each X object stores an integer called number.

When you use this **new** operator, you pass the argument location which indicates the array location of buffer where you want to "create" your new X

object. If the array location is not "filled" (the data member of filled is equal to false at that array location), the new operator returns a pointer pointing to the X object located at buffer[location].

```
#include <new>
#include <iostream>
using namespace std;
class X;
struct Node {
  X* data;
  bool filled;
  Node() : filled(false) { }
class X {
  static Node buffer[];
public:
  int number;
  enum \{size = 3\};
  void* operator new(size t sz) throw (const char*) {
    void* p = malloc(sz);
    if (sz == 0) throw "Error: malloc() failed";
   cout << "X::operator new(size_t)" << endl;</pre>
    return p;
  void *operator new(size t sz, int location) throw (const char*) {
    cout << "X::operator new(size_t, " << location << ")" << endl;</pre>
    void* p = 0;
    if (location < 0 || location >= size || buffer[location].filled == true) {
      throw "Error: buffer location occupied";
    else {
   p = malloc(sizeof(X));
  if (p == 0) throw "Error: Creating X object failed";
  buffer[location].filled = true;
     buffer[location].data = (X*) p;
    return p;
  static void printbuffer() {
    for (int i = 0; i < size; i++) {
      cout << buffer[i].data->number << endl;</pre>
};
Node X::buffer[size];
int main() {
  try {
   X* ptr1 = new X;
   X* ptr2 = new(0) X;
    X* ptr3 = new(1) X;
    X* ptr4 = new(2) X;
    ptr2->number = 10000;
    ptr3->number = 10001;
    ptr4->number = 10002;
```

```
X::printbuffer();
  X* ptr5 = new(0) X;
}
catch (const char* message) {
  cout << message << endl;
}</pre>
```

The following is the output of the above example:

```
X::operator new(size_t)
X::operator new(size_t, 0)
X::operator new(size_t, 1)
X::operator new(size_t, 2)
10000
10001
10002
X::operator new(size_t, 0)
Error: buffer location occupied
```

The statement X* ptr1 = new X calls X::operator new(sizeof(X)). The statement X* ptr2 = new(0) X calls X::operator new(sizeof(X),0).

The **delete** operator destroys an object created by the **new** operator. The operand of **delete** must be a pointer returned by **new**. If **delete** is called for an object with a destructor, the destructor is invoked before the object is deallocated.

If you destroy a class object with the **delete** operator, the operator function **operator delete()** or **operator delete()** (if they have been declared) is called to destroy the object. An **operator delete()** or **operator delete()** for a class is always a static member, even if it is not declared with the keyword **static**. Its first parameter must have type **void***. Because **operator delete()** and **operator delete()** have a return type **void**, they cannot return a value.

The following example shows the declaration and use of the operator functions **operator new()** and **operator delete()**:

```
#include <cstdlib>
#include <iostream>
using namespace std;
class X {
public:
  void* operator new(size t sz) throw (const char*) {
    void* p = malloc(sz);
    if (p == 0) throw "malloc() failed";
    return p;
  // single argument
  void operator delete(void* p) {
cout << "X::operator delete(void*)" << endl;</pre>
    free(p);
};
class Y {
 int filler[100];
public:
  // two arguments
  void operator delete(void* p, size_t sz) throw (const char*) {
    cout << "Freeing " << sz << " byte(s)" << endl;</pre>
    free(p);
```

```
};
};
int main() {
 X* ptr = new X;
 // call X::operator delete(void*)
 delete ptr;
 Y* yptr = new Y;
 // call Y::operator delete(void*, size t)
 // with size of Y as second argument
 delete yptr;
```

The above example will generate output similar to the following:

```
X::operator delete(void*)
Freeing 400 byte(s)
```

The statement delete ptr calls X::operator delete(void*). The statement delete yptr calls Y::operator delete(void*, size t).

The result of trying to access a deleted object is undefined because the value of the object can change after deletion.

If new and delete are called for a class object that does not declare the operator functions new and delete, or they are called for a nonclass object, the global operators **new** and **delete** are used. The global operators **new** and **delete** are provided in the C++ library.

The C++ operators for allocating and deallocating arrays of class objects are operator new[]() and operator delete[]().

You cannot declare the **delete** operator as virtual. However you can add polymorphic behavior to your delete operators by declaring the destructor of a base class as virtual. The following example demonstrates this:

```
#include <iostream>
using namespace std;
struct A {
  virtual A() { cout << "A()" << endl; };</pre>
  void operator delete(void* p) {
    cout << "A::operator delete" << endl;</pre>
    free(p);
};
struct B : A {
  void operator delete(void* p) {
    cout << "B::operator delete" << endl;</pre>
    free(p);
};
int main() {
 A* ap = new B;
  delete ap;
```

The following is the output of the above example:

```
A()
B::operator delete
```

The statement delete ap uses the **delete** operator from class B instead of class A because the destructor of A has been declared as virtual.

Although you can get polymorphic behavior from the **delete** operator, the **delete** operator that is statically visible must still be accessible even though another **delete** operator might be called. For example, in the above example, the function A::operator delete(void*) must be accessible even though the example calls B::operator delete(void*) instead.

Virtual destructors do not have any affect on deallocation operators for arrays (**operator delete[]()**). The following example demonstrates this:

```
#include <iostream>
using namespace std;
struct A {
   virtual A() { cout << "A()" << endl; }</pre>
  void operator delete[](void* p, size_t) {
    cout << "A::operator delete[]" << endl;</pre>
    ::delete [] p;
};
struct B : A {
  void operator delete[](void* p, size t) {
    cout << "B::operator delete[]" << endl;</pre>
    ::delete [] p;
};
int main() {
  A* bp = new B[3];
  delete[] bp;
};
```

The behavior of the statement delete[] bp is undefined.

When you overload the **delete** operator, you must declare it as class member, returning type **void**, with the first parameter having type **void***, as described above. You can add a second parameter of type **size_t** to the declaration. You can only have one **operator delete()** or **operator delete[]()** for a single class.

RELATED REFERENCES

- "C++ new Operator" on page 92
- "C++ delete Operator" on page 95
- "Allocation and Deallocation Functions" on page 96

Temporary Objects

It is sometimes necessary for the compiler to create temporary objects. They are used during reference initialization and during evaluation of expressions including standard type conversions, argument passing, function returns, and evaluation of the **throw** expression.

When a temporary object is created to initialize a reference variable, the name of the temporary object has the same scope as that of the reference variable. When a temporary object is created during the evaluation of a full-expression (an expression that is not a subexpression of another expression), it is destroyed as the last step in its evaluation that lexically contains the point where it was created.

There are two exceptions in the destruction of full-expressions:

- The expression appears as an initializer for a declaration defining an object: the temporary object is destroyed when the initialization is complete.
- · A reference is bound to a temporary object: the temporary object is destroyed at the end of the reference's lifetime.

If a temporary object is created for a class with constructors, the compiler calls the appropriate (matching) constructor to create the temporary object.

When a temporary object is destroyed and a destructor exists, the compiler calls the destructor to destroy the temporary object. When you exit from the scope in which the temporary object was created, it is destroyed. If a reference is bound to a temporary object, the temporary object is destroyed when the reference passes out of scope unless it is destroyed earlier by a break in the flow of control. For example, a temporary object created by a constructor initializer for a reference member is destroyed on leaving the constructor.

The ANSI/ISO C++ definition permits an implementation that eliminates the construction of such temporary objects in cases in which they are redundant.

VAC++ The VisualAge C++ compiler takes advantage of this fact to create more efficient optimized code. Take this into consideration when debugging your programs, especially for memory problems.

RELATED REFERENCES

- "Arguments of catch Blocks" on page 324
- "Initializing References" on page 69
- "Cast Expressions" on page 108
- "Function Return Values" on page 140

User-Defined Conversions

User-defined conversions allow you to specify object conversions with constructors or with conversion functions. User-defined conversions are implicitly used in addition to standard conversions for conversion of initializers, functions arguments, function return values, expression operands, expressions controlling iteration, selection statements, and explicit type conversions.

There are two types of user-defined conversions:

- Conversion by constructor
- Conversion functions

The compiler can use only one user-defined conversion (either a conversion constructor or a conversion function) when implicitly converting a single value. The following example demonstrates this:

```
class A {
 int x;
public:
 operator int() { return x; };
```

```
class B {
    A y;
public:
    operator A() { return y; };
};
int main () {
    B b_obj;
// int i = b_obj;
    int j = A(b_obj);
}
```

The compiler would not allow the statement int $i = b_obj$. The compiler would have to implicitly convert b_obj into an object of type A (with B::operator A()), then implicitly convert that object to an integer (with A::operator int()). The statement int $j = A(b_obj)$ explicitly converts b_obj into an object of type A, then implicitly converts that object to an integer.

User-defined conversions must be unambiguous, or they are not called. A conversion function in a derived class does not hide another conversion function in a base class unless both conversion functions convert to the same type. Function overload resolution selects the most appropriate conversion function. The following example demonstrates this:

```
class A {
 int a int;
 char* a_carp;
public:
 operator int() { return a int; }
 operator char*() { return a_carp; }
class B : public A {
  float b float;
 char* b carp;
 operator float() { return b float; }
 operator char*() { return b carp; }
};
int main () {
 B b obj;
// long a = b_obj;
 char* c_p = b_obj;
```

The compiler would not allow the statement long a = b_obj. The compiler could either use A::operator int() or B::operator float() to convert b_obj into a long. The statement char* c_p = b_obj uses B::operator char*() to convert b_obj into a char* because B::operator char*() hides A::operator char*().

When you call a constructor with an argument and you have not defined a constructor accepting that argument type, only standard conversions are used to convert the argument to another argument type acceptable to a constructor for that class. No other constructors or conversions functions are called to convert the argument to a type acceptable to a constructor defined for that class. The following example demonstrates this:

```
class A {
public:
   A() { }
   A(int) { }
};
```

```
int main() {
  A a1 = 1.234;
  A moocow = "text string";
```

The compiler allows the statement A a1 = 1.234. The compiler uses the standard conversion of converting 1.234 into an int, then implicitly calls the converting constructor A(int). The compiler would not allow the statement A moocow = "text string"; converting a text string to an integer is not a standard conversion.

RELATED REFERENCES

• "Standard Type Conversions" on page 116

Conversion by Constructor

A converting constructor is a constructor that can be called with one parameter, but is not declared without the function specifier explicit. The compiler uses converting constructors to convert objects from the type of the first parameter to the type of the converting constructor's class. The following example demonstrates this:

```
class Y {
  int a, b;
  char* name;
public:
  Y(int i) { };
  Y(const char* n, int j = 0) { };
void add(Y) { };
int main() {
  // equivalent to
  // \text{ obj1} = Y(2)
  Y \text{ obj1} = 2;
  // equivalent to
  // obj2 = Y("somestring",0)
Y obj2 = "somestring";
  // equivalent to
  // \text{ obj1} = Y(10)
  ob.j1 = 10;
  // equivalent to
  // add(Y(5))
  add(5);
```

The above example has the following two converting constructors:

- Y(int i)which is used to convert integers to objects of class Y.
- $Y(const\ char*\ n$, int j=0) which is used to convert pointers to strings to objects of class Y.

The compiler will not implicitly convert types as demonstrated above with constructors declared with the explicit keyword. The compiler will only use explicitly declared constructors in **new** expressions, the **static_cast** expressions and explicit casts, and the initialization of bases and members. The following example demonstrates this:

```
class A {
public:
    explicit A() { };
    explicit A(int) { };
};
int main() {
    A z;
// A y = 1;
    A x = A(1);
    A w(1);
    A * v = new A(1);
    A u = (A)1;
    A t = static_cast<A>(1);
}
```

The compiler would not allow the statement A y = 1 because this is an implicit conversion; class A has no conversion constructors.

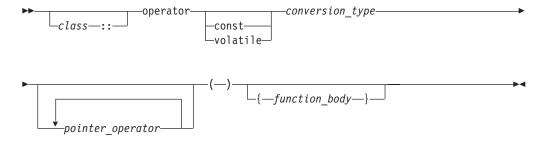
A copy constructor is a converting constructor.

RELATED REFERENCES

- "The explicit Keyword" on page 121
- "C++ new Operator" on page 92
- "static_cast Operator" on page 82
- "Cast Expressions" on page 108

Conversion Functions

You can define a member function of a class, called a *conversion function*, that converts from the type of its class to another specified type.



A conversion function that belongs to a class X specifies a conversion from the class type X to the type specified by the *conversion_type*. The following code fragment shows a conversion function called operator int():

```
class Y {
  int b;
public:
  operator int();
};
Y::operator int() {
  return b;
}
void f(Y obj) {
  int i = int(obj);
  int j = (int)obj;
  int k = i + obj;
}
```

All three statements in function f(Y) use the conversion function Y::operator

Classes, enumerations, typedef names, function types, or array types cannot be declared or defined in the conversion_type. You cannot use a conversion function to convert an object of type A to type A, a base class of A, or void.

Conversion functions have no arguments, and the return type is implicitly the conversion type. Conversion functions can be inherited. You can have virtual conversion functions but not static ones.

Copy Constructors

The copy constructor lets you create a new object from an existing one by initialization. A copy constructor of a class A is a non-template constructor in which the first parameter is of type A&, const A&, volatile A&, or const volatile A&, and the rest of its parameters (if there are any) have default values.

If you do not declare a copy constructor for a class A, the compiler will implicitly declare one for you, which will be an inline public member.

The following example demonstrates implicitly defined and user-defined copy constructors:

```
#include <iostream>
using namespace std;
struct A {
 int i;
 A() : i(10) \{ \}
struct B {
 int j;
 B() : j(20) {
   cout << "Constructor B(), j = " << j << endl;
 B(B& arg) : j(arg.j) {
    cout << "Copy constructor B(B&), j = " << j << endl;</pre>
 B(const B&, int val = 30): j(val) {
    cout << "Copy constructor B(const B&, int), j = " << j << endl;</pre>
};
struct C {
 C() { }
 C(C&) { }
};
int main() {
 A a;
 A a1(a);
 B b;
 const B b const;
 B b1(b);
 B b2(b_const);
 const C c const;
// C c1(c const);
```

The following is the output of the above example:

```
Constructor B(), j = 20

Constructor B(), j = 20

Copy constructor B(B&), j = 20

Copy constructor B(const B&, int), j = 30
```

The statement A a1(a) creates a new object from a with an implicitly defined copy constructor. The statement B b1(b) creates a new object from b with the user-defined copy constructor B::B(B&). The statement B b2(b_const) creates a new object with the copy constructor B::B(const B&, int). The compiler would not allow the statement C c1(c_const) because a copy constructor that takes as its first parameter an object of type const C& has not been defined.

The implicitly declared copy constructor of a class A will have the form A::A(const A&) if the following are true:

- The direct and virtual bases of A have copy constructors whose first parameters have been qualified with **const** or **const volatile**
- The non-static class type or array of class type data members of A have copy constructors whose first parameters have been qualified with const or const volatile

If the above are not true for a class A, the compiler will implicitly declare a copy constructor with the form A::A(A&).

The compiler cannot allow a program in which the compiler must implicitly define a copy constructor for a class A and one or more of the following are true:

- Class A has a nonstatic data member of a type which has an inaccessible or ambiguous copy constructor.
- Class A is derived from a class which has an inaccessible or ambiguous copy constructor.

The compiler will implicitly define an implicitly declared constructor of a class A if you initialize an object of type A or an object derived from class A.

An implicitly defined copy constructor will copy the bases and members of an object in the same order that a constructor would initialize the bases and members of the object.

RELATED REFERENCES

"Constructors and Destructors Overview" on page 257

Copy Assignment Operators

The *copy assignment operator* lets you create a new object from an existing one by initialization. A copy assignment operator of a class A is a nonstatic nontemplate member function that has one of the following forms:

- A::operator=(A)
- A::operator=(A&)
- A::operator=(const A&)
- A::operator=(volatile A&)
- A::operator=(const volatile A&)

If you do not declare a copy assignment operator for a class A, the compiler will implicitly declare one for you which will be inline public.

The following example demonstrates implicitly defined and user-defined copy assignment operators:

```
#include <iostream>
using namespace std;
struct A {
 A& operator=(const A&) {
   cout << "A::operator=(const A&)" << endl;</pre>
    return *this;
 A& operator=(A&) {
    cout << "A::operator=(A&)" << endl;</pre>
    return *this;
};
class B {
 A a;
};
struct C {
 C& operator=(C&) {
   cout << "C::operator=(C&)" << endl;</pre>
    return *this;
 C() { }
};
int main() {
 В х, у;
 x = y;
 Aw,z;
 W = Z;
 C i;
 const C j();
// i = j;
```

The following is the output of the above example:

```
A::operator=(const A&)
A::operator=(A&)
```

The assignment x = y calls the implicitly defined copy assignment operator of B, which calls the user-defined copy assignment operator A::operator=(const A&). The assignment w = z calls the user-defined operator A::operator=(A&). The compiler will not allow the assignment i = j because an operator C::operator=(const C&) has not been defined.

The implicitly declared copy assignment operator of a class A will have the form A& A::operator=(const A&) if the following are true:

- A direct or virtual base B of class A has a copy assignment operator whose parameter is of type const B&, const volatile B&, or B.
- A non-static class type data member of type X that belongs to class A has a copy constructor whose parameter is of type const X&, const volatile X&, or X.

If the above are not true for a class A, the compiler will implicitly declare a copy assignment operator with the form A& A::operator=(A&).

The implicitly declared copy assignment operator returns a reference to the operator's argument.

The copy assignment operator of a derived class hides the copy assignment operator of its base class.

The compiler cannot allow a program in which the compiler must implicitly define a copy assignment operator for a class A and one or more of the following are true:

- Class A has a nonstatic data member of a const type or a reference type
- Class A has a nonstatic data member of a type which has an inaccessible copy assignment operator
- Class A is derived from a base class with an inaccessible copy assignment operator.

An implicitly defined copy assignment operator of a class A will first assign the direct base classes of A in the order that they appear in the definition of A. Next, the implicitly defined copy assignment operator will assign the nonstatic data members of A in the order of their declaration in the definition of A.

RELATED REFERENCES

• "Constructors and Destructors Overview" on page 257

Chapter 16. Templates

A template describes a set of related classes or set of related functions in which a list of parameters in the declaration describe how the members of the set vary. The compiler generates new classes or functions when you supply arguments for these parameters; this process is called *template instantiation*. This class or function definition generated from a template and a set of template parameters is called a *specialization*.

Syntax — Template Declaration

```
template—<—template_parameter_list—>—declaration————
```

The *template_parameter_list* is a comma-separated list of the following kinds of template parameters:

- non-type
- type
- template

The declaration is one of the following::

- · a declaration or definition of a function or a class
- a definition of a member function or a member class
- a definition of a static data member of a class template
- a definition of a static data member of a class nested within a class template
- a definition of a member template of a class or class template

The *identifier* of a *type* is defined to be a *type_name* in the scope of the template declaration. A template declaration can appear as a namespace scope or class scope declaration.

The following example demonstrates the use of a class template:

```
template<class L> class Key
{
   L k;
   L* kptr;
   int length;
public:
   Key(L);
   // ...
};
```

Suppose the following declarations appear later:

```
Key<int> i;
Key<char*> c;
Key<mytype> m;
```

The compiler would create three objects. The following table shows the definitions of these three objects if they were written out in source form:

class Key <int> i;</int>	class Key <char*> c;</char*>	class Key <mytype> m;</mytype>
class Key <int></int>	class Key	class Key
{	{	{
int k;	char* k;	mytype k;
int * kptr;	char** kptr;	mytype* kptr;
int length;	int length;	int length;
public:	public:	public:
<pre>Key(int);</pre>	<pre>Key(char*);</pre>	<pre>Key(mytype);</pre>
//	//	//
};	};	};

Note that these three classes have different names. The arguments contained within the angle braces are not just the arguments to the class names, but part of the class names themselves. Key<int> and Key<char*> are class names.

RELATED REFERENCES

- "Template Instantiation" on page 303
- "Template Specialization" on page 305
- "Template Parameters"

Template Parameters

There are three kinds of template parameters:

- type
- · non-type
- template

You may interchange the keywords class and typename in a template parameter declaration. You cannot use storage class specifiers (static and auto) in a template parameter declaration.

RELATED REFERENCES

- "Type Template Parameters"
- "Non-Type Template Parameters" on page 285
- "Template Template Parameters" on page 285

Type Template Parameters

The following is the syntax of a type template parameter declaration:

Syntax —Type Template Parameter Declaration



The *identifier* is the name of a type.

RELATED REFERENCES

"The Keyword typename" on page 314

Non-Type Template Parameters

The syntax of a non-type template parameter is the same as a declaration of one of the following types:

- integral or enumeration
- pointer to object or pointer to function
- · reference to object or reference to function
- pointer to member

Non-type template parameters that are declared as arrays or functions are converted to pointers or pointer to functions, respectively. The following example demonstrates this:

```
template<int a[4]> struct A { };
template<int f(int)> struct B { };
int q;
int g(int) {return 0;}

A<&q> x;
B<&g> y;
```

The type of &q is int *, and the type of &g is int (*)(int).

You may qualify a non-type template parameter with **const** or **volatile**.

You cannot declare a non-type template parameter as a floating point, class, or void type.

Non-type template parameters are not lvalues.

RELATED REFERENCES

- "volatile and const Qualifiers" on page 50
- "Lvalues and Rvalues" on page 75

Template Template Parameters

The following is the syntax of a template template parameter declaration:

Syntax — Template Template Parameter Declaration

```
▶▶template—<—template-parameter-list—>—class——identifier——=—id-expression——
```

The following example demonstrates a declaration and use of a template template parameter:

```
template<template <class T> class X> class A { };
template<class T> class B { };
A<B> a;
```

RELATED REFERENCES

"Template Parameters" on page 284

Default Template Parameters

C++ Template parameters may have default arguments. The set of default template arguments accumulates over all declarations of a given template. The following example demonstrates this:

```
template<class T, class U = int> class A;
template<class T = float, class U> class A;
template<class T, class U> class A {
  public:
     T x;
     Uy;
};
A<> a;
```

The type of member a.x is float, and the type of a.y is int.

You cannot give default arguments to the same template parameters in different declarations in the same scope. For example, the compiler will not allow the following:

```
template<class T = char> class X;
template<class T = char> class X { };
```

If one template parameter has a default argument, then all template parameters following it must also have default arguments. For example, the compiler will not allow the following:

```
template<class T = char, class U, class V = int> class X { };
```

Template parameter U needs a default argument or the default for T must be removed.

The scope of a template parameter starts from the point of its declaration to the end of its template definition. This implies that you may use the name of a template parameter in other template parameter declarations and their default arguments. The following example demonstrates this:

```
template<class T = int> class A;
template<class T = float> class B;
template<class T> T f(T* p = new T);
template<class T, class U = T> class C { };
```

RELATED REFERENCES

• "Template Parameters" on page 284

Template Arguments

There are three kinds of template arguments corresponding to the three types of template parameters:

- type
- · non-type
- template

A template argument must match the type and form specified by the corresponding parameter declared in the template.

To use the default value of a template parameter, you omit the corresponding template argument. However, even if all template parameters have defaults, you still must use the <> brackets. For example, the following will yield a syntax error:

```
template<class T = int> class X { };
X<> a;
X b;
```

The last declaration, X b, will yield an error.

RELATED REFERENCES

- "Template Type Arguments"
- "Template Non-Type Arguments"
- "Template Template Arguments" on page 289

Template Type Arguments

You cannot use one of the following as a template argument for a type template parameter:

- · a local type
- a type with no linkage
- an unnamed type
- a type compounded from any of the above types

If it is ambiguous whether a template argument is a type or an expression, the template argument is considered to be a type. The following example demonstrates this:

```
template<class T> void f() { };
template<int i> void f() { };
int main() {
  f<int()>();
}
```

The function call f<int()>() calls the function with T as a template argument – the compiler considers int() as a type – and therefore implicitly instantiates and calls the first f().

RELATED REFERENCES

- "Local Scope" on page 1
- "No Linkage" on page 6
- "Declaring and Using Bit Fields in Structures" on page 39
- "typedef" on page 29

Template Non-Type Arguments

A non-type template argument provided within a template argument list is an expression whose value can be determined at compile time. Such arguments must be constant expressions, addresses of functions or objects with external linkage, or addresses of static class members. Non-type template arguments are normally used to initialize a class or to specify the sizes of class members.

For non-type integral arguments, the instance argument matches the corresponding template argument as long as the instance argument has a value and sign appropriate to the argument type.

For non-type address arguments, the type of the instance argument must be of the form identifier or &identifier, and the type of the instance argument must match the template argument exactly, except that a function name is changed to a pointer to function type before matching.

The resulting values of non-type template arguments within a template argument list form part of the template class type. If two template class names have the same template name and if their arguments have identical values, they are the same class.

In the following example, a class template is defined that requires a non-type template int argument as well as the type argument:

```
template<class T, int size> class myfilebuf
     T* filepos;
     static int array[size];
public:
     myfilebuf() { /* ... */ }
      ~myfilebuf();
     advance(); // function defined elsewhere in program
};
```

In this example, the template argument size becomes a part of the template class name. An object of such a template class is created with both the type argument T of the class and the value of the non-type template argument size.

An object x, and its corresponding template class with arguments **double** and size=200, can be created from this template with a value as its second template argument:

```
myfilebuf<double,200> x;
```

x can also be created using an arithmetic expression:

```
myfilebuf<double,10*20> x;
```

The objects created by these expressions are identical because the template arguments evaluate identically. The value 200 in the first expression could have been represented by an expression whose result at compile time is known to be equal to 200, as shown in the second construction.

Note: Arguments that contain the < symbol or the > symbol must be enclosed in parentheses to prevent it from being parsed as a template argument list delimiter when it is being used as a relational operator or a nested template delimiter. For example, the arguments in the following definition are valid:

```
myfilebuf<double, (75>25)> x;
                                    // valid
```

The following definition, however, is not valid because the greater than operator (>) is interpreted as the closing delimiter of the template argument

```
myfilebuf<double, 75>25> x;
                               // error
```

If the template arguments do not evaluate identically, the objects created are of different types:

```
myfilebuf<double,200> x;
                                 // create object x of class
                                 // myfilebuf<double,200>
myfilebuf<double,200.0> y;
                                 // error, 200.0 is a double,
                                  // not an int
```

The instantiation of y fails because the value 200.0 is of type **double**, and the template argument is of type **int**.

The following two objects:

```
myfilebuf<double, 128> x myfilebuf<double, 512> y
```

are objects of separate template specializations. Referring either of these objects later with myfilebuf<double> is an error.

A class template does not need to have a type argument if it has non-type arguments. For example, the following template is a valid class template:

```
template<int i> class C
{
    public:
        int k;
        C() { k = i; }
};
```

This class template can be instantiated by declarations such as:

```
class C<100>;
class C<200>;
```

Again, these two declarations refer to distinct classes because the values of their non-type arguments differ.

RELATED REFERENCES

- "Integer Constant Expressions" on page 75
- "References" on page 68
- "External Linkage" on page 5
- "Static Members" on page 218

Template Template Arguments

A template argument for a template template parameter is the name of a class template.

When the compiler tries to find a template to match the template template argument, it only considers primary class templates. (A *primary template* is the template which is being specialized.) The compiler will not consider any partial specialization even if their parameter lists match that of the template template parameter. For example, the compiler will not allow the following code:

```
template<class T, int i> class A {
   int x;
};

template<class T> class A<T, 5> {
   short x;
};

template<template<class T> class U> class B1 { };

B1<A> c;
```

The compiler will not allow the declaration B1<A> c. Although the partial specialization of A seems to match the template template parameter of B1, the compiler considers only the primary template of A, which has different template parameters than B1.

The compiler considers the partial specializations based on a template template argument once you have instantiated a specialization based on the corresponding template template parameter. The following example demonstrates this:

```
#include <iostream>
using namespace std;
template<class T, class U> class A {
template<class U> class A<int, U> {
   short x;
};
template<template<class T, class U> class V> class B {
   V<int, char> i;
   V<char, char> j;
};
B<A> c;
int main() {
   cout << typeid(c.i.x).name() << endl;</pre>
   cout << typeid(c.j.x).name() << endl;</pre>
```

The following is the output of the above example:

short int

The declaration V<int, char> i uses the partial specialization while the declaration V<char, char> j uses the primary template.

RELATED REFERENCES

- "Partial Specialization" on page 311
- "Template Instantiation" on page 303

Class Templates

The relationship between a class template and an individual class is like the relationship between a class and an individual object. An individual class defines how a group of objects can be constructed, while a class template defines how a group of classes can be generated.

Note the distinction between the terms *class template* and *template class*:

Class template

is a template used to generate template classes. You cannot declare an object of a class template.

Template class

is an instance of a class template.

A template definition is identical to any valid class definition that the template might generate, except for the following:

• The class template definition is preceded by

```
template < template-parameter-list >
```

where *template-parameter-list* is a comma-separated list of one or more of the following kinds of template parameters:

- type
- non-type
- template
- Types, variables, constants and objects within the class template can be declared using the template parameters as well as explicit types (for example, int or char).

A class template be declared without being defined by using an elaborated type specifier. For example:

```
template <class L,class T> class key;
```

This reserves the name as a class template name. All template declarations for a class template must have the same types and number of template arguments. Only one template declaration containing the class definition is allowed.

Note: When you have nested template argument lists, you must have a separating space between the > at the end of the inner list and the one at the end of the outer list. Otherwise, there is an ambiguity between the output operator >> and two template list delimiters >.

```
template <class L, class T> class key { /* ... */ };
template <class L> class vector { /* ... */ };;
int main ()
{
class key <int, vector<int> >; // implicitly instantiates template
}
```

Objects and function members of individual template classes can be accessed by any of the techniques used to access ordinary class member objects and functions. Given a class template:

the constructor, the constructed object, and the member function drive() can be accessed with any of the following (assuming the standard header file <string.h> is included in the program file):

constructor	vehicle <char> bicycle;</char>	
	<pre>// constructor called automatically, // object bicycle created</pre>	
object bicycle	strcpy (bicycle.kind, "10 speed"); bicycle.kind[0] = '2';	
function drive()	<pre>char* n = bicycle.drive();</pre>	
function roadmap()	vehicle <char>::roadmap();</char>	

RELATED REFERENCES

- "Declaring Class Types" on page 199
- "Scope of Class Names" on page 203

Class Template Declarations and Definitions

A class template must be declared before any declaration of a corresponding template class. A class template definition can only appear once in any single compilation unit. A class template must be defined before any use of a template class that requires the size of the class or refers to members of the class.

None of these order restrictions apply to the incremental compiler with unordered programming.

In the following example, the class template key is declared before it is defined. The declaration of the pointer keyiptr is valid because the size of the class is not needed. The declaration of keyi, however, causes an error.

```
template <class L> class key;
                                   // class template declared.
                                   // not defined yet
class key<int> *keyiptr;
                                  // declaration of pointer
class key<int> keyi;
                                  // error, cannot declare keyi
                                   // without knowing size
template <class L> class key
                                   // now class template defined
{ /* ... */ };
```

If a template class is used before the corresponding class template is defined, the compiler issues an error. A class name with the appearance of a template class name is considered to be a template class. In other words, angle brackets are valid in a class name only if that class is a template class.

The definition of a class template is not compiled until the definition of a template class is required. At that point, the class template definition is compiled using the argument list of the template class to instantiate the template arguments. Any errors in the class definition are flagged at this time.

RELATED REFERENCES

• "Class Templates" on page 290

Static Data Members and Templates

Each class template instantiation has its own copy of any static data members. The static declaration can be of template argument type or of any defined type.

You must separately define static members. The following example demonstrates this:

```
template <class T> class K
{
public:
        static T x;
};
template <class T> T K<T> ::x;
int main()
{
        K<int>::x = 0;
}
```

The statement template<class T> T K<T>::x defines the static member of class T, while the statement in the main() function initializes the data members for K.

RELATED REFERENCES

• "Static Members" on page 218

Member Functions of Class Templates

You may define a template member function outside of its class template definition.

When you call a member function of a class template specialization, the compiler will use the template arguments that you used to generate the class template. The following example demonstrates this:

```
template<class T> class X {
   public:
        T operator+(T);
};

template<class T> T X<T>::operator+(T arg1) {
   return arg1;
};

int main() {
   X<char> a;
   X<int> b;
   a +'z';
   b + 4;
}
```

The overloaded addition operator has been defined outside of class X. The statement a + 'z' is equivalent to a operator+('z'). The statement b + 4 is equivalent to b operator+(4).

RELATED REFERENCES

• "Member Functions" on page 211

Friends and Templates

There are three kinds of relationships between template classes and their functions declared as friends:

- One-to-many: A non-template function may be a friend to all template class instantiations.
- One-to-one: A template function instantiated with one set of template arguments may be a friend to one template class instatiated with the same set of template arguments
- Many-to-many: All instantiations of a template function may be a friend to all instantiations of the template class.

The following example demonstrates these relationships:

```
template<class S> g();
template<class T> class A {
   friend int e();
   friend int f(T);
   friend int g<T>();
   template<class U> friend int h();
};
```

- Function e() has a one-to-many relationship with class A. Function e() is a friend to all instantiations of class A.
- Function f() has a one-to-one relationship with class A. The compiler will give you a warning for this kind of declaration similar to the following:
 - The friend function declaration "f" will cause an error when the enclosing template class is instantiated with arguments that declare a friend function that does not match an existing definition. The function declares only one function because it is not a template but the function type depends on one or more template parameters.
- Function g() has a one-to-one relationship with class A. Function g() is a function template. It must be declared before here or else the compiler will not recognize g<T> as a template name. For each instantiation of A there is one matching instantiation of g(). For example, g<int> is a friend of A<int>.
- Function h() has a many-to-many relationship with class A. Function h() is a function template. For all instantiations of A all instantiations of h() are friends.

These relationships also apply to friend classes.

RELATED REFERENCES

• "Friends" on page 225

Function Templates

A function template defines how a group of functions can be generated.

A non-template function is not related to a function template, even though the non-template function may have the same name and type of a specialization the function template might generate. A non-template function is never considered to be a specialization of a function template.

The following example implements the QuickSort algorithm with a function template named quicksort:

```
#include <iostream>
#include <cstdlib>
using namespace std;
template<class T> void quicksort(T a[], const int& leftarg, const int& rightarg) {
```

```
if (leftarg < rightarg) {</pre>
    T pivotvalue = a[leftarg];
   int left = leftarg - 1;
   int right = rightarg + 1;
  for(;;) {
    while (a[--right] > pivotvalue);
    while (a[++left] < pivotvalue);</pre>
   if (left >= right) break;
   T temp = a[right];
    a[right] = a[left];
   a[left] = temp;
 int pivot = right;
 quicksort(a, leftarg, pivot);
 quicksort(a, pivot + 1, rightarg);
int main(void) {
 int sortme[10];
 for (int i = 0; i < 10; i++) {
    sortme[i] = rand();
   cout << sortme[i] << " ";</pre>
 };
 cout << endl;
 quicksort<int>(sortme, 0, 10 - 1);
 for (int i = 0; i < 10; i++) cout << sortme[i] << " ";
 cout << endl;</pre>
 return 0;
```

The above example will have output similar to the following: 16838 5758 10113 17515 31051 5627 23010 7419 16212 4086 4086 5627 5758 7419 10113 16212 16838 17515 23010 31051

This QuickSort algorithm will sort an array of type T (whose relational and assignment operators have been defined). The template function takes one template argument and three function arguments:

- the type of the array to be sorted, T
- the name of the array to be sorted, a
- the lower bound of the array, leftarg
- the upper bound of the array, rightarg

In the above example, you can also call the quicksort() template function with the following statement:

```
quicksort(sortme, 0, 10 - 1);
```

You may omit any template argument if the compiler can deduce it by the usage and context of the template function call. In this case, the compiler deduces that sortme is an array of type int.

RELATED REFERENCES

• "Template Argument Deduction"

Template Argument Deduction

C++ When you call a template function, you may omit any template argument that the compiler can determine or deduce by the usage and context of that template function call.

The compiler tries to deduce a template argument by comparing the type of the corresponding template parameter with the type of the argument used in the function call. The two types that the compiler compares (the template parameter and the argument used in the function call) must be of a certain structure in order for template argument deduction to work. The following lists these type structures:

```
const T
volatile T
T&
T*
T[10]
A<T>
C(*)(T)
T(*)()
T(*)(U)
T C::*
C T::*
T U::*
T (C::*)()
C (T::*)()
D (C::*)(T)
C (T::*)(U)
T (C::*)(U)
T (U::*)()
T (U::*)(V)
E[10][i]
B<i>
TT<T>
TT<i>>
TT<C>
```

- T, U, and V represent a template type argument
- 10 represents any integer constant
- i represents a template non-type argument
- [i] represents an array bound of a reference or pointer type, or a non-major array bound of a normal array.
- TT represents a template template argument
- (T), (U), and (V) represents an argument list that has at least one template type
- () represents an argument list that has no template arguments
- <T> represents a template argument list that has at least one template type argument
- <i>represents a template argument list that has at least one template non-type argument
- <C> represents a template argument list that has no template arguments dependent on a template parameter

The following example demonstrates the use of each of these type structures. The example declares a template function using each of the above structures as an argument. These functions are then called (without template arguments) in order of declaration. The example outputs the same list of type structures:

```
#include <iostream>
using namespace std;
template<class T> class A { };
template<int i> class B { };
class C {
  public:
      int x;
};
class D {
   public:
      Су;
      int z;
};
                                        {cout << "T" << endl; };
template<class T> void f (T)
                                       {cout << "const T" << endl; };
template<class T> void f1(const T)
template<class T> void f2(volatile T) {cout << "volatile T" << endl; };</pre>
                                       {cout << "T*" << endl; };
template<class T> void g (T*)
                                       {cout << "T&" << endl; };
template<class T> void g (T&)
                                       {cout << "T[10]" << endl;};
template<class T> void g1(T[10])
template<class T> void h1(A<T>)
                                       {cout << "A<T>" << endl; };
void test 1() {
  A<char> a;
   C c;
   f(c); f1(c);
                    f2(c);
   g(c);
           g(&c);
                    g1(&c);
   h1(a);
}
                            void j(C(*)(T)) {cout << "C(*)(T)" << endl; };
template<class T>
                           void j(T(*)()) {cout << "T(*)()" << endl; }</pre>
template<class T>
template<class T, class U> void j(T(*)(U)) {cout << "T(*)(U)" << endl; };
void test 2() {
   C (*c pfunct1)(int);
   C (*c pfunct2)(void);
   int (*c pfunct3)(int);
   j(c pfunct1);
   j(c pfunct2);
   j(c pfunct3);
                            void k(T C::*) {cout << "T C::*" << endl; };</pre>
template<class T>
                            void k(C T::*) {cout << "C T::*" << endl; };</pre>
template<class T>
template<class T, class U> void k(T U::*) {cout << "T U::*" << endl; };
void test 3() {
   k(\&C::x);
   k(\&D::y);
   k(\&D::z);
                                     void m(T (C::*)() ) {cout << "T (C::*)()" << endl; };</pre>
template<class T>
                                     void m(C (T::*)() ) {cout << "C (T::*)()" << endl; };</pre>
template<class T>
                                     void m(D(C::*)(T)) {cout << "D(C::*)(T)" << endl; };
template<class T>
                                     void m(C (T::*)(U)) {cout << "C (T::*)(U)" << endl; };</pre>
template<class T, class U>
                                     void m(T (C::*)(U)) {cout << "T (C::*)(U)" << endl; };</pre>
template<class T, class U>
template<class T, class U>
                                     void m(T (U::*)() ) {cout << "T (U::*)()" << endl; };</pre>
template<class T, class U, class V> void m(T(U::*)(V)) {cout << "T(U::*)(V)" << end1; };
void test 4() {
   int (C::*f membp1)(void);
```

```
C (D::*f membp2)(void);
  D (C::*f membp3)(int);
  m(f membp1);
  m(f_membp2);
  m(f_membp3);
  C (D::*f membp4)(int);
  int (C::*f_membp5)(int);
  int (D::*f_membp6)(void);
  m(f_membp4);
  m(f membp5);
  m(f membp6);
  int (D::*f membp7)(int);
  m(f membp7);
template<int i> void n(C[10][i]) {cout << "E[10][i]" << endl; };
template<int i> void n(B<i>) {cout << "B<i>" << endl; };</pre>
void test 5() {
  C array [10][20];
  n(array);
  B<20> b;
  n(b);
template < template < class TT, class T> void p1(TT<T>) {cout << "TT<T>" << endl; };
template<template<int> class TT, int i> void p2(TT<i>) {cout << "TT<i>" << end]; };</pre>
                                            void p3(TT<C>) {cout << "TT<C>" << end1; };</pre>
template<template<class> class TT>
void test 6() {
  A<char> a;
  B<20> b;
  A<C>c;
  p1(a);
  p2(b);
  p3(c);
int main() { test 1(); test 2(); test 3(); test 4(); test 5(); test 6(); }
```

RELATED REFERENCES

- "Deducing Type Template Arguments"
- "Deducing Non-Type Template Arguments" on page 300

Deducing Type Template Arguments

The compiler can deduce template arguments from a type composed of several of the listed type structures. The following example demonstrates template argument deduction for a type composed of several type structures:

```
template<class T> class Y { };
template<class T, int i> class X {
  public:
     Y<T> f(char[20][i]) { return x; };
     Y < T > x;
};
template<template<class> class T, class U, class V, class W, int i>
  void g( T<U> (V::*)(W[20][i]) ) { };
int main()
```

```
{
    Y<int> (X<int, 20>::*p)(char[20][20]) = &X<int, 20>::f;
    g(p);
}
```

The type Y<int> (X<int, 20>::*p)(char[20][20])T<U> (V::*)(W[20][i]) is based on the type structure T (V::*)(V):

- T is Y<int>
- U is X<int, 20>
- V is char[20][20]

If you qualify a type with the class to which that type belongs, and that class (a *nested name specifier*) depends on a template parameter, the compiler will not deduce a template argument for that parameter. If a type contains a template argument that cannot be deduced for this reason, all template arguments in that type will not be deduced. The following example demonstrates this:

```
template<class T, class U, class V>
  void h(typename Y<T>::template Z<U>, Y<T>, Y<V>) { };
int main() {
  Y<int>::Z<char> a;
  Y<int> b;
  Y<float> c;

  h<int, char, float>(a, b, c);
  h<int, char>(a, b, c);
  // h<int>(a, b, c);
}
```

The compiler will not deduce the template arguments T and U in typename Y<T>::template <math>Z<U> (but it will deduce the T in Y<T>). The compiler would not allow the template function call h<int>(a, b, c) because U is not deduced by the compiler.

The compiler can deduce a function template argument from a pointer to function or pointer to member function argument given several overloaded function names. However, none of the overloaded functions may be function templates, nor can more than one overloaded function match the required type. The following example demonstrates this:

```
template < class T > void f(void(*)(T,int)) { };
template < class T > void g1(T, int) { };
void g2(int, int) { };
void g2(char, int) { };
void g3(int, int, int) { };
void g3(float, int) { };
int main() {
// f(&g1);
// f(&g2);
 f(&g3);
}
```

The compiler would not allow the call f(&g1) because g1() is a function template. The compiler would not allow the call f(&g2) because both functions named g2() match the type required by f().

The compiler cannot deduce a template argument from the type of a default argument. The following example demonstrates this:

```
template<class T> void f(T = 2, T = 3) \{ \};
int main() {
  f(6);
// f();
  f<int>();
```

The compiler allows the call f(6) because the compiler deduces the template argument (int) by the value of the function call's argument. The compiler would not allow the call f() because the compiler cannot deduce template argument from the default arguments of f().

The compiler cannot deduce a template type argument from the type of a non-type template argument. For example, the compiler will not allow the following:

```
template<class T, T i> void f(int[20][i]) { };
int main() {
  int a[20][30];
   f(a);
```

The compiler cannot deduce the type of template parameter T.

RELATED REFERENCES

• "Template Argument Deduction" on page 296

Deducing Non-Type Template Arguments

The compiler cannot deduce the value of a major array bound unless the bound refers to a reference or pointer type. Major array bounds are not part of function parameter types. The following code demonstrates this:

```
template<int i> void f(int a[10][i]) { };
template<int i> void g(int a[i]) { };
template<int i> void h(int (&a)[i]) { };
int main ()
   int b[10][20];
   int c[10];
  f(b);
  // g(c);
  h(c);
```

The compiler would not allow the call g(c); the compiler cannot deduce template argument i.

The compiler cannot deduce the value of a non-type template argument used in an expression in the template function's parameter list. The following example demonstrates this:

```
template<int i> class X { };
template<int i> void f(X<i - 1>) { };
int main () {
```

```
X<0> a;
f<1>(a);
// f(a);
}
```

In order to call function f() with object a, the function must accept an argument of type X<0>. However, the compiler cannot deduce that the template argument i must be equal to 1 in order for the function template argument type X<i-1> to be equivalent to X<0>. Therefore the compiler would not allow the function call f(a).

If you want the compiler to deduce a non-type template argument, the type of the parameter must match exactly the type of value used in the function call. For example, the compiler will not allow the following:

```
template<int i> class A { };
template<short d> void f(A<d>) { };
int main() {
    A<1> a;
    f(a);
}
```

The compiler will not convert int to short when the example calls f().

However, deduced array bounds may be of any integral type.

RELATED REFERENCES

• "Template Argument Deduction" on page 296

Overloading Function Templates

You may overload a function template either by a non-template function or by another function template.

If you call the name of an overloaded function template, the compiler will try to deduce its template arguments and check its explicitly declared template arguments. If successful, it will instantiate a function template specialization, then add this specialization to the set of *candidate functions* used in overload resolution. The compiler proceeds with overload resolution, choosing the most appropriate function from the set of candidate functions. Non-template functions take precedence over template functions. The following example describes this:

```
#include <iostream>
using namespace std;

template<class T> void f(T x, T y) { cout << "Template" << endl; }

void f(int w, int z) { cout << "Non-template" << endl; }

int main() {
   f(1, 2);
   f('a', 'b');
   f(1, 'b');
}</pre>
```

The following is the output of the above example:

```
Non-template
Template
Non-template
```

The function call f(1, 2) could match the argument types of both the template function and the non-template function. Because in overload resolution non-template functions take precendence, the non-template function is called.

The function call f('a', 'b') can only match the argument types of the template function. The template function is called.

Argument deduction fails for the function call f(1, 'b'); the compiler does not generate any template function specialization and overload resolution does not take place. The non-template function resolves this function call after using the standard conversion from **char** to **int** for the function argument 'b'.

RELATED REFERENCES

• "Overload Resolution" on page 196

Partial Ordering of Function Templates

C++ A function template specialization might be ambiguous because template argument deduction might associate the specialization with more than one of the overloaded definitions. The compiler will then choose the definition that is the most specialized. This process of selecting a function template definition is called partial ordering.

A template X is more specialized than a template Y if every argument list that matches the one specified by X also matches the one specified by Y, but not the other way around. The following example demonstrates partial ordering:

```
template<class T> void f(T)
template<class T> void f(T*) { }
template<class T> void f(const T*) { }
template<class T> void g(T) { }
template<class T> void g(T&) { }
template<class T> void h(T) { }
template<class T> void h(T, ...) { }
int main() {
  const int *p;
   f(p);
  int q;
// g(q);
// h(q);
```

The declaration template<class T> void f(const T*) is more specialized than template<class T> void f(T*). Therefore, the function call f(p) calls template<class T> void f(const T*). However, neither void g(T) nor void g(T)& are more specialized than the other. Therefore, the function call g(q) would be ambiguous.

Ellipses do not affect partial ordering. Therefore, the function call h(q) would also be ambiguous.

The compiler uses partial ordering in the following cases:

- Calling a function template specialization that requires overload resolution.
- Taking the address of a function template specialization.

- When a friend function declaration, an explicit instantiation, or explicit specialization refers to a function template specialization.
- Determining the appropriate deallocation function that is also a function template for a given placement operator new.

RELATED REFERENCES

- "Template Specialization" on page 305
- "C++ new Operator" on page 92

Template Instantiation

The act of creating a new definition of a function, class, or member of a class from a template declaration and one or more template arguments is called *template instantiation*. The definition created from a template instantiation is called a *specialization*.

RELATED REFERENCES

• "Template Specialization" on page 305

Implicit Instantiation

Unless a template specialization has been explicitly instantiated or explicitly specialized, the compiler will generate a specialization for the template only when it needs the definition. This is called *implicit instantiation*.

If the compiler must instantiate a class template specialization and the template is declared, you must also define the template.

For example, if you declare a pointer to a class, the definition of that class is not needed and the class will not be implicitly instantiated. The following example demonstrates when the compiler instantiates a template class:

```
template<class T> class X {
  public:
          X* p;
          void f();
          void g();
};

X<int>* q;
X<int> r;
X<float>* s;
r.f();
s->g();
```

The compiler requires the instantiation of the following classes and functions:

- X<int> when the object r is declared
- X<int>::f() at the member function call r.f()
- X<float> and X<float>::g() at the class member access function call s->g()

Therefore, the functions X<T>::g() and X<T>::g() must be defined in order for the above example to compile. (The compiler will use the default constructor of class X when it creates object r.) The compiler does not require the instantiation of the following definitions:

- · class X when the pointer p is declared
- X<int> when the pointer q is declared
- X<float> when the pointer s is declared

The compiler will implicitly instantiate a class template specialization if it is involved in pointer conversion or pointer to member conversion. The following example demonstrates this:

```
template<class T> class B { };
template<class T> class D : public B<T> { };
void g(D<double>* p, D<int>* q)
 B < double > * r = p;
 delete q;
```

The assignment B<double>* r = p converts p of type D<double> to a type of B<double>; the compiler must instantiate D<double>. The compiler must instantiate D<int> when it tries to delete q.

If the compiler implicitly instantiates a class template that contains static members, those static members are not implicitly instantiated. The compiler will instantiate a static member only when the compiler needs the static member's definition. Every instantiated class template specialization has its own copy of static members. The following example demonstrates this:

```
template<class T> class X {
public:
   static T v;
};
template<class T> T X<T>::v = 0;
X<char*> a;
X<float> b;
X<float> c;
```

Object a has a static member variable v of type char*. Object b has a static variable v of type float. Objects b and c share the single static data member v.

An implicitly instantiated template is in the same namespace where you defined the template.

If a function template or a member function template specialization is involved with overload resolution, the compiler implicitly instantiates a declaration of the specialization.

RELATED REFERENCES

• "Template Instantiation" on page 303

Explicit Instantiation

You can explicitly tell the compiler when it should generate a definition from a template. This is called *explicit instantiation*.

Syntax — Explicit Instantiation Declaration

```
▶▶—template—template_declaration—
```

The following are examples of explicit instantiations:

```
template<class T> class Array { void mf(); };
template class Array<char>;
template void Array<int>::mf();
template<class T> void sort(Array<T>& v) { }
template void sort(Array<char>&);
namespace N {
  template<class T> void f(T&) { }
template void N::f<int>(int&);
int* p = 0;
template<class T> T g(T = &p);
template char g(char);
template <class T> class X {
  private:
     T v(T arg) { return arg; };
template int X<int>::v(int);
template<class T> T g(T val) {return val;}
template<class T> void Array<T>::mf() { }
```

A template declaration must be in scope at the point of instantiation of the template's explicit instantiation. An explicit instantiation of a template specialization is in the same namespace where you defined the template.

Access checking rules do not apply to names in explicit instantiations. Template arguments and names in a declaration of an explicit instantiation may be private types or objects. In the above example, the compiler allows the the explicit instantiation template int X<int>::v(int) even though the member function is declared private.

The compiler does not use default arguments when you explicitly instantiate a template. In the above example the compiler allows the explicit instantiation template char g(char) even though the default argument is an address of type int.

RELATED REFERENCES

"Template Instantiation" on page 303

Template Specialization

The act of creating a new definition of a function, class, or member of a class from a template declaration and one or more template arguments is called *template instantiation*. The definition created from a template instantiation is called a *specialization*. A *primary template* is the template which is being specialized.

RELATED REFERENCES

"Template Instantiation" on page 303

Explicit Specialization

■ C++■ When you instantiate a template with a given set of template arguments the compiler generates a new definition based on those template arguments. You can override this behavior of definition generation. You can instead specify the definition the compiler uses for a given set of template arguments. This is called explicit specialization. You can explicitly specialize any of the following:

Function or class template

using namespace std;

- Member function of a class template
- Static data member of a class template
- Member class of a class template
- Member function template of a class template
- Member class template of a class template

Syntax — Explicit Specialization Declaration

```
▶▶template—<->—declaration name
                                                                    —declaration body-►◀
                                   _<-template_argument_list-->-
```

The **template**<> prefix indicates that the following template declaration takes no template parameters. The declaration_name is the name of a previously declared template. The following example demonstrates explicit specialization:

```
template<class T = float, int i = 5> class A
   public:
      A();
      int value;
};
template<> class A<> { public: A(); };
template<> class A<double, 10> { public: A(); };
template<class T, int i> A<T, i>::A() : value(i) {
   cout << "Primary template,"</pre>
        << "non-type argument is " << value << endl;</pre>
A <> :: A()  {
   cout << "Explicit specialization "</pre>
        << "default arguments" << endl;</pre>
A<double, 10>::A() {
  cout << "Explicit specialization "</pre>
        << "<double, 10>" << endl;
int main() {
   A<int,6> x;
   A <> y;
```

The following is the output of the above example:

```
Primary template non-type argument is: 6
Explicit specialization default arguments
Explicit specialization <double, 10>
```

A<double, 10> z;

This example declared two explicit specializations for the primary template (the template which is being specialized) class A. Object x uses the constructor of the primary template. Object y uses the explicit specialization A <> :: A(). Object z uses the explicit specialization A < 0.

RELATED REFERENCES

- "Function Templates" on page 294
- "Class Templates" on page 290
- "Member Functions of Class Templates" on page 293
- "Static Data Members and Templates" on page 293

Definition and Declaration of Explicit Specializations

The definition of an explicitly specialized class is unrelated to the definition of the primary template. You do not have to define the primary template in order to define the specialization (nor do you have to define the specialization in order to define the primary template). For example, the compiler will allow the following:

```
template<class T> class A;
template<> class A<int>;
template<> class A<int> { /* ... */ };
```

The primary template is not defined, but the explicit specialization is.

You can use the name of an explicit specialization that has been declared but not defined the same way as an incompletely defined class. The following example demonstrates this:

```
template<class T> class X { };
template<> class X<char>;
X<char>* p;
X<int> i;
// X<char> j;
```

The compiler does not allow the declaration X<char> j because the explicit specialization of X<char> is not defined.

RELATED REFERENCES

"Explicit Specialization" on page 306

Explicit Specialization and Scope

A declaration of a primary template must be in scope at the *point of declaration* of the explicit specialization. In other words, an explicit specialization declaration must appear after the declaration of the primary template. For example, the compiler will not allow the following:

```
template<> class A<int>;
template<class T> class A;
```

An explicit specialization is in the same namespace as the definition of the primary template.

RELATED REFERENCES

"Explicit Specialization" on page 306

Class Members of Explicit Specializations

C++ A member of an explicitly specialized class is not implicitly instantiated from the member declaration of the primary template. You have to explicitly define members of a class template specialization. You define members of an explicitly specialized template class as you would normal classes, without the template<> prefix. In addition, you can define the members of an explicit specialization inline; no special template syntax is used in this case. The following example demonstrates a class template specialization:

```
template<class T> class A {
  public:
      void f(T);
template<> class A<int> {
  public:
     int g(int);
int A<int>::g(int arg) { return 0; }
int main() {
  A<int> a;
  a.g(1234);
```

The explicit specialization A<int> contains the member function g(), which the primary template does not.

If you explicitly specialize a template, a member template, or the member of a class template, then you must declare this specialization before that specialization is implicitly instantiated. For example, the compiler will not allow the following code:

```
template<class T> class A { };
void f() { A<int> x; }
template<> class A<int> { };
int main() { f(); }
```

The compiler will not allow the the explicit specialization template<> class A<int> { }; because function f() uses this specialization (in the construction of x) before the specialization.

The unordered incremental compiler allows the above example. In this case the compiler will use the explicit specialization when it calls function f().

RELATED REFERENCES

"Explicit Specialization" on page 306

Explicit Specialization of Function Templates

In a function template specialization, a template argument is optional if the compiler can deduce it from the type of the function arguments. The following example demonstrates this:

```
template<class T> class X { };
template<class T> void f(X<T>);
template<> void f(X<int>);
```

The explicit specialization template<> void f(X<int>) is equivalent to template<> void f<int>(X<int>).

You cannot specify default function arguments in a declaration or a definition for any of the following:

- Explicit specialization of a function template
- Explicit specialization of a member function template

For example, the compiler will not allow the following code:

```
template<class T> void f(T a) { };
template<> void f<int>(int a = 5) { };

template<class T> class X {
  void f(T a) { }
};
template<> void X<int>::f(int a = 10) { };
```

RELATED REFERENCES

- "Explicit Specialization" on page 306
- "Function Templates" on page 294

Explicit Specialization of Members of Class Templates

Each instantiated class template specialization has its own copy of any static members. You may explicitly specialize static members. The following example demonstrates this:

```
template<class T> class X {
public:
    static T v;
    static void f(T);
};

template<class T> T X<T>::v = 0;
template<class T> void X<T>::f(T arg) { v = arg; }

template<> char* X<char*>::v = "Hello";
template<> void X<float>::f(float arg) { v = arg * 2; }

int main() {
    X<char*> a, b;
    X<float> c;
    c.f(10);
}
```

This code explicitly specializes the initialization of static data member X::v to point to the string "Hello" for the template argument char*. The function X::f() is explicitly specialized for the template argument float. The static data member v in objects a and b point to the same string, "Hello". The value of c.v is equal to 20 after the call function call c.f(10).

You can nest member templates within many enclosing class templates. If you explicitly specialize a template nested within several enclosing class templates, you must prefix the declaration with **template<>** for every enclosing class template you specialize. You may leave some enclosing class templates unspecialized, however you cannot explicitly specialize a class template unless its enclosing class templates are also explicitly specialized. The following example demonstrates explicit specialization of nested member templates:

```
#include <iostream>
using namespace std;
```

```
template<class T> class X {
public:
  template<class U> class Y {
 public:
    template<class V> void f(U,V);
   void g(U);
 };
};
template<class T> template<class U> template<class V>
 void X<T>::Y<U>::f(U, V) { cout << "Template 1" << endl; }</pre>
template<class T> template<class U>
 void X<T>::Y<U>::g(U) { cout << "Template 2" << endl; }
template<> template<>
 void X<int>::Y<int>::g(int) { cout << "Template 3" << endl; }</pre>
template<> template<> template<class V>
 void X<int>::Y<int>::f(int, V) {cout << "Template 4" << endl; }</pre>
template<> template<> template<>
 void X<int>:::f<int>(int, int) { cout << "Template 5" << endl; }</pre>
// template<> template<class U> template<class V>
      void X<char>::Y<U>::f(U, V) { cout << "Template 6" << endl; }</pre>
// template<class T> template<>
      void X<T>::Y<float>::g(float) { cout << "Template 7" << endl; }</pre>
int main() {
 X<int>::Y<int> a;
 X<char>::Y<char> b;
 a.f(1, 2);
 a.f(3, 'x');
 a.g(3);
b.f('x', 'y');
 b.g('z');
```

The following is the output of the above program:

Template 5 Template 4 Template 3 Template 1 Template 2

- The compiler would not allow the the template specialization definition that would output "Template 6" because it is attempting to specialize a member (function f()) without specializatin its containing class (Y).
- The compiler would not allow the template specialization definition that would output "Template 7" because the enclosing class of class Y (which is class X) is not explicitly specialized.

A friend declaration cannot declare an explicit specialization.

RELATED REFERENCES

- "Explicit Specialization" on page 306
- "Static Data Members and Templates" on page 293

Partial Specialization

When you instantiate a class template, the compiler creates a definition based on the template arguments you have passed. Alternatively, if all those template arguments match those of an explicit specialization, the compiler uses the definition defined by the explicit specialization.

A partial specialization is a generalization of explicit specialization. An explicit specialization only has a template argument list. A partial specialization has both a template argument list and a template parameter list. The compiler uses the partial specialization if its template argument list matches a subset of the template arguments of a template instantiatation. The compiler will then generate a new definition from the partial specialization with the rest of the unmatched template arguments of the template instantiation.

You cannot partially specialize function templates.

Syntax — Partial Specialization

```
\textcolor{red}{\blacktriangleright \texttt{template}} - \textcolor{red}{template} - parameter\_list \rightarrow \textcolor{red}{-} declaration\_name - \textcolor{red}{-} - template\_argument\_list - \textcolor{red}{-} - declaration\_body - \textcolor{red}{\blacktriangleright} - \textcolor{red}{-} - \textcolor{r
```

The *declaration_name* is a name of a previously declared template. The following demonstrates the use of partial specializations:

```
#include <iostream>
using namespace std;
template<class T, class U, int I> struct X
  { void f() { cout << "Primary template" << endl; } };
template<class T, int I> struct X<T, T*, I>
  { void f() { cout << "Partial specialization 1" << endl; } };
template<class T, class U, int I> struct X<T*, U, I>
  { void f() { cout << "Partial specialization 2" << endl; } };
template<class T> struct X<int, T*, 10>
  { void f() { cout << "Partial specialization 3" << endl; } };
template<class T, class U, int I> struct X<T, U*, I>
  { void f() { cout << "Partial specialization 4" << endl; } };
int main() {
  X<int, int, 10> a;
  X<int, int*, 5> b;
  X<int*, float, 10> c;
  X<int, char*, 10> d;
  X<float, int*, 10> e;
// X<int, int*, 10> f;
  a.f(); b.f(); c.f(); d.f(); e.f();
```

The following is the output of the above example:

```
Primary template
Partial specialization 1
Partial specialization 2
Partial specialization 3
Partial specialization 4
```

The compiler would not allow the declaration X<int, int*, 10> f because it can match template struct X<T, T*, I>, template struct X<int, T*, 10>, or template struct X<T, U*, I>, and none of these declarations are a better match than the others.

Each class template partial specialization is a separate template. You must provide definitions for each member of a class template partial specialization.

RELATED REFERENCES

"Template Specialization" on page 305

Template Parameter and Argument Lists of Partial **Specializations**

Primary templates do not have template argument lists; this list is implied in the template parameter list.

Template parameters specified in a primary template but not used in a partial specialization are omitted from the template parameter list of the partial specialization. The order of a partial specialization's argument list is the same as the order of the primary template's implied argument list.

In a template argument list of a partial template parameter, you cannot have an expression that involves non-type arguments unless that expression is only an identifier. In the following example, the compiler will not allow the first partial specialization, but will allow the second one:

```
template<int I, int J> class X { };
// Invalid partial specialization
template<int I> class \langle I * 4, I + 3 \rangle \{ \};
// Valid partial specialization
template <int I> class <I, I> { };
```

The type of a non-type template argument cannot depend on a template parameter of a partial specialization. The compiler will not allow the following partial specialization:

```
template<class T, T i> class X { };
// Invalid partial specialization
template<class T> class X<T, 25> { };
```

A partial specialization's template argument list cannot be the same as the list implied by the primary template.

You cannot have default values in the template parameter list of a partial specialization.

RELATED REFERENCES

- "Template Parameters" on page 284
- "Template Arguments" on page 286

Matching of Class Template Partial Specializations

The compiler determines whether to use the primary template or one of its partial specializations by matching the template arguments of the class template specialization with the template argument lists of the primary template and the partial specializations:

- If the compiler finds only one specialization, then the compiler generates a definition from that specialization.
- If the compiler finds more than one specialization, then the compiler tries to determine which of the specializations is the most specialized. A template X is more specialized than a template Y if every argument list that matches the one specified by X also matches the one specified by Y, but not the other way around. If the compiler cannot find the most specialized specialization, then the use of the class template is ambiguous; the compiler will not allow the program.
- If the compiler does not find any matches, then the compiler generates a definition from the primary template.

RELATED REFERENCES

• "Partial Specialization" on page 311

Name Binding and Dependent Names

Name binding is the process of finding the declaration for each name that is explicitly or implicitly used in a template. The compiler may bind a name in the definition of a template, or it may bind a name at the instantiation of a template.

A *dependent name* is a name that depends on the type or the value of a template argument. For example:

```
template<class T> class U : A<T>
{
  typename T::B x;
  void f(A<T>& y)
  {
    *y++;
  }
};
```

The dependent names in this example are the base class A<T>, the type name T::B, and the variable y.

The compiler binds dependent names when a template is instantiated. The compiler binds non-dependent names when a template is defined. For example:

```
void f(double) { cout << "Function f(double)" << endl; }

template<class T> void g(T a) {
   f(123);
   h(a);
}

void f(int) { cout << "Function f(int)" << endl; }

void h(double) { cout << "Function h(double)" << endl; }

void i() {
   extern void h(int);
   g<int>(234);
}

void h(int) { cout << "Function h(int)" << endl; }</pre>
```

The following is the output if you call function i(): Function f(double) Function h(double)

The unordered incremental compiler will not give this result; the compiler will call the best match of the overloaded functions regardless of their location.

The point of definition of a template is located immediately before its definition. In this example, the point of definition of the function template g(T) is located immediately before the keyword template. Because the function call f(123) does not depend on a template argument, the compiler will consider names declared before the definition of function template g(T). Therefore, the call f(123) will call f(double). Although f(int) is a better match, it is not in scope at the point of definition of g(T).

The point of instantiation of a template is located immediately before the declaration that encloses its use. In this example, the point of instantiation of the call to g<int>(234) is located immediately before i(). Because the function call h(a) depends on a template argument, the compiler will consider names declared before the instantiation of function template g(T). Therefore, the call h(a) will call h(double). It will not consider h(int), because this function was not in scope at the point of instantiation of g<int>(234).

Point of instantiation binding implies the following:

- A template parameter cannot depend on any local name or class member.
- An unqualified name in a template cannot depend on a local name or class member.

RELATED REFERENCES

"Template Instantiation" on page 303

The Keyword typename

Solution Use the keyword **typename** if you have a qualified name that refers to a type and depends on a template parameter. Only use the keyword typename in template declarations and definitions. The following example illustrates the use of the keyword typename:

```
template<class T> class A
 T::x(y);
 typedef char C;
 A::C d;
```

The statement T::x(y) is ambiguous. It could be a call to function x() with a non-local argument y, or it could be a declaration of variable y with type T::x. C++ will interpret this statement as a function call. In order for the compiler to interpret this statement as a declaration you would add the keyword typename to the beginning of it. The statement A:: C d; is ill-formed. The class A also refers to A<T> and thus depends on a template parameter. You must add the keyword typename to the beginning of this declaration:

```
typename A::C d;
```

You can also use the keyword typename in place of the keyword class in template parameter declarations.

RELATED REFERENCES

• "Template Parameters" on page 284

The Keyword template as Qualifier

Use the keyword **template** as a qualifier to distinguish member templates from other names. The following example illustrates when you must use **template** as a qualifier:

```
class A
{
  public:     template<class T> T function_m() { };
};

template<class U> void function_n(U argument)
{
  char object_x = argument.function_m<char>();
}
```

The declaration char object_x = argument.function_m<char>(); is ill-formed. The compiler assumes that the < is a less-than operator. In order for the compiler to regcognize the function template call, you must add the **template** quantifier:

```
char object_x = argument.template function_m<char>();
```

If the name of a member template specialization appears after a ., ->, or :: operator, and that name has explicitly qualified template parameters, prefix the member template name with the keyword **template**. The following example demonstrates this use of the keyword **template**:

```
#include <iostream>
using namespace std;
class X {
   public:
      template <int j> struct S {
         void h() {
            cout << "member template's member function: " << j << endl;</pre>
      };
      template <int i> void f() {
        cout << "Primary: " << i << endl;</pre>
};
template<> void X::f<20>() {
   cout << "Specialized, non-type argument = 20" << endl;</pre>
template<class T> void g(T* p) {
   p->template f<100>();
   p->template f<20>();
   typename T::template S<40> s; // use of scope operator on a member template
   s.h();
int main()
   X temp;
   g(&temp);
```

The following is the output of the above example:

Primary: 100 Specialized, non-type argument = 20

member template's member function: 40

If you do not use the keyword template in these cases, the compiler will interpret the < as a less-than operator. For example, the following line of code is ill-formed: p->f<100>();

The compiler interprets f as a non-template member, and the < as a less-than operator.

RELATED REFERENCES

• "Chapter 16. Templates" on page 283

Chapter 17. Exception Handling

Exception handling is a mechanism that separates code that detects and handles exceptional circumstances from the rest of your program. Note that an exceptional circumstance is not necessarily an error.

When a function detects an exceptional situation, you represent this with an object. This object is called an *exception object*. In order to deal with the exceptional situation you *throw the exception*. This passes control, as well as the exception, to a designated block of code in a direct or indirect caller of the function that threw the exception. This block of code is called a *handler*. In a handler, you specify the types of exceptions that it may process. The C++ runtime, together with the generated code, will pass control to the first appropriate handler that is able to process the the exception thrown. When this happens, an exception is *caught*. A handler may *rethrow* an exception so it can be caught by another handler.

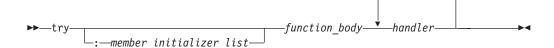
The exception handling mechanism is made up of the following elements:

- try blocks: a block of code that may throw an exception that you want to handle with special processing
- catch blocks or handlers: a block of code that is executed when a try block encounters an exception
- throw expression: indicates when your program encounters an exception
- exception specifications: specify which exceptions (if any) a function may throw
- **unexpected()** function: called when a function throws an exception not specified by an exception specification
- terminate() function: called for exceptions that are not caught

Syntax — try Block



Syntax — Function try Block



Syntax — Exception Handler or catch Block



Syntax — throw Expression



RELATED REFERENCES

- · "The try Keyword"
- "catch Blocks" on page 320
- "The throw Expression" on page 326
- "Exception Specifications" on page 329
- "unexpected()" on page 332
- "terminate()" on page 333

The try Keyword

Sou use a try block to indicate which areas in your program that might throw exceptions you want to handle immediately. You use a function try block to indicate that you want to detect exceptions in the entire body of a function.

Syntax — try Block



Syntax — Function try Block



The following is an example of a function try block with a member initializer, a function try block and a try block:

```
#include <iostream>
using namespace std;
class E {
   public:
      const char* error;
      E(const char* arg) : error(arg) { }
};
class A {
   public:
      int i;
      // A function try block with a member
      // initializer
      A() try : i(0) {
         throw E("Exception thrown in A()");
      catch (E& e) {
         cout << e.error << endl;</pre>
```

```
};
// A function try block
void f() try {
   throw E("Exception thrown in f()");
catch (E& e) {
   cout << e.error << endl;</pre>
void g() {
   throw E("Exception thrown in g()");
int main() {
   f();
   // A try block
   try {
      g();
   catch (E& e) {
      cout << e.error << endl;</pre>
   A x;
}
```

The following is the output of the above example:

```
Exception thrown in f()
Exception thrown in g()
Exception thrown in A()
```

The constructor of class A has a function try block with a member initializer. Function f() has a function try block. The main() function contains a try block.

RELATED REFERENCES

• "Initializing Base Classes and Members" on page 262

Nested Try Blocks

When try blocks are nested and a **throw** occurs in a function called by an inner try block, control is transferred outward through the nested try blocks until the first catch block is found whose argument matches the argument of the throw expression.

```
For example:

try
{
    func1();
    try
    {
        func2();
    }
    catch (spec_err) { /* ... */ }
    func3();
}
catch (type_err) { /* ... */ }
// if no throw is issued, control resumes here.
```

In the above example, if spec_err is thrown within the inner try block (in this case, from func2()), the exception is caught by the inner catch block, and, assuming this

catch block does not transfer control, func3() is called. If spec err is thrown after the inner try block (for instance, by func3()), it is not caught and the function terminate() is called. If the exception thrown from func2() in the inner try block is type err, the program skips out of both try blocks to the second catch block without invoking func3(), because no appropriate catch block exists following the inner try block.

You can also nest a try block within a catch block.

RELATED REFERENCES

- "terminate()" on page 333
- "unexpected()" on page 332
- "Special Exception Handling Functions" on page 332

catch Blocks

The following is the syntax for an exception handler or a catch block:

You can declare a handler to catch many types of exceptions. The allowable objects that a function can catch are declared in the parentheses following the catch keyword (the exception_declaration). You can catch objects of the fundamental types, base and derived class objects, references, and pointers to all of these types. You can also catch const and volatile types. The exception_declaration cannot be an incomplete type, or a reference or pointer to an incomplete type other than one of the following:

- void*
- · const void*
- volatile void*
- const volatile void*

You cannot define a type in an exception_declaration.

You can also use the catch(...) form of the handler to catch all thrown exceptions that have not been caught by a previous catch block. The ellipsis in the catch argument indicates that any exception thrown can be handled by this handler.

If an exception is caught by a catch(...) block, there is no direct way to access the object thrown. Information about an exception caught by catch(...) is very limited.

You can declare an optional variable name if you want to access the thrown object in the catch block.

A catch block can only catch accessible objects. The object caught must have an accessible copy constructor.

RELATED REFERENCES

- "volatile and const Qualifiers" on page 50
- "Member Access" on page 223

Function try block Handlers

The scope and lifetime of the parameters of a function or constructor extend into the handlers of a function try block. The following example demonstrates this:

```
void f(int &x) try {
    throw 10;
}
catch (const int &i)
{
    x = i;
}
int main() {
    int v = 0;
    f(v);
}
```

The value of v after f() is called is 10.

A function try block on **main()** does not catch exceptions thrown in destructors of objects with static storage duration, or constructors of namespace scope objects.

The following example throws an exception from a destructor of a static object:

```
#include <iostream>
using namespace std;
class E {
public:
 const char* error;
  E(const char* arg) : error(arg) { }
class A {
  public: A() { throw E("Exception in A()"); }
class B {
 public: B() { throw E("Exception in B()"); }
int main() try {
 cout << "In main" << endl;</pre>
  static A cow;
 B bull;
catch (E& e) {
  cout << e.error << endl;</pre>
```

The following is the output of the above example:

```
In main Exception in \tilde{B}()
```

The runtime will not catch the exception thrown when object cow is destroyed at the end of the program.

The following example throws an exception from a constructor of a namespace scope object:

```
#include <iostream>
using namespace std;
```

```
class E {
public:
  const char* error;
  E(const char* arg) : error(arg) { }
};
namespace N {
  class C {
  public:
    C() {
      cout << "In C()" << endl;</pre>
       throw E("Exception in C()");
  };
  C calf;
};
int main() try {
  cout << "In main" << endl;</pre>
catch (E& e) {
  cout << e.error << endl;</pre>
```

The following is the output of the above example: In C()

The compiler will not catch the exception thrown when object calf is created.

In a function try block's handler, you cannot have a jump into the body of a constructor or destructor.

A return statement cannot appear in a function try block's handler of a constructor.

When the function try block's handler of an object's constructor or destructor is entered, fully constructed base classes and members of that object are destroyed. The following example demonstrates this:

```
#include <iostream>
using namespace std;
class E {
   public:
      const char* error;
      E(const char* arg) : error(arg) { };
};
class B {
   public:
      B() { cout << "B() called" << endl; };
};
class D : public B {
   public:
       D() { cout << "D() called" << endl; };
};
D::D() try : B() {
   throw E("Exception in D()");
catch(E& e) {
   cout << "Handler of function try block of D(): " << e.error << endl;</pre>
```

```
};
int main() {
    D val;
};
```

The following is the output of the above example:

```
B() called
Handler of function try block of D(): Exception in D()
```

When the function try block's handler of D() is entered, the runtime first calls the destructor of the base class of D, which is B. The destructor of D is not called because val is not fully constructed.

The runtime will rethrow an exception at the end of a function try block's handler of a constructor or destructor. All other functions will return once they have reached the end of their function try block's handler. The following example demonstrates this:

```
#include <iostream>
using namespace std;
class E {
   public:
      const char* error;
      E(const char* arg) : error(arg) { };
};
class A {
      A() try { throw E("Exception in A()"); }
      catch(E& e) { cout << "Handler in A(): " << e.error << endl; }</pre>
};
int f() try {
   throw E("Exception in f()");
   return 0;
catch(E& e) {
   cout << "Handler in f(): " << e.error << endl;</pre>
   return 1;
int main() {
   int i = 0;
   try { A cow; }
   catch(E& e) {
      cout << "Handler in main(): " << e.error << endl;</pre>
   try { i = f(); }
   catch(E& e) {
      cout << "Another handler in main(): " << e.error << endl;</pre>
   cout << "Returned value of f(): " << i << endl;</pre>
}
The following is the output of the above example:
Handler in A(): Exception in A()
Handler in main(): Exception in A()
Handler in f(): Exception in f()
Returned value of f(): 1
```

RELATED REFERENCES

- "The main() Function" on page 132
- "static Storage Class Specifier" on page 28
- "Chapter 10. Namespaces" on page 177
- "Destructors" on page 266

Arguments of catch Blocks

If you specify a class type for the argument of a catch block (the exception_declaration), the compiler will use a copy constructor to initialize that argument. If that argument does not have a name, the compiler initializes a temporary object. The compiler destroys this object when the handler exits.

The ANSI/ISO C++ definition permits an implementation that eliminates the construction of such temporary objects in cases in which they are redundant.

YAC++ The VisualAge C++ compiler takes advantage of this fact to create more efficient optimized code. Take this into consideration when debugging your programs, especially for memory problems.

RELATED REFERENCES

• "Temporary Objects" on page 273

Matching between Exceptions Thrown and Caught

SEC++ An argument in the catch argument of a handler matches an argument in the assignment_expression of the throw expression (throw argument) if any of the following conditions is met:

- The catch argument type matches the type of the thrown object.
- The catch argument is a public base class of the thrown class object.
- The catch specifies a pointer type, and the thrown object is a pointer type that can be converted to the pointer type of the catch argument by standard pointer conversion.

Note: If the type of the thrown object is const or volatile, the catch argument must also be a const or volatile for a match to occur. However, a const, volatile, or reference type catch argument can match a nonconstant, nonvolatile, or nonreference object type. A nonreference catch argument type matches a reference to an object of the same type.

RELATED REFERENCES

- "Pointer Conversions" on page 118
- "volatile and const Qualifiers" on page 50
- "References" on page 68
- "Special Exception Handling Functions" on page 332

Order of Catching

If the compiler encounters an exception in a try block, it will try each handler in order of appearance.

Always place a catch block that catches a derived class before a catch block that catches the base class of that derived class (following a try block). If a catch block for objects of a base class is followed by a catch block for objects of a derived class of that base class, the compiler issues a warning and continues to compile the program despite the unreachable code in the derived class handler.

A catch block of the form **catch(...)** must be the last catch block following a try block or an error occurs. This placement ensures that the **catch(...)** block does not prevent more specific catch blocks from catching exceptions intended for them.

If the runtime cannot cannot find a matching handler in the current scope, the runtime will continue to find a matching handler in a dynamically surrounding try block. The following example demonstrates this:

```
#include <iostream>
using namespace std;
class E {
public:
  const char* error;
  E(const char* arg) : error(arg) { };
class F : public E {
public:
  F(const char* arg) : E(arg) { };
void f() {
  try {
    cout << "In try block of f()" << endl;</pre>
    throw E("Class E exception");
  catch (F& e) {
    cout << "In handler of f()";</pre>
    cout << e.error << endl;</pre>
};
int main() {
  try {
    cout << "In main" << endl;</pre>
    f();
  catch (E& e) {
    cout << "In handler of f(): ";</pre>
    cout << e.error << endl;</pre>
  };
  cout << "Resume execution in main" << endl;</pre>
```

The following is the output of the above example:

```
In try block of main
In try block of f()
In handler of main: Class E exception
Resume execution in main
```

In function f(), the runtime could not find a handler to handle the exception of type E thrown. The runtime finds a matching handler in a dynamically surrounding try block: the try block in the main() function.

If the runtime cannot find a matching handler in the program, it calls the **terminate()** function.

RELATED REFERENCES

- "The try Keyword" on page 318
- "terminate()" on page 333

The throw Expression

You use a throw expression to indicate when your program encounters an exception.

Syntax — throw Expression



The type of assignment_expression cannot be an incomplete type, or a pointer to an incomplete type other than one of the following:

- void*
- const void*
- volatile void*
- const volatile void*

The assignment_expression is treated the same way as a function argument in a call or the operand of a return statement.

If the assignment_expression is a class object, that object's copy constructor and destructor must be accessible. For example, you cannot throw a class object that has its copy constructor declared as private.

RELATED REFERENCES

"Incomplete Types" on page 51

Rethrowing an Exception

GET +-- If a catch block cannot handle the particular exception it has caught, you can rethrow the exception. The rethrow expression (throw without assignment_expression) causes the originally thrown object to be rethrown.

Because the exception has already been caught at the scope in which the rethrow expression occurs, it is rethrown out to the next dynamically enclosing try block. Therefore, it cannot be handled by catch blocks at the scope in which the rethrow expression occurred. Any catch blocks for the dynamically enclosing try block have an opportunity to catch the exception.

The following example demonstrates rethrowing an exception:

```
#include <iostream>
using namespace std;
struct E {
 const char* message;
 E() : message("Class E") { }
};
struct E1 : E {
 const char* message;
 E1() : message("Class E1") { }
```

```
struct E2 : E {
  const char* message;
  E2() : message("Class E2") { }
void f() {
  try {
    cout << "In try block of f()" << endl;</pre>
    cout << "Throwing exception of type E1" << endl;</pre>
    E1 myException;
    throw myException;
  catch (E2& e) {
    cout << "In handler of f(), catch (E2& e)" << endl;</pre>
    cout << "Exception: " << e.message << endl;</pre>
  catch (E1& e) {
    cout << "In handler of f(), catch (E1& e)" << endl;</pre>
    cout << "Exception: " << e.message << endl;</pre>
    throw;
  catch (E& e) {
    cout << "In handler of f(), catch (E& e)" << endl;</pre>
    cout << "Exception: " << e.message << endl;</pre>
    throw;
  }
}
int main() {
  try {
    cout << "In try block of main()" << endl;</pre>
    f();
  catch (E2& e) {
    cout << "In handler of main(), catch (E2& e)" << endl;</pre>
    cout << "Exception: " << e.message << endl;</pre>
  catch (...) {
    cout << "In handler of main(), catch (...)" << endl;</pre>
The following is the output of the above example:
In try block of main()
In try block of f()
Throwing exception of type E1
In handler of f(), catch (E1& e)
Exception: Class E1
In handler of main(), catch (...)
```

The try block in the main() function calls function f(). The try block in function f() throws an object of type E1 named myException. The handler catch (E1 &e) catches myException. The handler then rethrows myException with the statement throw to the next dynamically enclosing try block: the try block in the main() function. The handler catch(...) catches myException.

RELATED REFERENCES

• "The throw Expression" on page 326

Stack Unwinding

■ C++■ When an exception is thrown and control passes from a try block to a handler, the C++ runtime calls destructors for all automatic objects constructed since the beginning of the try block. This process is called *stack unwinding*. The automatic objects are destroyed in reverse order of their construction. (Automatic objects are local objects that have been declared auto or register, or not declared static or extern. An automatic object x is deleted whenever the program exits the block in which x is declared.)

If an exception is thrown during construction of an object consisting of subobjects or array elements, destructors are only called for those subobjects or array elements successfully constructed before the exception was thrown. A destructor for a local static object will only be called if the object was successfully constructed.

If during stack unwinding a destructor throws an exception and that exception is not handled, the terminate() function is called. The following example demonstrates this:

```
#include <iostream>
using namespace std;
struct E {
 const char* message;
 E(const char* arg) : message(arg) { }
};
void my terminate() {
 cout << "Call to my_terminate" << endl;</pre>
};
 A() { cout << "In constructor of A" << endl; }
  A() {
   cout << "In destructor of A" << endl;</pre>
    throw E("Exception thown in A()");
};
struct B {
 B() { cout << "In constructor of B" << endl; }
   B() { cout << "In destructor of B" << endl; }
int main() {
  set terminate(my terminate);
   cout << "In try block" << endl;</pre>
    A a;
    B b;
    throw("Exception thrown in try block of main()");
 catch (const char* e) {
    cout << "Exception: " << e << endl;</pre>
 catch (...) {
    cout << "Some exception caught in main()" << endl;</pre>
 cout << "Resume execution of main()" << endl;</pre>
```

The following is the output of the above example:

```
In try block
In constructor of A
In constructor of B
In destructor of B
In destructor of A
Call to my terminate
```

In the try block, two automatic objects are created: a and b. The try block throws an exception of type const char*. The handler catch (const char* e) catches this exception. The C++ runtime unwinds the stack, calling the destructors for a and b in reverse order of their construction. The destructor for a throws an exception. Since there is no handler in the program that can handle this exception, the C++ runtime calls terminate(). (The function terminate() calls the function specified as the argument to **set_terminate()**. In this example, **terminate()** has been specified to call my_terminate().)

RELATED REFERENCES

- "terminate()" on page 333
- "set_unexpected() and set_terminate()" on page 334

Exception Specifications

C++ provides a mechanism to ensure that a given function is limited to throwing only a specified list of exceptions. An exception specification at the beginning of any function acts as a guarantee to the function's caller that the function will throw only the exceptions contained in the exception specification.

```
For example, a function:
```

```
void translate() throw(unknown word,bad grammar) { /* ... */ }
```

explicitly states that it will only throw exception objects whose types are unknown word or bad grammar, or any type derived from unknown word or bad grammar.

Syntax – Exception Specification



The type_id_list is a comma-separated list of types. In this list you cannot specify an incomplete type, a pointer or a reference to an incomplete type, other than a pointer to void, optionally qualified with const and/or volatile. You cannot define a type in an exception specification.

A function with no exception specification allows all exceptions. A function with an exception specification that has an empty type_id_list, throw(), does not allow any exceptions to be thrown.

An exception specification is not part of a function's type.

An exception specification may only appear at the end of a function declarator of a function, pointer to function, reference to function, pointer to member function

Exception Specifications

declaration, or pointer to member function definition. An exception specification cannot appear in a typedef declaration. The following declarations demonstrate this:

```
void f() throw(int);
void (*g)() throw(int);
void h(void i() throw(int));
// typedef int (*j)() throw(int);
```

The compiler would not allow the last declaration, typedef int (*j)() throw(int).

Suppose that class A is one of the types in the *type_id_list* of an exception specification of a function. That function may throw exception objects of class A, or any class publicly derived from class A. The following example demonstrates this:

```
class A { };
class B : public A { };
class C { };
void f(int i) throw (A) {
  switch (i) {
     case 0: throw A();
      case 1: throw B();
      default: throw C();
}
void g(int i) throw (A*) {
  A* a = new A();
  B*b = new B();
  C* c = new C();
  switch (i) {
     case 0: throw a;
      case 1: throw b;
      default: throw c;
  }
}
```

Function f() can throw objects of types A or B. If the function tries to throw an object of type C, the compiler will call unexpected() because type C has not been specified in the function's exception specification, nor does it derive publicly from A. Similarly, function g() cannot throw pointers to objects of type C; the function may throw pointers of type A or pointers of objects that derive publicly from A.

A function that overloads a virtual function can only throw exceptions specified by the virtual function. The following example demonstrates this:

```
class A {
   public:
      virtual void f() throw (int, char);
};
class B : public A {
   public: void f() throw (int) { }
};
   class C : public A {
      public: void f() { }
   };
```

```
class D : public A {
    public: void f() throw (int, char, double) { }
};
*/
```

The compiler allows B::f() because the member function may throw only exceptions of type int. The compiler would not allow C::f() because the member function may throw any kind of exception. The compiler would not allow D::f() because the member function can throw more types of exceptions (int, char, and double) than A::f().

Suppose that you assign or initialize a pointer to function named x with a function or pointer to function named y. The pointer to function x can only throw exceptions specified by the exception specifications of y. The following example demonstrates this:

```
void (*f)();
void (*g)();
void (*h)() throw (int);

void i() {
   f = h;
// h = g;
}
```

The compiler allows the assignment f = h because f can throw any kind of exception. The compiler would not allow the assignment h = g because h can only throw objects of type f int, while g can throw any kind of exception.

Implicitly declared special functions (default constructors, copy constructors, destructors, and copy assignment operators) have exception specifications. An implicitly declared special function will have in its exception specification the types declared in the functions' exception specifications that the special function invokes. If any function that a special function invokes allows all exceptions, then that special function allows all exceptions. If all the functions that a special function invokes allow no exceptions, then that special function will allow no exceptions. The following example demonstrates this:

```
class A {
   public:
        A() throw (int);
        A(const A&) throw (float);
        A() throw();
};

class B : public A {
   public:
        B() throw (char);
        B(const A&);
        B() throw();
};

class C : public B { };
```

The following special functions in the above example have been implicitly declared:

```
C::C() throw (int, char);
C::C(const C&);
C::C() throw();
```

Exception Specifications

The default constructor of C can throw exceptions of type int or char. The copy constructor of C can throw any kind of exception. The destructor of C cannot throw any exceptions.

RELATED REFERENCES

- "Incomplete Types" on page 51
- "Function Declarations" on page 124
- "Pointers to Functions" on page 141
- "Chapter 15. Special Member Functions" on page 257
- "unexpected()"

Special Exception Handling Functions

Not all thrown errors can be caught and successfully dealt with by a catch block. In some situations, the best way to handle an exception is to terminate the program. Two special library functions are implemented in C++ to process exceptions not properly handled by catch blocks or exceptions thrown outside of a valid try block. These functions are **unexpected()** and **terminate()**.

RELATED REFERENCES

- "unexpected()"
- "terminate()" on page 333

unexpected()

■ C++■ When a function with an exception specification throws an exception that is not listed in its exception specification, the C++ runtime does the following:

- 1. The **unexpected()** function is called.
- 2. The **unexpected()** function calls the function pointed to by unexpected_handler. By default, unexpected_handler points to the function terminate().

You can replace the default value of unexpected_handler with the function set_unexpected().

Although **unexpected()** cannot return, it may throw (or rethrow) an exception. Suppose the exception specification of a function f() has been violated. If **unexpected()** throws an exception allowed by the exception specification of f(), then the C++ runtime will search for another handler at the call of f(). The following example demonstrates this:

```
#include <iostream>
using namespace std;
struct E {
 const char* message;
 E(const char* arg) : message(arg) { }
};
void my unexpected() {
 cout << "Call to my_unexpected" << endl;</pre>
 throw E("Exception thrown from my_unexpected");
void f() throw(E) {
 cout << "In function f(), throw const char* object" << endl;</pre>
 throw("Exception, type const char*, thrown from f()");
```

```
int main() {
    set_unexpected(my_unexpected);
    try {
       f();
    }
    catch (E& e) {
       cout << "Exception in main(): " << e.message << endl;
    }
}</pre>
```

The following is the output of the above example:

```
In function f(), throw const char* object
Call to my_unexpected
Exception in main(): Exception thrown from my unexpected
```

The main() function's try block calls function f(). Function f() throws an object of type const char*. However the exception specification of f() allows only objects of type E to be thown. The function unexpected() is called. The function unexpected() calls my_unexpected(). The function my_unexpected() throws an object of type E. Since unexpected() throws an object allowed by the exception specification of f(), the handler in the main() function may handle the exception.

If **unexpected()** did not throw (or rethrow) an object alloed by the exception specification of f(), then the C++ runtime does one of two things:

- If the exception specification of f() included the class **std::bad_exception**, **unexpected()** will throw an object of type **std::bad_exception**, and the C++ runtime will search for another handler at the call of f().
- If the exception specification of f() did not include the class **std::bad_exception**, the function **terminate()** is called.

RELATED REFERENCES

"set_unexpected() and set_terminate()" on page 334

terminate()

In some cases, the exception handling mechanism fails and a call to **void terminate()** is made. This **terminate()** call occurs in any of the following situations:

- The exception handling mechanism cannot find a handler for a thrown exception. The following are more specific cases of this:
 - During stack unwinding, a destructor throws an exception and that exception is not handled.
 - The expression that is thrown also throws an exception, and that exception is not handled.
 - The constructor or desturctor of a non-local static object throws an exception, and the exception is not handled.
 - A function registered with atexit() throws an exception, and the exception is not handled. The following demonstrates this:

```
extern "C" printf(char* ...);
#include <exception>
#include <cstdlib>
using namespace std;

void f() {
  printf("Function f()\n");
  throw "Exception thrown from f()";
}
```

```
void g() { printf("Function g()\n"); } void h() { printf("Function h()\n"); }
void my_terminate() {
  printf("Call to my terminate\n");
  abort();
int main() {
  set_terminate(my_terminate);
  atexit(f);
  atexit(g);
  atexit(h);
  printf("In main\n");
```

The following is the output of the above example:

```
Function h()
Function g()
Function f()
Call to my terminate
```

To register a function with atexit(), you pass a parameter to atexit() a pointer to the function you want to register. At normal program termination, atexit() calls the functions you have registered with no arguments in reverse order. The **atexit()** function is in the **<cstdlib>** library.

- A throw expression without an operand tries to rethrow an exception, and no exception is presently being handled.
- A function f() throws an exception that violates its exception specification. The **unexpected()** function then throws an exception which violates the exception specification of f(), and the exception specification of f() did not include the class std::bad_exception.
- The default value of **unexpected_handler** is called.

The **terminate()** function calls the function pointed to by **terminate_handler**. By default, terminate_handler points to the function abort(), which exits from the program. You can replace the default value of terminate_handler with the function set_terminate().

A terminate function cannot return to its caller, either by using **return** or by throwing an exception.

RELATED REFERENCES

"set_unexpected() and set_terminate()"

set unexpected() and set terminate()

The function **unexpected()**, when invoked, calls the function most recently supplied as an argument to **set_unexpected()**. If **set_unexpected()** has not yet been called, unexpected() calls terminate().

The function terminate(), when invoked, calls the function most recently supplied as an argument to **set_terminate()**. If **set_terminate()** has not yet been called, **terminate()** calls **abort()**, which ends the program.

You can use **set_unexpected()** and **set_terminate()** to register functions you define to be called by **unexpected()** and **terminate()**. The functions **set_unexpected()** and

set_terminate() are included in the standard header files. Each of these functions has as its return type and its argument type a pointer to function with a **void** return type and no arguments. The pointer to function you supply as the argument becomes the function called by the corresponding special function: the argument to **set_unexpected()** becomes the function called by **unexpected()**, and the argument to **set_terminate()** becomes the function called by **terminate()**.

Both **set_unexpected()** and **set_terminate()** return a pointer to the function that was previously called by their respective special functions (**unexpected()** and **terminate()**). By saving the return values, you can restore the original special functions later so that **unexpected()** and **terminate()** will once again call **terminate()** and **abort()**.

If you use **set_terminate()** to register your own function, the function should no return to its caller but terminate execution of the program.

If you attempt to return from the function called by **terminate()**, **abort()** is called instead and the program ends.

Example of Using the Exception Handling Functions

The following example shows the flow of control and special functions used in exception handling:

```
#include <iostream>
#include <exception>
using namespace std;
class X { };
class Y { };
class A { };
// pfv type is pointer to function returning void
typedef void (*pfv)();
void my terminate() {
  cout << "Call to my terminate" << endl;</pre>
  abort();
void my_unexpected() {
  cout << "Call to my unexpected()" << endl;</pre>
  throw;
void f() throw(X,Y, bad_exception) {
  throw A();
void g() throw(X,Y) {
  throw A();
int main()
  pfv old term = set terminate(my terminate);
  pfv old unex = set unexpected(my unexpected);
    cout << "In first try block" << endl;</pre>
    f();
  catch(X) {
    cout << "Caught X" << endl;</pre>
```

```
catch(Y) {
  cout << "Caught Y" << endl;</pre>
catch (bad exception& e1) {
  cout << "Caught bad_exception" << endl;</pre>
catch (...) {
  cout << "Caught some exception" << endl;</pre>
cout << endl;
try {
 cout << "In second try block" << endl;</pre>
  g();
catch(X) {
  cout << "Caught X" << endl;</pre>
catch(Y) {
  cout << "Caught Y" << endl;</pre>
catch (bad exception& e2) {
  cout << "Caught bad exception" << endl;</pre>
catch (...) {
  cout << "Caught some exception" << endl;</pre>
```

The following is the output of the above example:

```
In first try block
Call to my_unexpected()
Caught bad exception
In second try block
Call to my unexpected()
Call to my terminate
```

At run time, this program behaves as follows:

- 1. The call to **set_terminate()** assigns to old_term the address of the function last passed to **set_terminate()** when **set_terminate()** was previously called.
- 2. The call to **set_unexpected()** assigns to old_unex the address of the function last passed to **set_unexpected()** when **set_unexpected()** was previously called.
- 3. Within the first try block, function f() is called. Because f() throws an unexpected exception, a call to unexpected() is made. unexpected() in turn calls my_unexpected(), which prints a message to standard output. The function my unexpected() tries to rethow the exception of type A. Because class A has not been specified in the exception specification of function f(), my unexpected() throws an exception of type **bad_exception**.
- 4. Because bad_exception has been specified in the exception specification of function f(), the handler catch (bad exception e1) is able to handle the exception.
- 5. Within the second try block, function q() is called. Because q() throws an unexpected exception, a call to **unexpected()** is made. The **unexpected()** throws an exception of type bad_exception. Because bad_exception has not been specified in the exception specification of g(), unexpected() calls terminate(), which calls the function my terminate().

6. my_terminate() displays a message then calls abort(), which terminates the program.

Noe that the catch blocks following the second try block are not entered, because the exception was handled by my_unexpected() as an unexpected throw, not as a valid exception.

RELATED REFERENCES

- "unexpected()" on page 332
- "terminate()" on page 333
- "set_unexpected() and set_terminate()" on page 334

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SC09-4815-00

